



Consortium des Equipements
de Calcul Intensif
en Fédération Wallonie-Bruxelles

Introduction to Parallel Computing

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Agenda

1. General concepts, definitions, challenges
2. Hardware for parallel computing
3. Programming models
4. User tools

1.

General concepts

Why parallel? (simplified)

Speed up – Solve a problem faster

→ more processing power
(a.k.a. strong scaling)

Scale up – Solve a larger problem

→ more memory and network capacity
(a.k.a. weak scaling)

Scale out – Solve many problems

→ more storage capacity

Also: energy consumption is a cubic function of clock frequency so using 2
compute units is 8 times cheaper than using one unit with double the frequency

Parallelization involves:

- *decomposition* of the work
 - **distributing instructions** to processors
 - **distributing data** to memories
- *collaboration* of the workers
 - **synchronization** of the distributed work
 - **communication** of data

Parallelization involves:

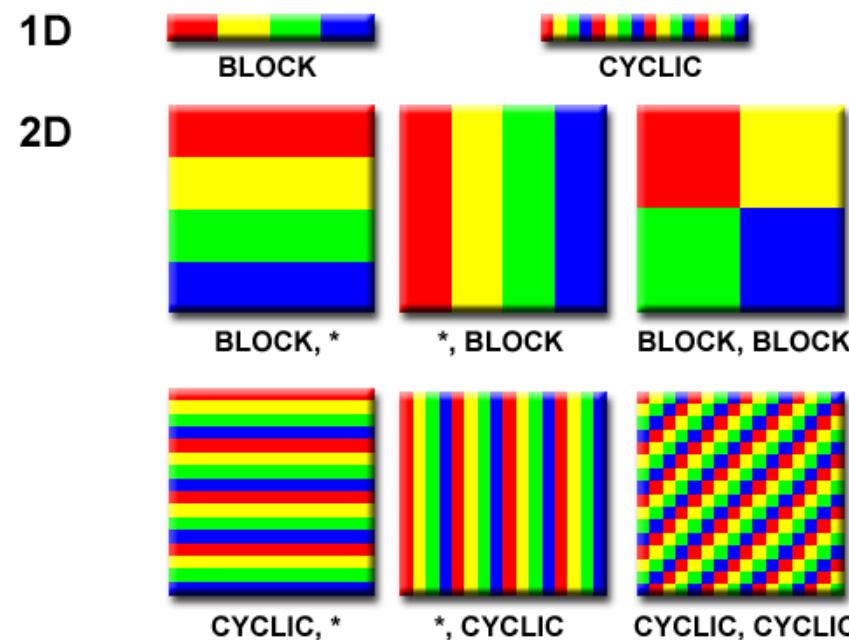
Decomposition of the work

- **Operation decomposition** : task-level parallelism
 - Multiple programs (functional decomposition)
 - Multiple instances of the same program (SPMD)
- **Data decomposition** : data-level parallelism

Parallelization involves:

Decomposition of the work

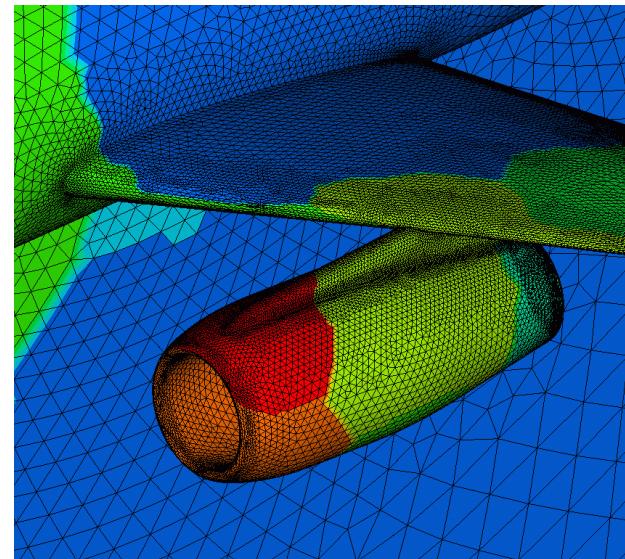
- **Operation decomposition** : task-level parallelism
- **Data decomposition** : data-level parallelism
 - **Block, cyclic**



Parallelization involves:

Decomposition of the work

- **Operation decomposition** : task-level parallelism
- **Data decomposition** : data-level parallelism
 - **Domain decomposition** : decomposition of work and data is done in a higher model, e.g. in the reality



Parallelization involves:

Collaboration of the workers

- **Synchronization of the workers**



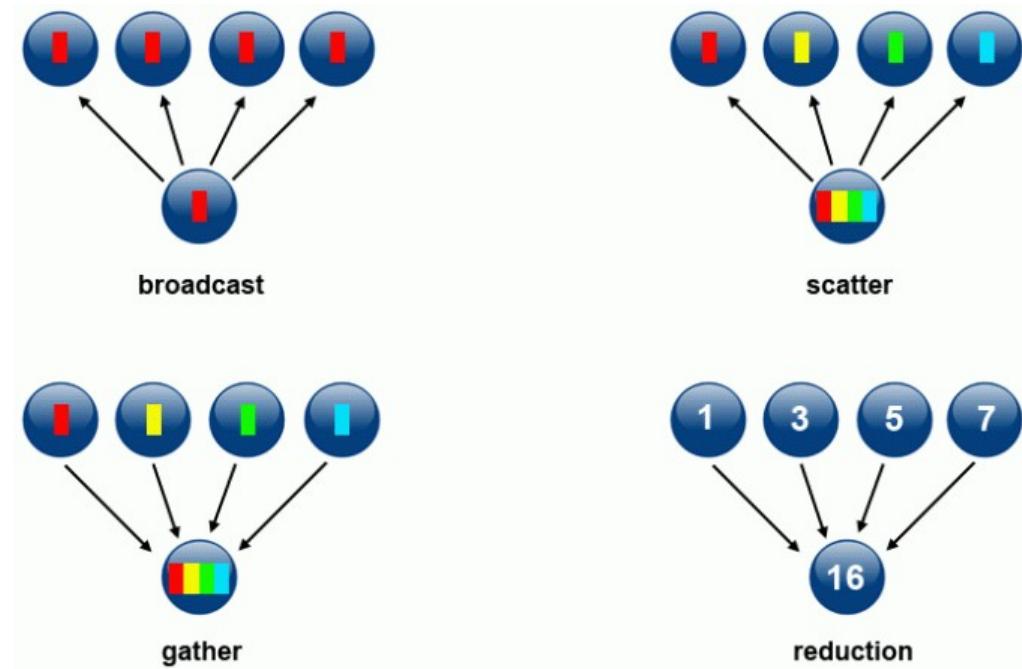
- **Synchronous** (SIMD) at the processor level ; the same processor instruction for each worker at any time ; (instruction level)
- **Fine-grained** parallelism : subtasks communicate many times per second (typically at the loop level)
- **Coarse-grained** parallelism : they do not communicate many times per second (typically function-call level)
- **Embarrassingly parallel** : they rarely or never have to communicate (asynchronous)

Parallelization involves:

Collaboration of the workers

- **Communication between workers**

- Point to point
- Broadcast
- Scatter
- Gather
- Reduction



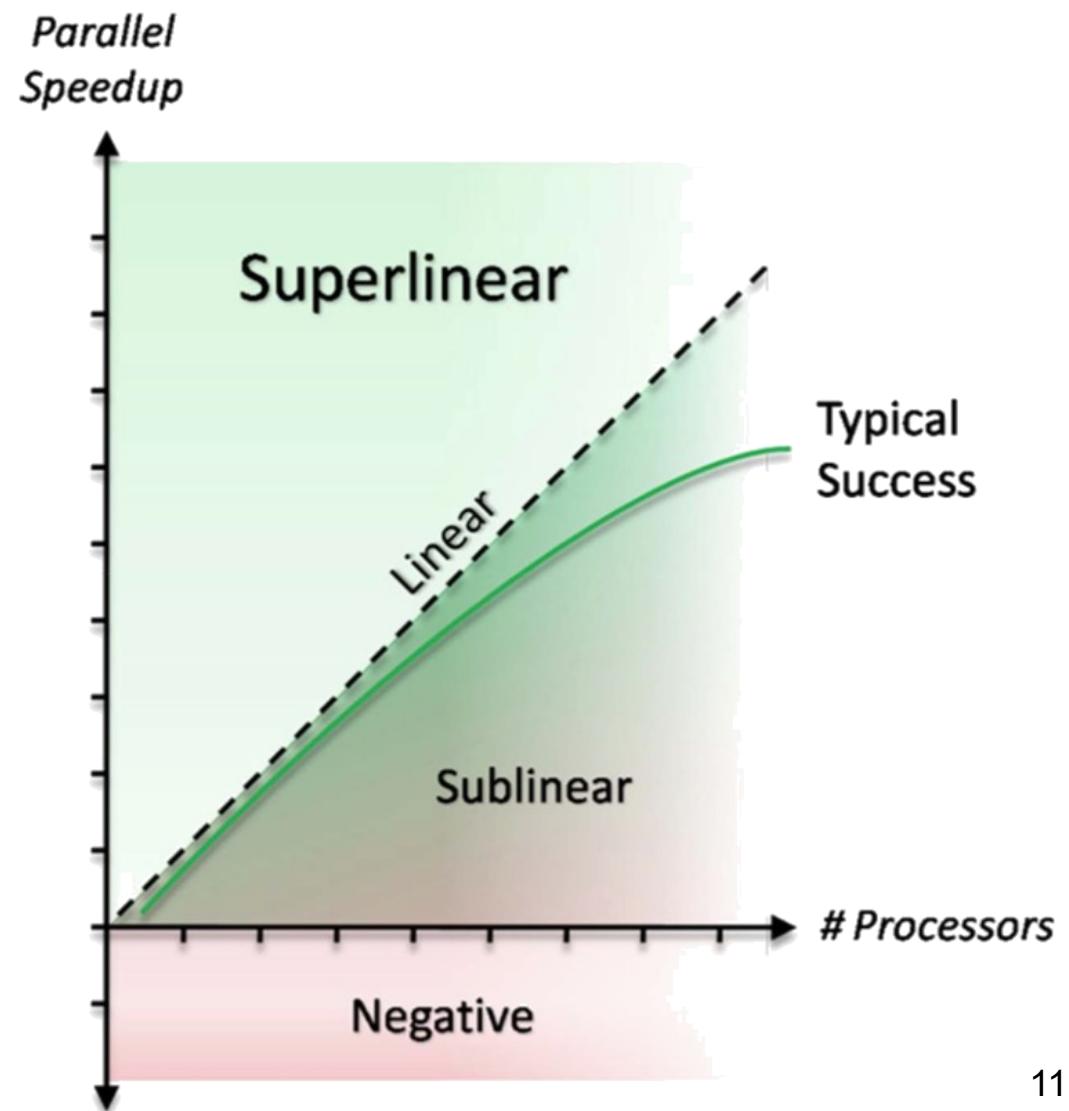
Does it work? Speedup, Efficiency, Scalability

$$\text{Speedup } S = \frac{T_S}{T_P}$$

Time for serial operations
Time for parallel operations

$$\text{Efficiency } E = \frac{S}{p} = \frac{T_S}{pT_p}$$

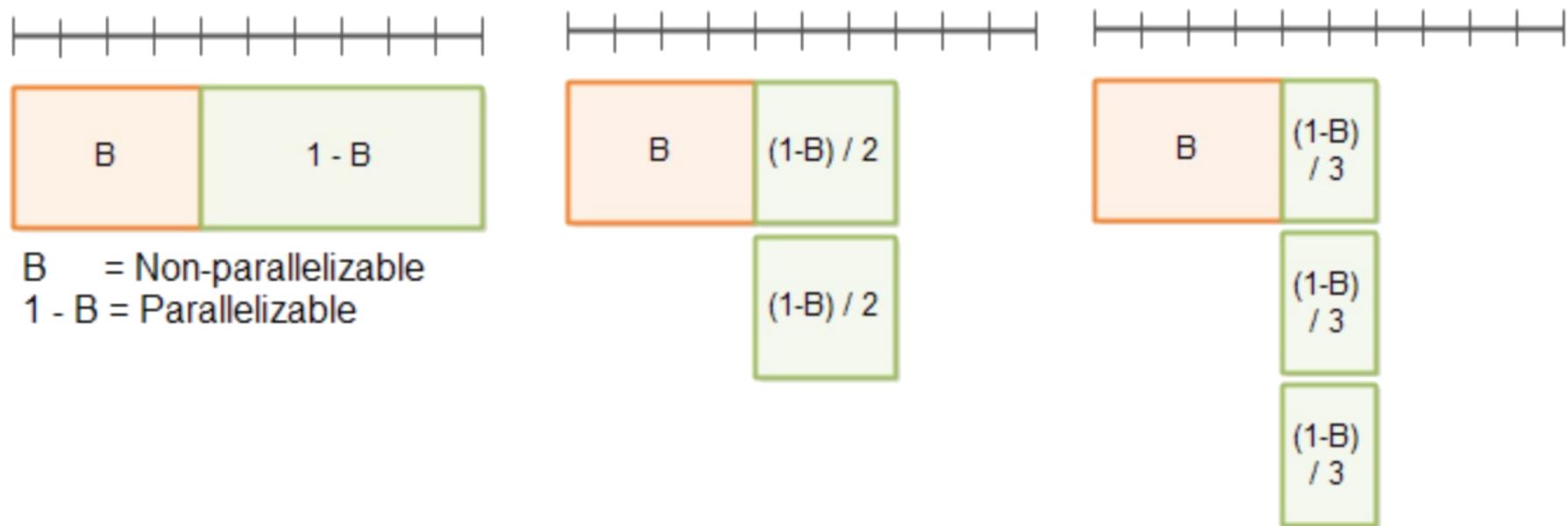
Number of processors



Why wouldn't it work?

Challenge 1: Amdahl's Law

Not all the work can be decomposed



In parallel computing, Amdahl's law is mainly used to predict the theoretical maximum speedup for programs using multiple processors.

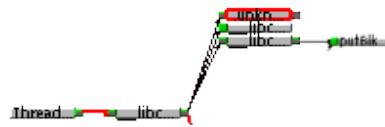
Why wouldn't it work?

Challenge 2: Parallel overhead

Collaboration means communication and extra work

```
void main (int argc, char *argv[])
{
    printf("Processor %d of %d: Hello World!\n",
}

```

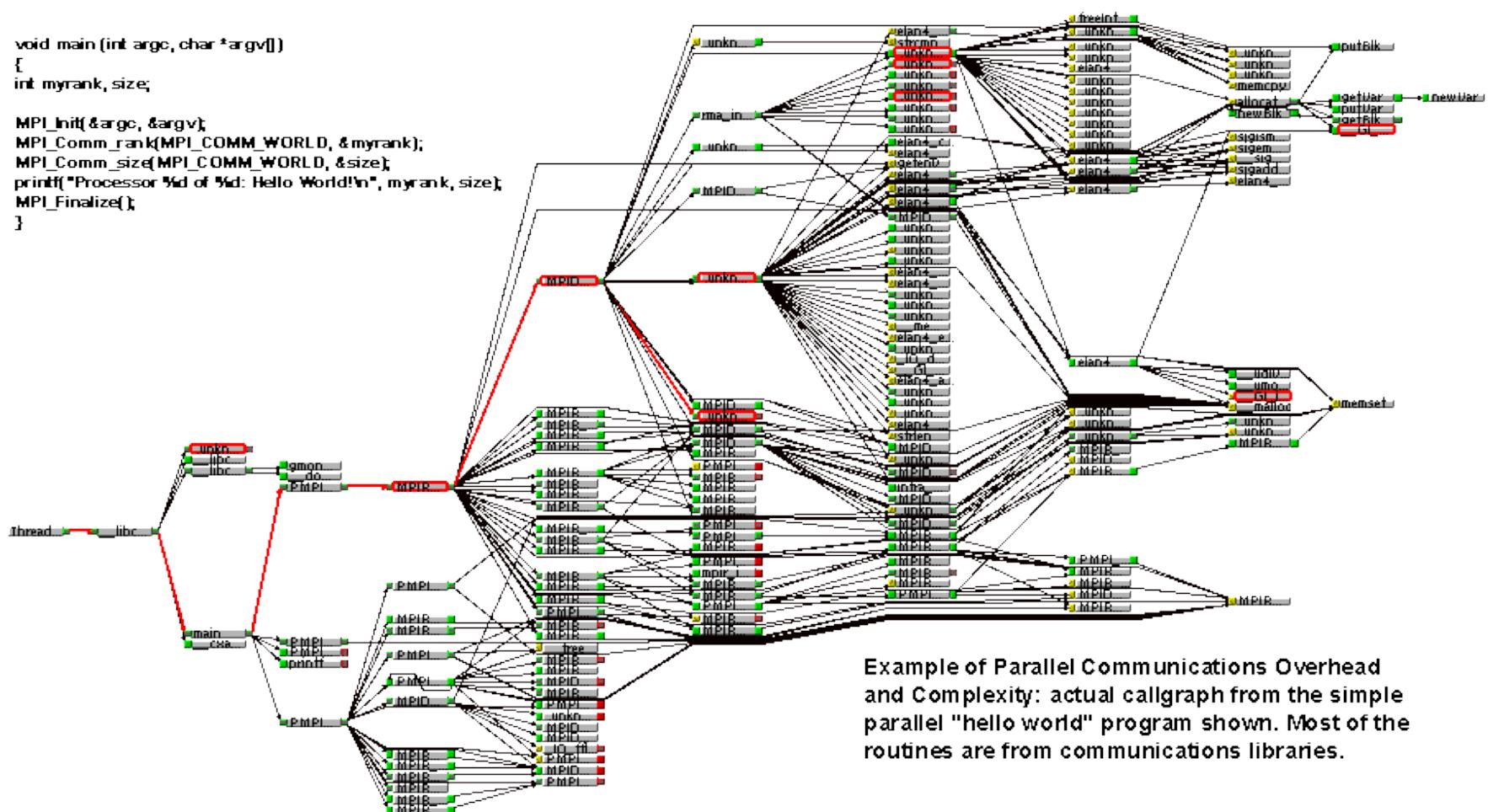


Why wouldn't it work?

Challenge 2: Parallel overhead

Collaboration means communication and extra work

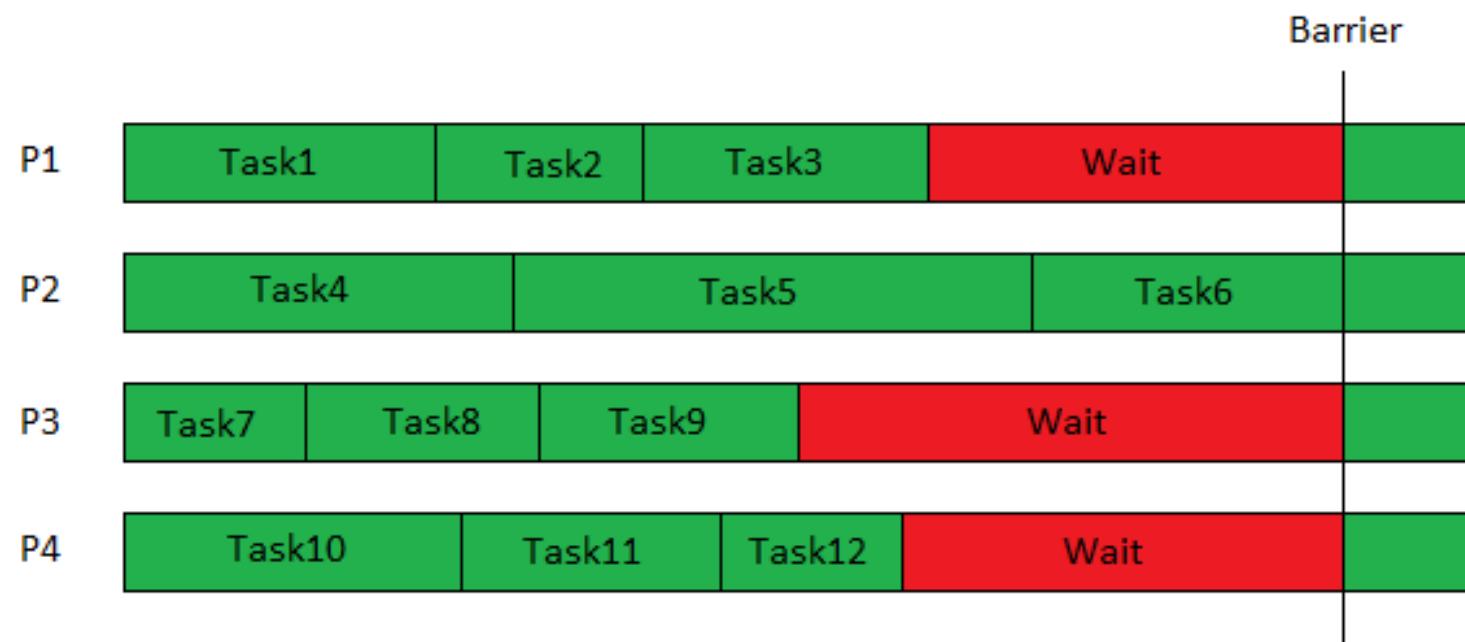
```
void main (int argc, char *argv[])
{
int myrank, size;
MPI_Init(&argc, &argv);
MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
MPI_Comm_size(MPI_COMM_WORLD, &size);
printf("Processor %d of %d: Hello World!\n", myrank, size);
MPI_Finalize();
}
```



Why wouldn't it work?

Challenge 3: Load imbalance

Parallelization is efficient only if every worker roughly has the same amount of work

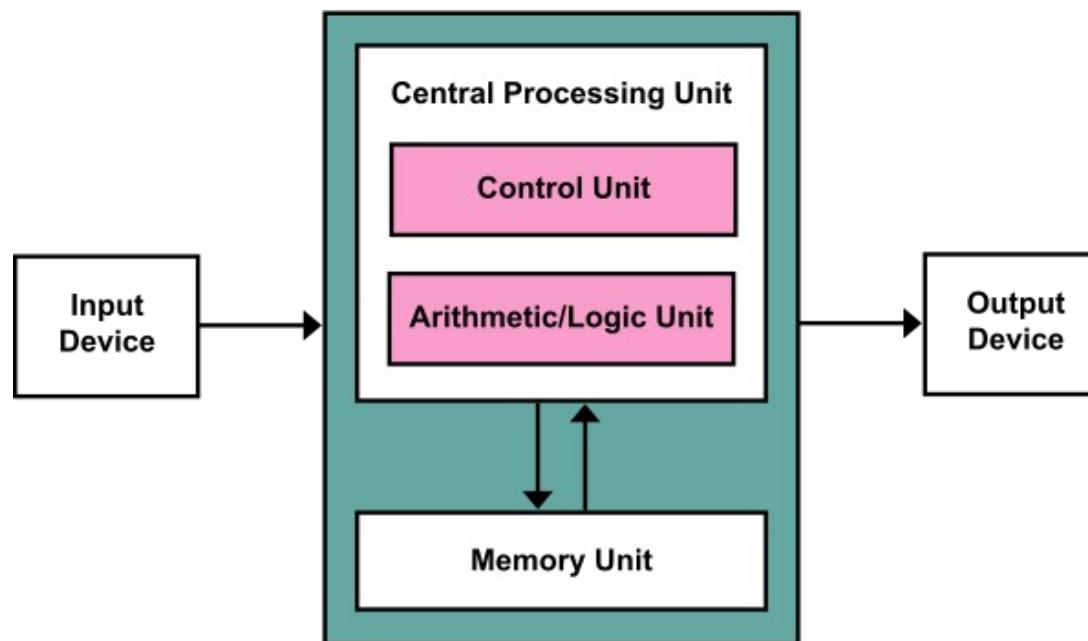


2.

Hardware for parallel computing

Von Neumann (serial) architecture

An abstract view of early computers



Parallelism at the CPU (core) level

An abstract view of modern CPUs

- Instruction-level parallelism (**ILP**)
 - Instruction pipelining
 - Out-of-order execution
 - Speculative execution
 - ...
- Single Instruction Multiple Data (**SIMD**)
- Simultaneous multithreading (**SMT**)

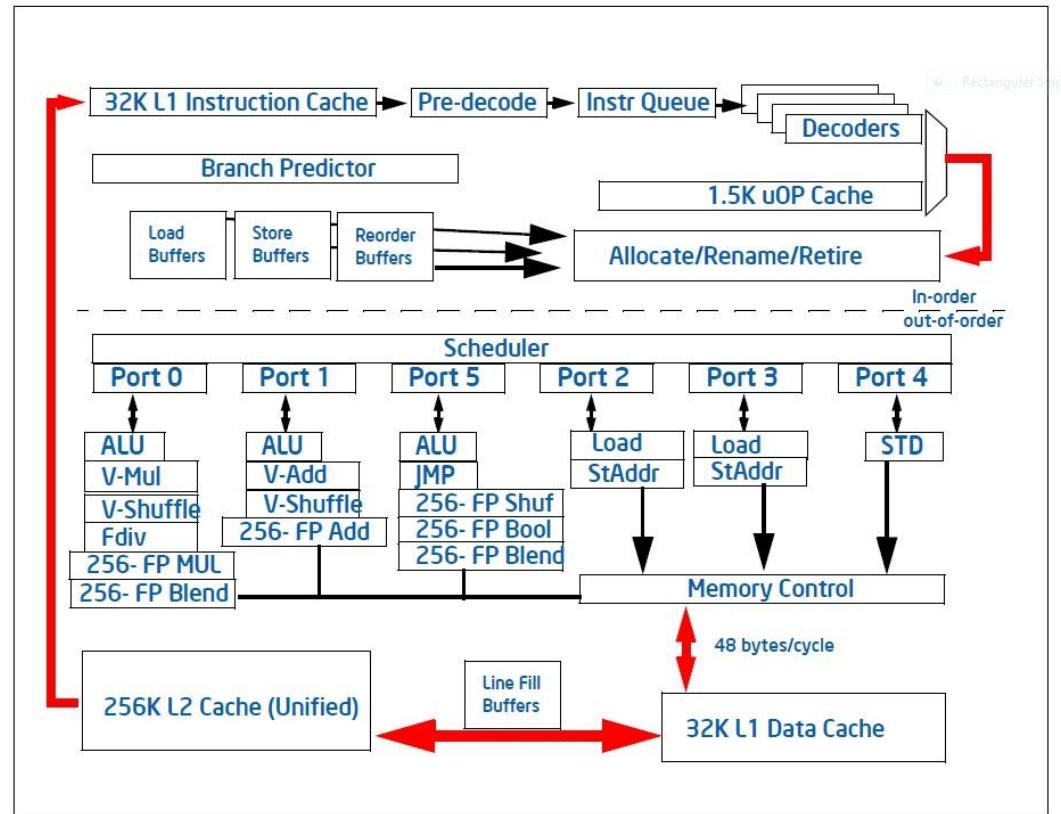
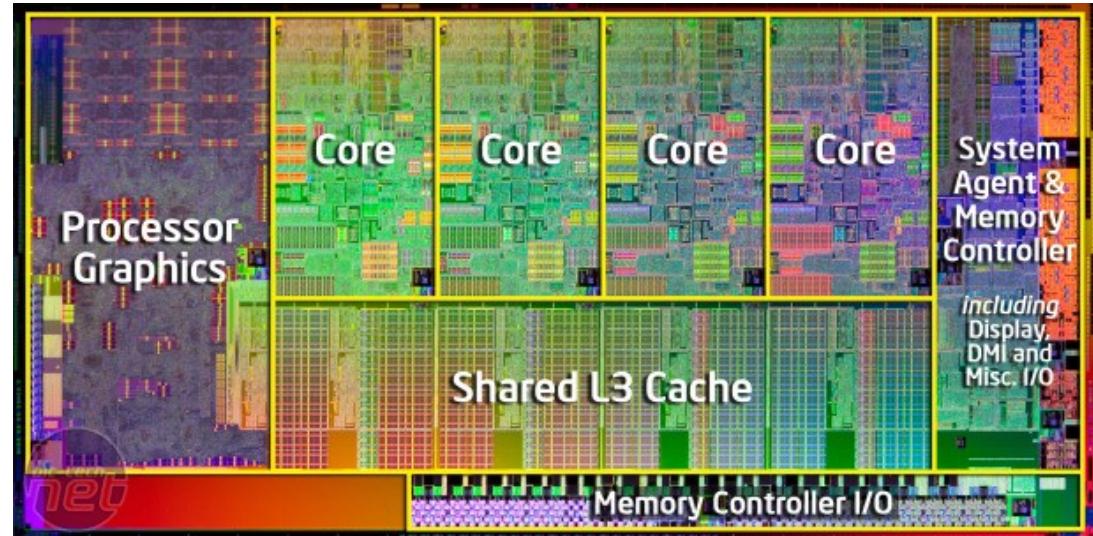


Figure 2-1. Intel microarchitecture code name Sandy Bridge Pipeline Functionality

Parallelism at the chip (socket) level



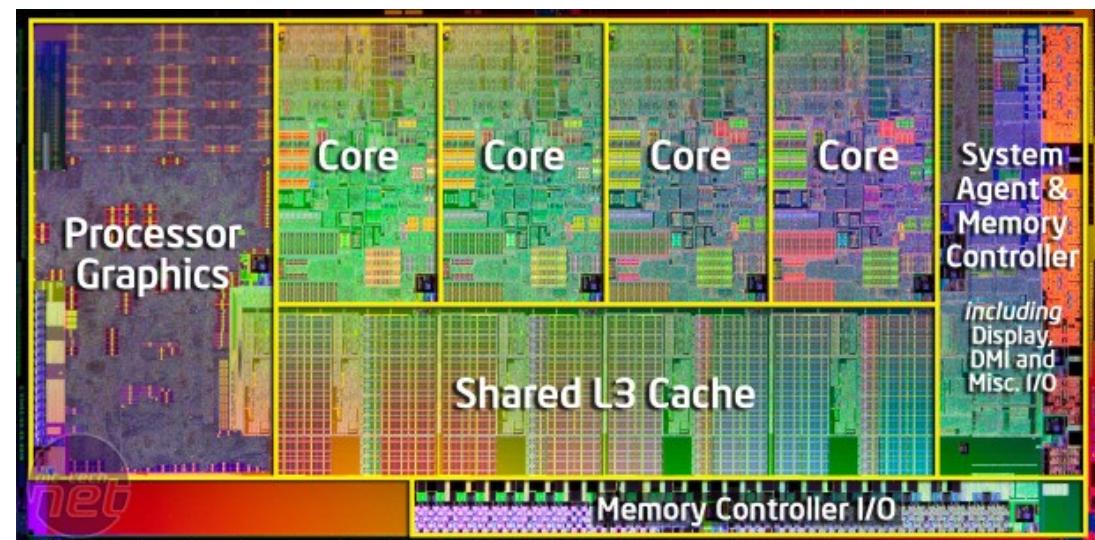
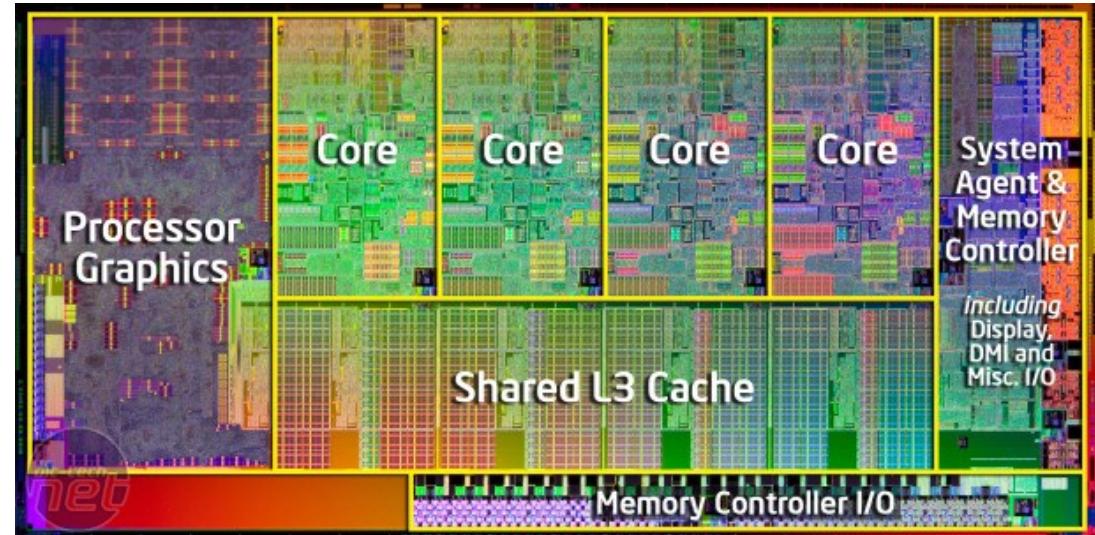
- **Multicore** parallelism



Parallelism at the computer level



- Multi-socket parallelism
 - SMP
 - NUMA
- Accelerators



Parallelism at the data center level



Multi-node parallelism



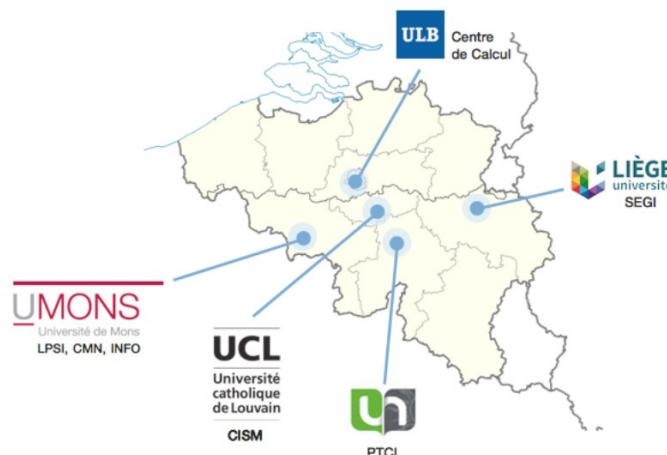
Parallelism at the data center level

Cluster computing



About

CÉCI is the 'Consortium des Équipements de Calcul Intensif'; a consortium of high-performance computing centers of [UCL](#), [ULB](#), [ULg](#), [UMons](#), and [UNamur](#). [Read more](#).



The common storage is functional!

Have you tried it yet? [More info...](#)

Latest News

SATURDAY, 23 SEPTEMBER 2017

A CECI user pictured in the ULiège news!

The ULiège website [published a story](#) (in French) about the work of Denis Baurain and his collaborators on the Tier-1 cluster Zenobe that lead to a publication in [Nature Ecology & Evolution](#).

TUESDAY, 01 AUGUST 2017

Ariel Lozano is the new CÉCI logisticien

We are happy to announce the hire of a new CECI logisticien: Ariel Lozano. Welcome Ariel!

Parallelism at the data center level

Cloud computing “someone else’s cluster”

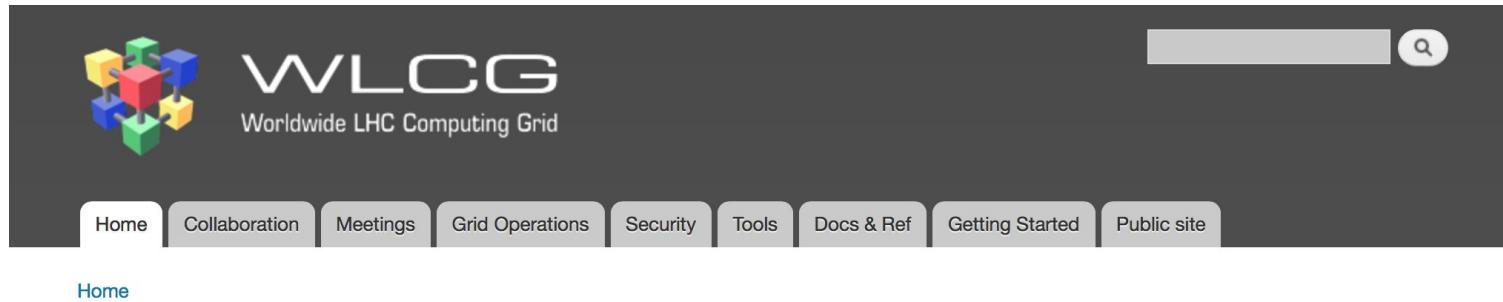


Google
Cloud Platform



Parallelism at the world level

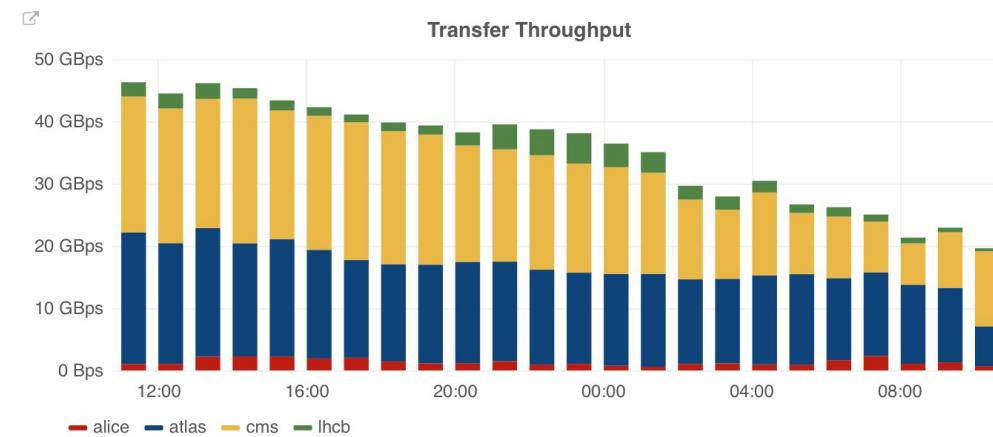
Grid computing – “cluster of clusters”



The screenshot shows the homepage of the Worldwide LHC Computing Grid (WLCG). The header features the WLCG logo (a cluster of colored cubes) and the text "WLCG Worldwide LHC Computing Grid". Below the header is a navigation bar with links: Home, Collaboration, Meetings, Grid Operations, Security, Tools, Docs & Ref, Getting Started, and Public site. The main content area has a "Home" link in blue. The main heading is "Welcome to the Worldwide LHC Computing Grid". Below it is a chart titled "Transfer Throughput" showing data for the last 24 hours. The chart is a stacked bar chart with four series: alice (red), atlas (dark blue), cms (yellow), and lhcb (green). The y-axis ranges from 0 Bps to 50 Gbps. The x-axis shows time intervals from 12:00 to 08:00. The total throughput fluctuates between approximately 20 and 45 Gbps. The bottom of the page contains a footer note about the WLCG project and a copyright notice.

Welcome to the Worldwide LHC Computing Grid

Last 24 hours



The Worldwide LHC Computing Grid (WLCG) project is a global collaboration of more than 170 computing centres in 42 countries, linking up national and international grid infrastructures.

About WLCG

What is WLCG? See our public site...

Hangout with CERN: LHC and Grid - the world is our calculator [☞](#)

Jobs

No jobs currently published.

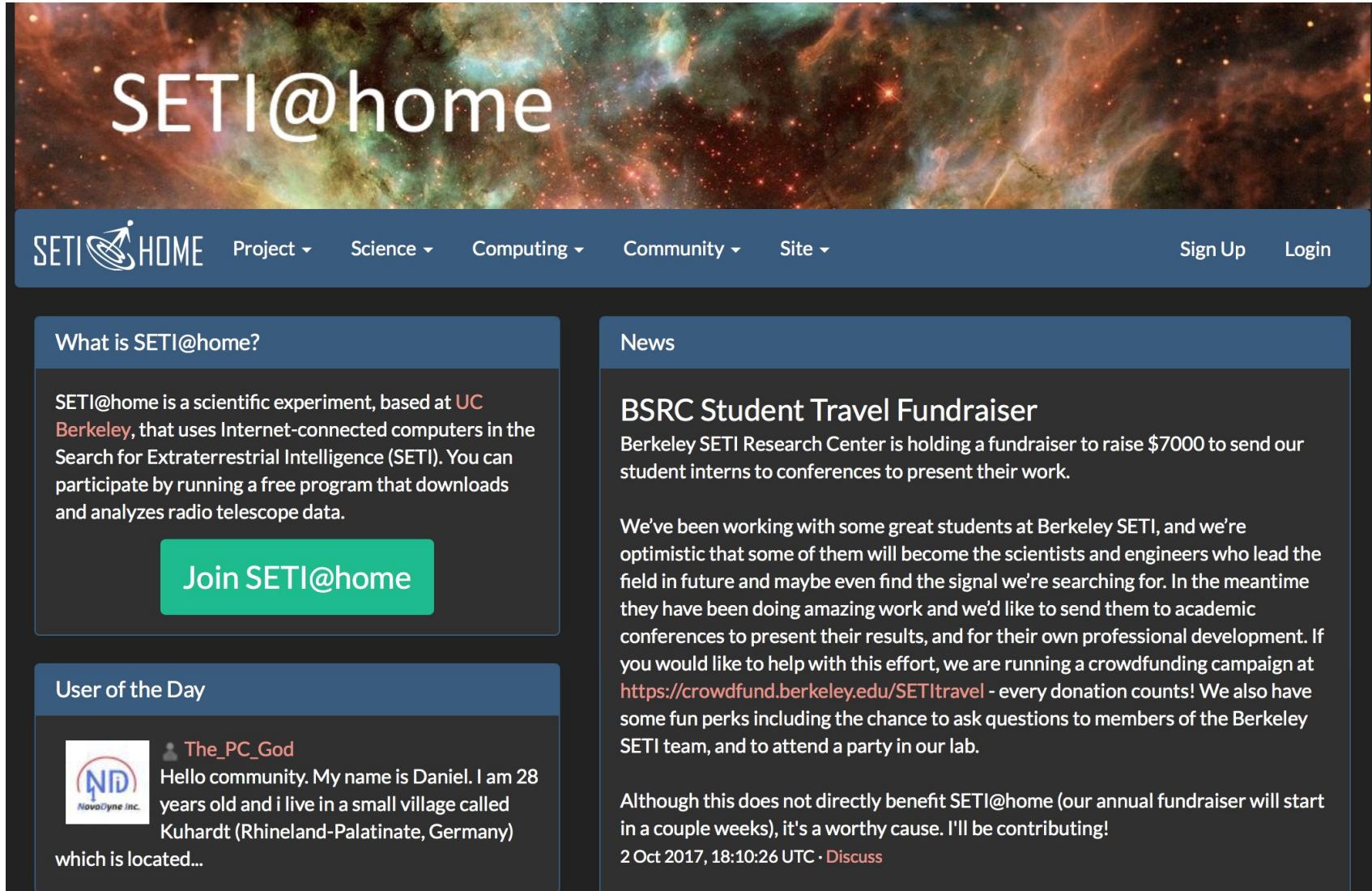
News

- [TEG Reports](#)
- [Computing Model Update available](#)
- [Technology Market Cost Trends](#)



Parallelism at the world level

Distributed computing – “no unused cycle”



The screenshot shows the homepage of the SETI@home website. The header features a background image of a colorful nebula. The main title "SETI@home" is prominently displayed in large white letters. Below the title is a navigation bar with links for "Project", "Science", "Computing", "Community", "Site", "Sign Up", and "Login".

What is SETI@home?

SETI@home is a scientific experiment, based at [UC Berkeley](#), that uses Internet-connected computers in the Search for Extraterrestrial Intelligence (SETI). You can participate by running a free program that downloads and analyzes radio telescope data.

[Join SETI@home](#)

User of the Day

 **The_PC_God**
Hello community. My name is Daniel. I am 28 years old and I live in a small village called Kuhardt (Rhineland-Palatinate, Germany) which is located...

News

BSRC Student Travel Fundraiser
Berkeley SETI Research Center is holding a fundraiser to raise \$7000 to send our student interns to conferences to present their work.

We've been working with some great students at Berkeley SETI, and we're optimistic that some of them will become the scientists and engineers who lead the field in future and maybe even find the signal we're searching for. In the meantime they have been doing amazing work and we'd like to send them to academic conferences to present their results, and for their own professional development. If you would like to help with this effort, we are running a crowdfunding campaign at <https://crowdfund.berkeley.edu/SETItravel> - every donation counts! We also have some fun perks including the chance to ask questions to members of the Berkeley SETI team, and to attend a party in our lab.

Although this does not directly benefit SETI@home (our annual fundraiser will start in a couple weeks), it's a worthy cause. I'll be contributing!

2 Oct 2017, 18:10:26 UTC · [Discuss](#)

3.

Programming paradigms
(logic of work sharing and organizing)

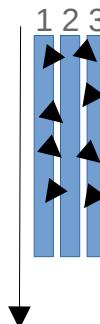
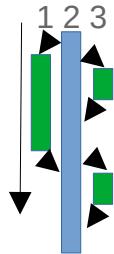
Programming models
(software libraries and APIs)

Is parallelization automagic?

- **ILP**: yes, by the CPU and/or the compiler
- **SIMD**: mostly, by the compiler
- **Intra-node**: can be if the library/software you use is designed for it
- **GPUs**: can be if the library/software you use is designed for it
- **Inter-node**: never automagic.

Main parallel programming paradigms

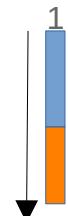
- **Task-farming:**
 - Master program distributes work to worker programs (*leader/follower*); or
 - Worker programs pick up tasks from pool (*work stealing*).
- **SPMD** (Single program multiple data)
 - A single program that contains both the logic for distributing work and computing
 - Multiple instances are started and “linked” together
 - Instances are identified with a distinct index



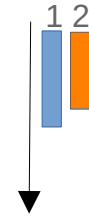
Other parallel programming paradigms

- **MPMD (Multiple program multiple data)**
- **Pipelining** : workers take care of a subtask in the processing chain and pass the intermediate result to the next worker
- **Divide and Conquer** :
 - workers are spawned at need and report their result to the parent
 - Speculative parallelism : workers are spawned and result possibly discarded

```
CPU1: If (very_long_computation())
CPU1: then
CPU1:   do A
CPU1: else
CPU1:   do B
```



```
CPU1: res=very_long_computation()
CPU2: do A
CPU3: do B
CPU1: if (res) discard B else discard A
```



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Main programming models

- Single computer:
 - **CPUs**: PThreads, *OpenMP*, TBB, OpenCL
 - **Accelerators**: CUDA, OpenCL, *OpenAcc/OpenMP*, SYCL, Hipp, ROCm
- Multi-computer:
 - **Clusters**:
 - Message passing: *MPI*, PVM
 - PGAS: CoArray Fortran, UPC, Global Arrays
 - **Clouds**: MapReduce, Spark RDD
 - **Distributed computing**: BOINC

Linux program starting process

Code (program.c)

Text file

Compiler

Binary (program.exe)

Executable file

Loader

Process (PID 1235)

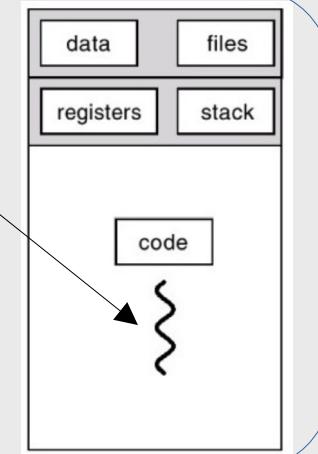
Running instance

```
#include <stdio.h>
```

```
int main(void)
{
    printf("Hello, World!\n");
}
```

Computer

One execution thread



Code (program.c)

Text file

Compiler

Binary (program.exe)

Executable file

Loader

Process (PID 1235)

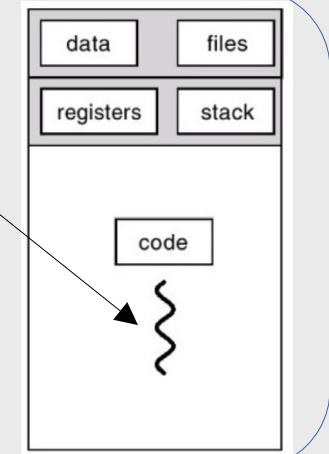
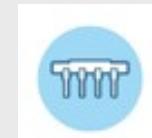
Running instance

```
#include <stdio.h>
```

```
int main(void)
{
    printf("Hello, World!\n");
}
```

Computer

One execution thread
is assigned by the OS
one CPU core



Code (program.c)

Text file

Compiler

Binary (program.exe)

Executable file

Loader, called multiple times

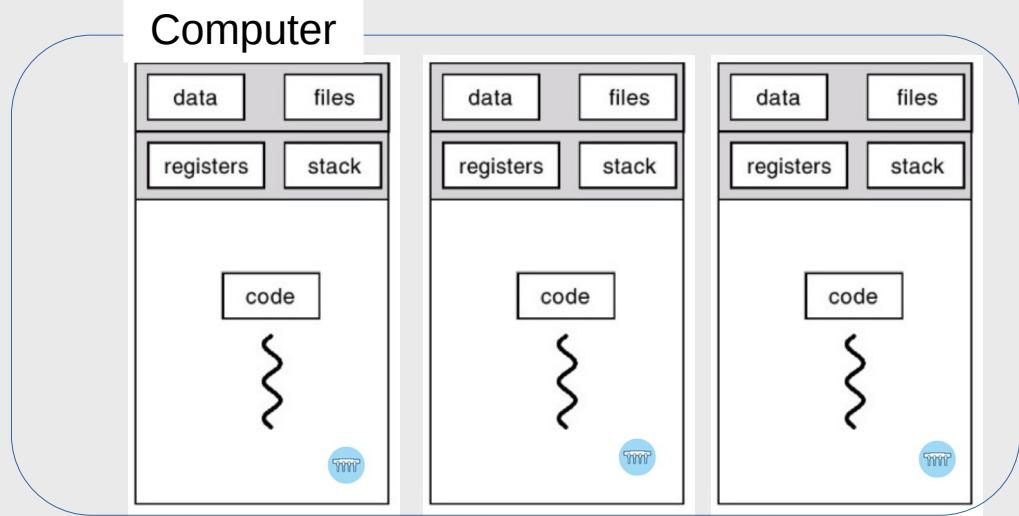
Multiple independent processes

Running instances

```
#include <stdio.h>
```

```
int main(void)
{
    printf("Hello, World!\n");
}
```

Computer



Code (program.c)

Text file

Compiler

Binary (program.exe)

Executable file

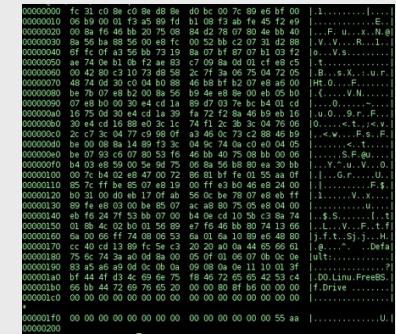
Loader,
called multiple times
with SSH on multiple
servers

Multiple **independent**
processes on multiple
servers

Running instances

```
#include <stdio.h>

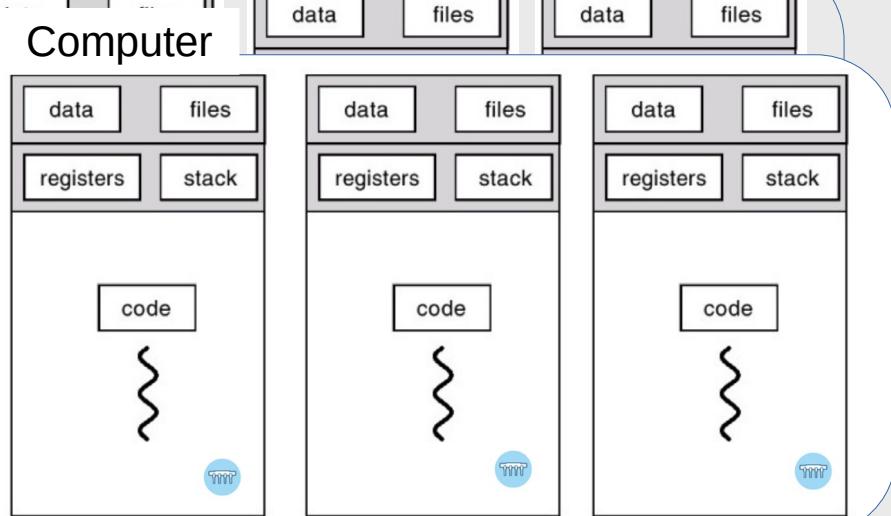
int main(void)
{
    printf("Hello, World!\n");
}
```



A screenshot of a hex editor showing the binary representation of the C code. The code is "Hello, World!\n" in red. The hex editor shows the memory address (00000000 to 00000020), the ASCII characters, and the corresponding hex values. The text is preceded by a null byte (00) and followed by a carriage return (0D) and a line feed (0A).

Computer

Computer



Forking Code

Text file

Compiler

Single binary

Executable file

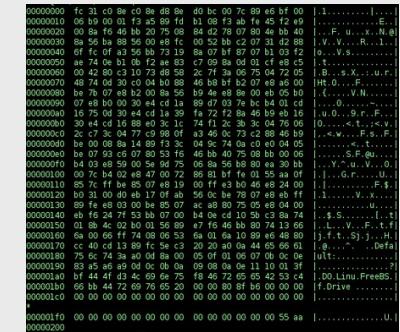
Loader, called once

Multiple Processes
with parent/child
relationship

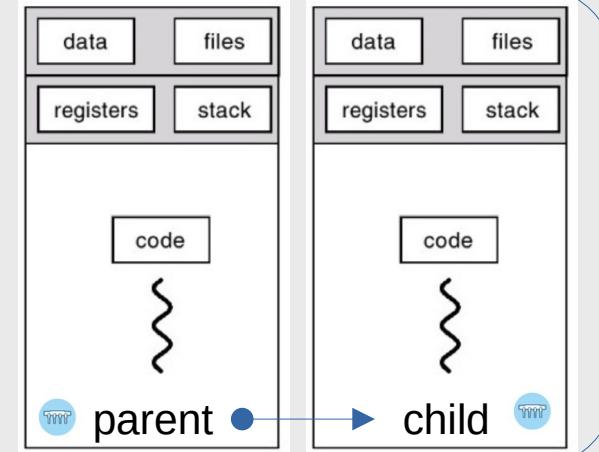
Running instances

```
#include <stdio.h>
#include <sys/types.h>;
#include <unistd.h>;
int main()
{
    // make two process which run same
    // program after this instruction
    fork();

    printf("Hello world!\n");
    return 0;
}
```



Computer



IPC – Inter-process communication

Multithreaded Code

Text file

Compiler

Single binary

Executable file

Loader, called once

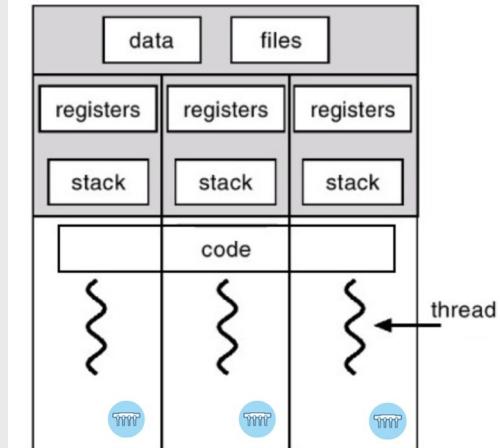
Single process with
multiple threads
(multithreaded – shared memory)

Running instance

```
void print_message_function( void *ptr );  
  
main()  
{  
    pthread_t thread1, thread2;  
    char *message1 = "Hello";  
    char *message2 = "World";  
  
    pthread_create( &thread1, pthread_attr_default,  
                    (void*)&print_message_function, (void*) message1);  
    pthread_create(&thread2, pthread_attr_default,  
                    (void*)&print_message_function, (void*) message2);  
  
    exit(0);  
}  
  
void print_message_function( void *ptr )  
{  
    char *message;  
    message = (char *) ptr;  
    printf("%s ", message);  
}
```



Computer



Multithreaded Code

Text file

Compiler

Single binary

Executable file

Loader, called once

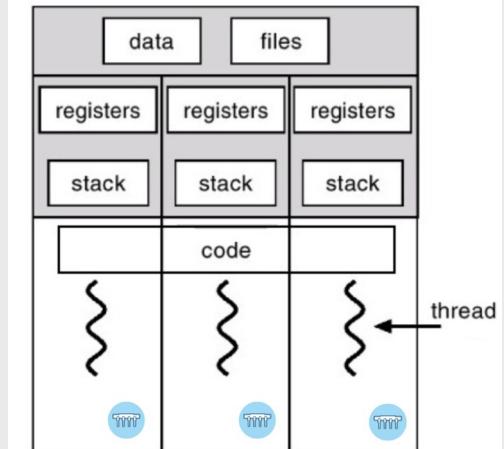
Single process with
multiple threads
(multithreaded – shared memory)

Running instance

```
void print_message_function( void *ptr );  
  
main()  
{  
    pthread_t thread1, thread2;  
    char *message1 = "Hello  
    char *message2 = "World  
  
    pthread_create( &thread1, pthread_attr_default,  
                    (void*)print_message_function, (void*) message1);  
    pthread_create(&thread2, pthread_attr_default,  
                    (void*)print_message_function, (void*) message2);  
  
    exit(0);  
}  
  
void print_message_function( void *ptr )  
{  
    char *message;  
    message = (char *) ptr;  
    printf("%s ", message);  
}
```



Computer



Code (program.c)

Text file

Compiler

```
int main() {
    struct sockaddr_in client, server;
    int lfd, n, confd;
    char r_buff[100] = "", s_buff[100] = "";
    lfd = socket(AF_INET, SOCK_STREAM, 0);
    server.sin_family = AF_INET;
    server.sin_port = 2000;
    server.sin_addr.s_addr = inet_addr("127.0.0.1");
    bind(lfd, (struct sockaddr *)&server, sizeof server);
    listen(lfd, 1);

    n = sizeof client;
    confd = accept(lfd, (struct sockaddr *)&client, &n);
    while (1) {
        recv(confd, r_buff, sizeof r_buff, 0);
        printf("\n[client] %s", r_buff);
        if (strcmp(r_buff, "exit") == 0)
            break;

        printf("\n[server: ");
        gets(s_buff);
        send(confd, s_buff, sizeof s_buff, 0);
        if (strcmp(s_buff, "exit") == 0)
            break;
        printf("\n");
    }
    printf("\n");
}
```

Binary (program.exe)

Executable file

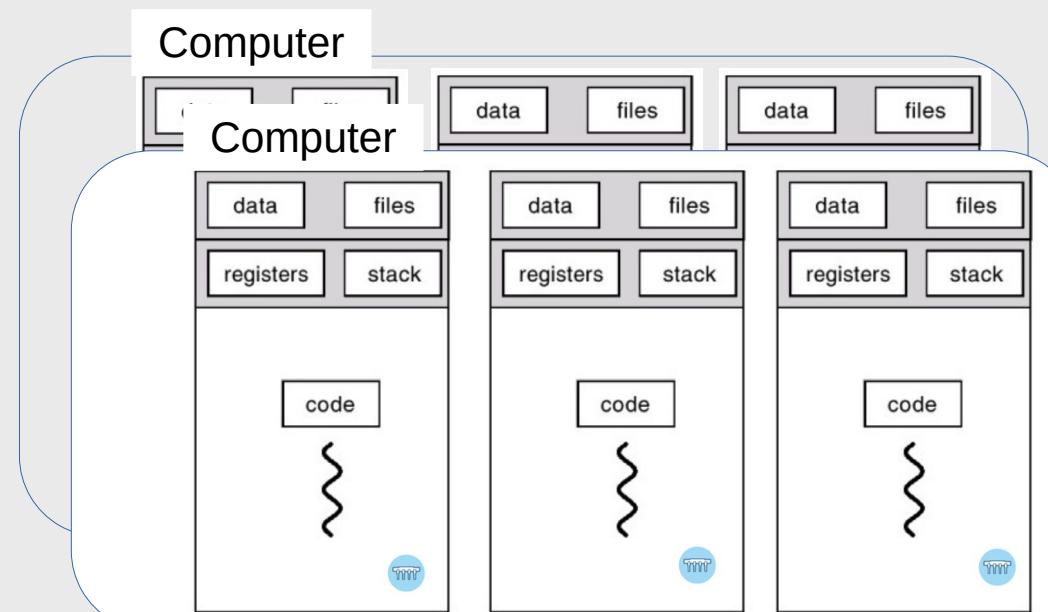


Hex dump of the binary program.exe, showing assembly code and data segments. The dump includes labels, assembly instructions, and memory addresses.

Loader,
called multiple times
with SSH on multiple
servers

Multiple **connected**
processes on multiple
servers

Running instances



Code (program.c)

Text file

Compiler

Binary (program.exe)

Executable file

Loader,
called multiple times
with SSH on multiple
servers

Multiple connected processes on multiple servers

Running instances

Running instances

```
int main() {
    struct sockaddr_in client, server;
    int lfd, n, confd;
    char r_buff[100] = "", s_buff[100] = "";

    lfd = socket(AF_INET, SOCK_STREAM, 0);
    server.sin_family = AF_INET;
    server.sin_port = 2000;
    server.sin_addr.s_addr = inet_addr("27.0.0.1");
    bind(lfd, (struct sockaddr *)&server, sizeof(server));
    listen(lfd, 1);

    n = sizeof(client);
    confd = accept(lfd, (struct sockaddr *)&client, &n);
    if (strcmp(r_buff, "exit") == 0)
        break;

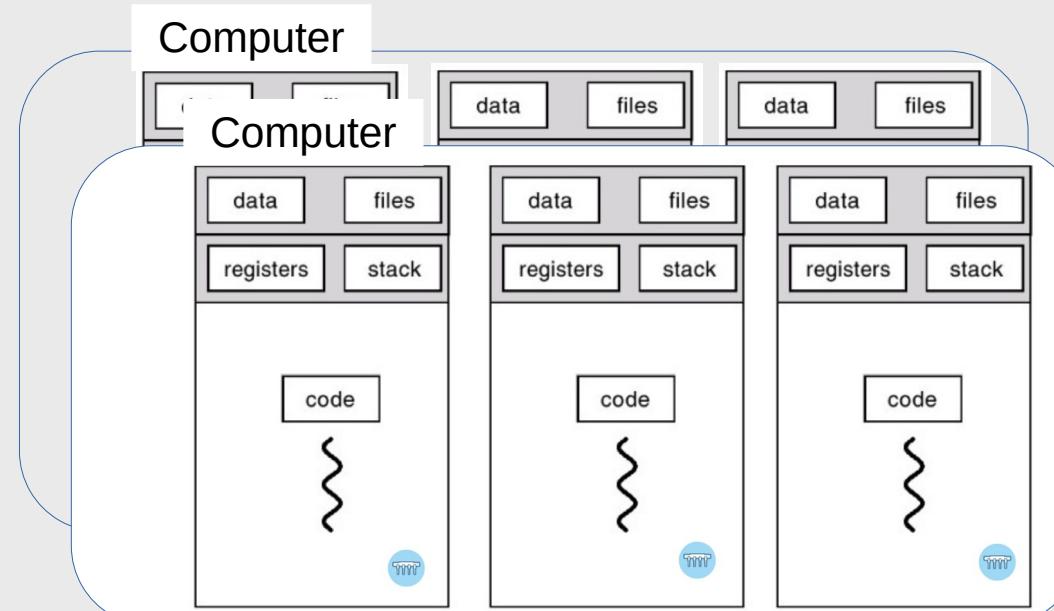
    while (1) {
        recvfrom(confd, r_buff, sizeof(r_buff), 0);
        printf("\n[client] %s", r_buff);
        if (strcmp(r_buff, "exit") == 0)
            break;

        printf("\nserver: ");
        gets(s_buff);
        send(confd, s_buff, sizeof(s_buff), 0);
        if (strcmp(s_buff, "exit") == 0)
            break;
    }
    printf("\n");
}

int main() {
    struct sockaddr_in server;
    int lfd;
    char r_buff[100] = "", s_buff[100] = "";

    lfd = socket(AF_INET, SOCK_STREAM, 0);
    server.sin_family = AF_INET;
    server.sin_port = 2000;
    server.sin_addr.s_addr = inet_addr("127.0.0.1");
    connect(lfd, (struct sockaddr *)&server, sizeof(server));

    while (1) {
        if (recv(lfd, r_buff, sizeof(r_buff), 0) < 0)
            break;
        if (strcmp(r_buff, "end") == 0)
            break;
        send(lfd, s_buff, sizeof(s_buff), 0);
        if (strcmp(s_buff, "end") == 0)
            break;
    }
    printf("\n");
}
```



xargs, split, GNU parallel → start multiple independent processes

OpenMP → write multithreaded programs

MPI → write multiprocess connected programs

4.

User tools that GNU/Linux offers

- 4.1 Parallelized tools
- 4.2 Job control and parallel processes
- 4.3 Basic tools
- 4.4 GNU Parallel

4.1. Parallelized utilities

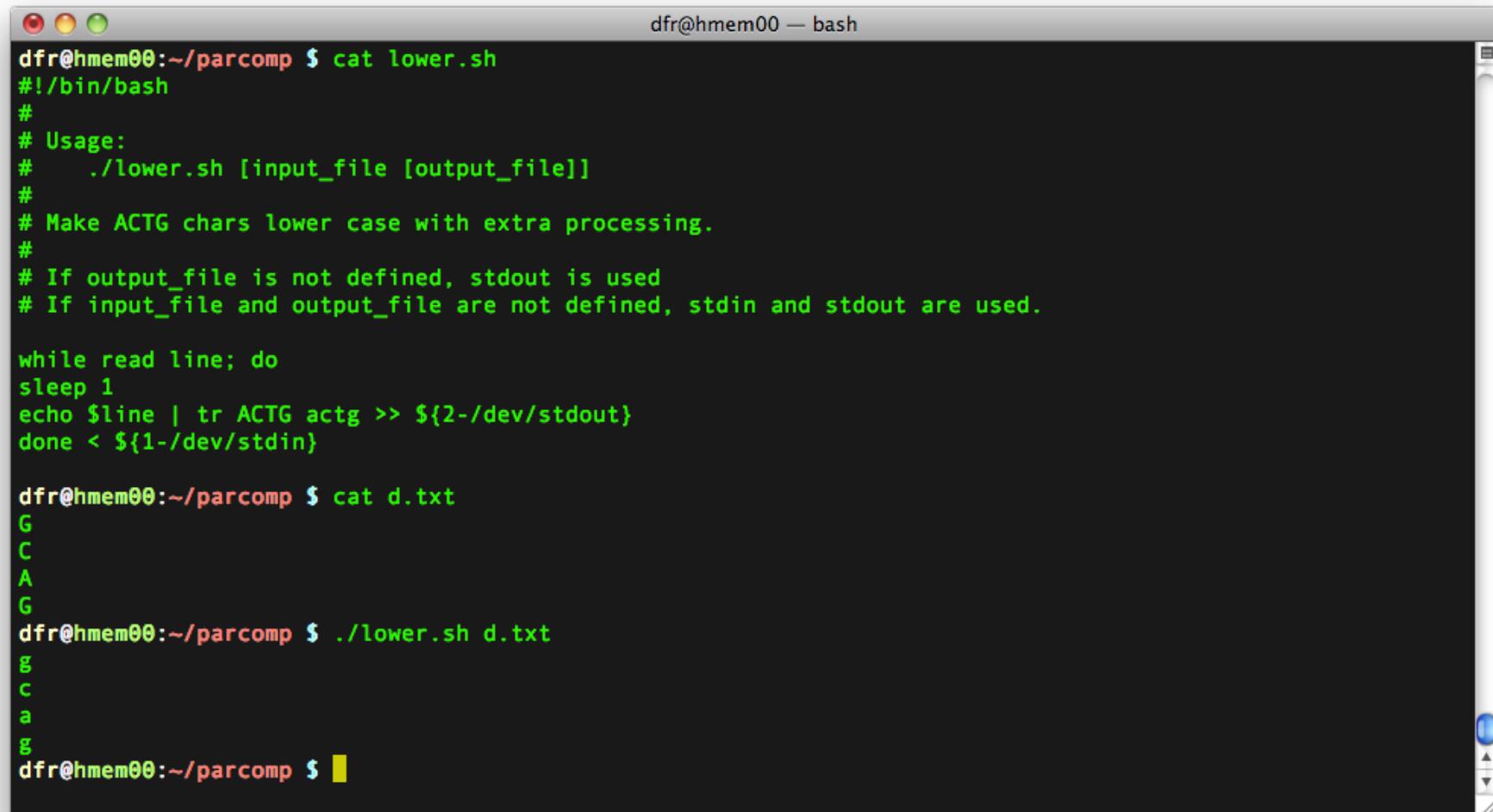
Some tools have a parallelized counterpart, or parallel options. Examples:

serial		parallel
gzip	↔	pigz
grep	↔	ripgrep, singrep
ssh	↔	clustershell
sort	↔	sort --parallel
scp	↔	bbcp
bc	↔	bcx
...		...

You might have to install them by yourself if they are not present on the clusters

4.2. Job control & Parallel processes in Bash

Consider the following example program



```
dfr@hmem00:~/parcomp $ cat lower.sh
#!/bin/bash
#
# Usage:
#   ./lower.sh [input_file [output_file]]
#
# Make ACTG chars lower case with extra processing.
#
# If output_file is not defined, stdout is used
# If input_file and output_file are not defined, stdin and stdout are used.

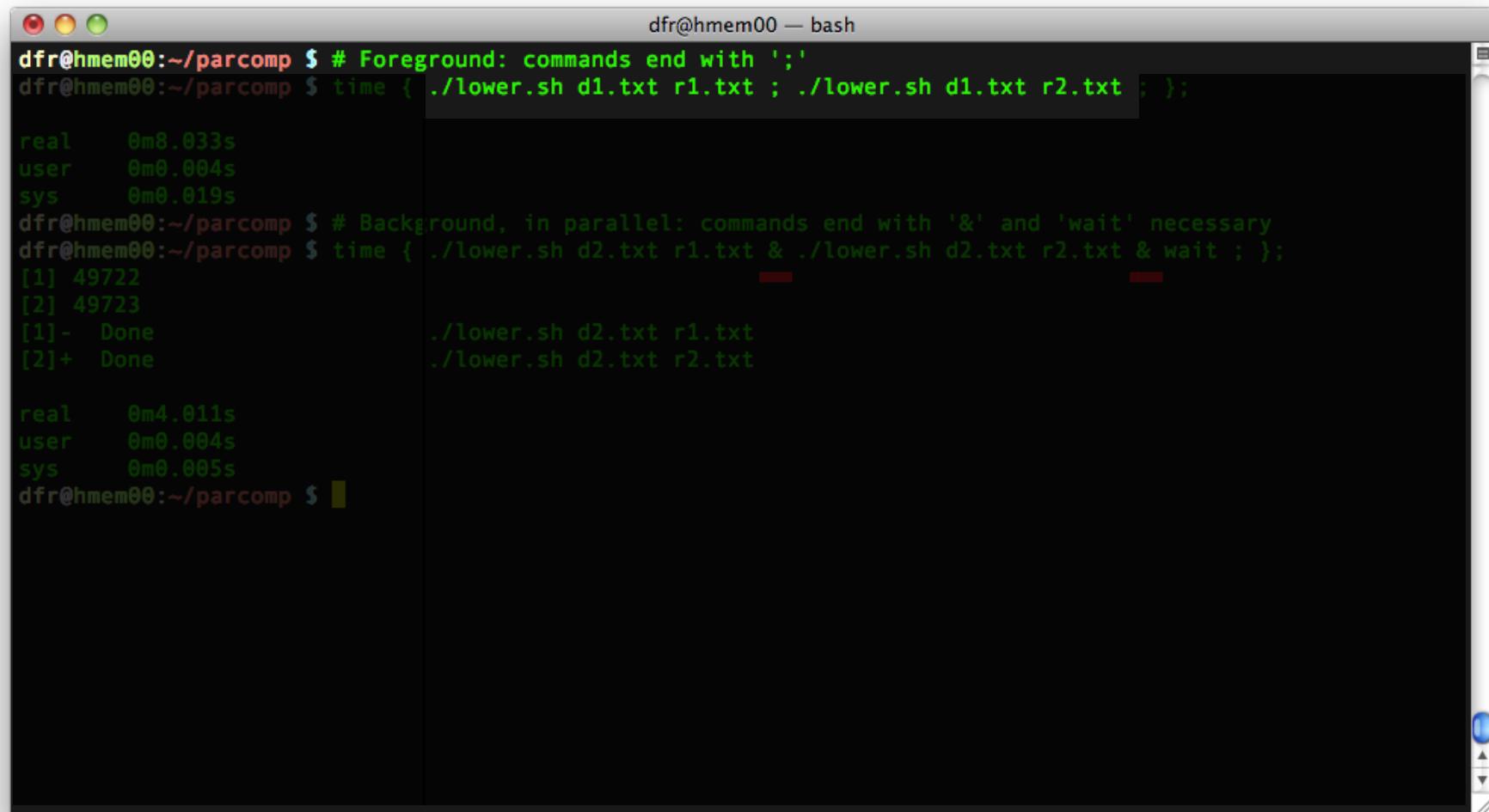
while read line; do
sleep 1
echo $line | tr ACTG actg >> ${2-/dev/stdout}
done < ${1-/dev/stdin}

dfr@hmem00:~/parcomp $ cat d.txt
G
C
A
G
dfr@hmem00:~/parcomp $ ./lower.sh d.txt
g
c
a
g
dfr@hmem00:~/parcomp $
```

It is written in Bash and just transforms some upper case letters to lower case

4.2. Job control & Parallel processes in Bash

Run the program twice

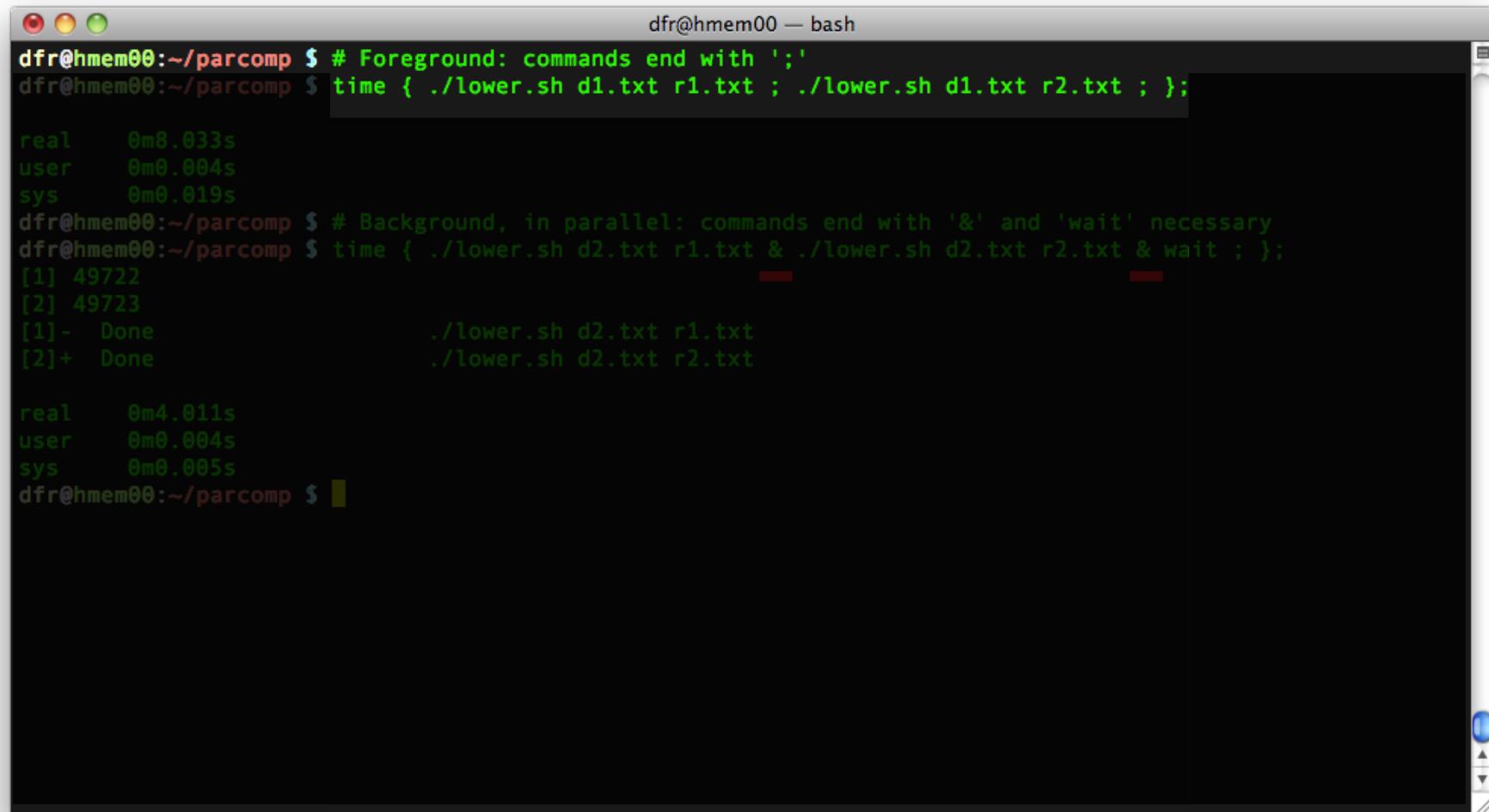


```
dfr@hmem00:~/parcomp $ # Foreground: commands end with ';'
dfr@hmem00:~/parcomp $ time { ./lower.sh d1.txt r1.txt ; ./lower.sh d1.txt r2.txt ; };
real    0m8.033s
user    0m0.004s
sys     0m0.019s
dfr@hmem00:~/parcomp $ # Background, in parallel: commands end with '&' and 'wait' necessary
dfr@hmem00:~/parcomp $ time { ./lower.sh d2.txt r1.txt & ./lower.sh d2.txt r2.txt & wait ; };
[1] 49722
[2] 49723
[1]-  Done                  ./lower.sh d2.txt r1.txt
[2]+  Done                  ./lower.sh d2.txt r2.txt

real    0m4.011s
user    0m0.004s
sys     0m0.005s
dfr@hmem00:~/parcomp $
```

4.2. Job control & Parallel processes in Bash

Run the program twice and measure the time it takes



The screenshot shows a terminal window with the following session:

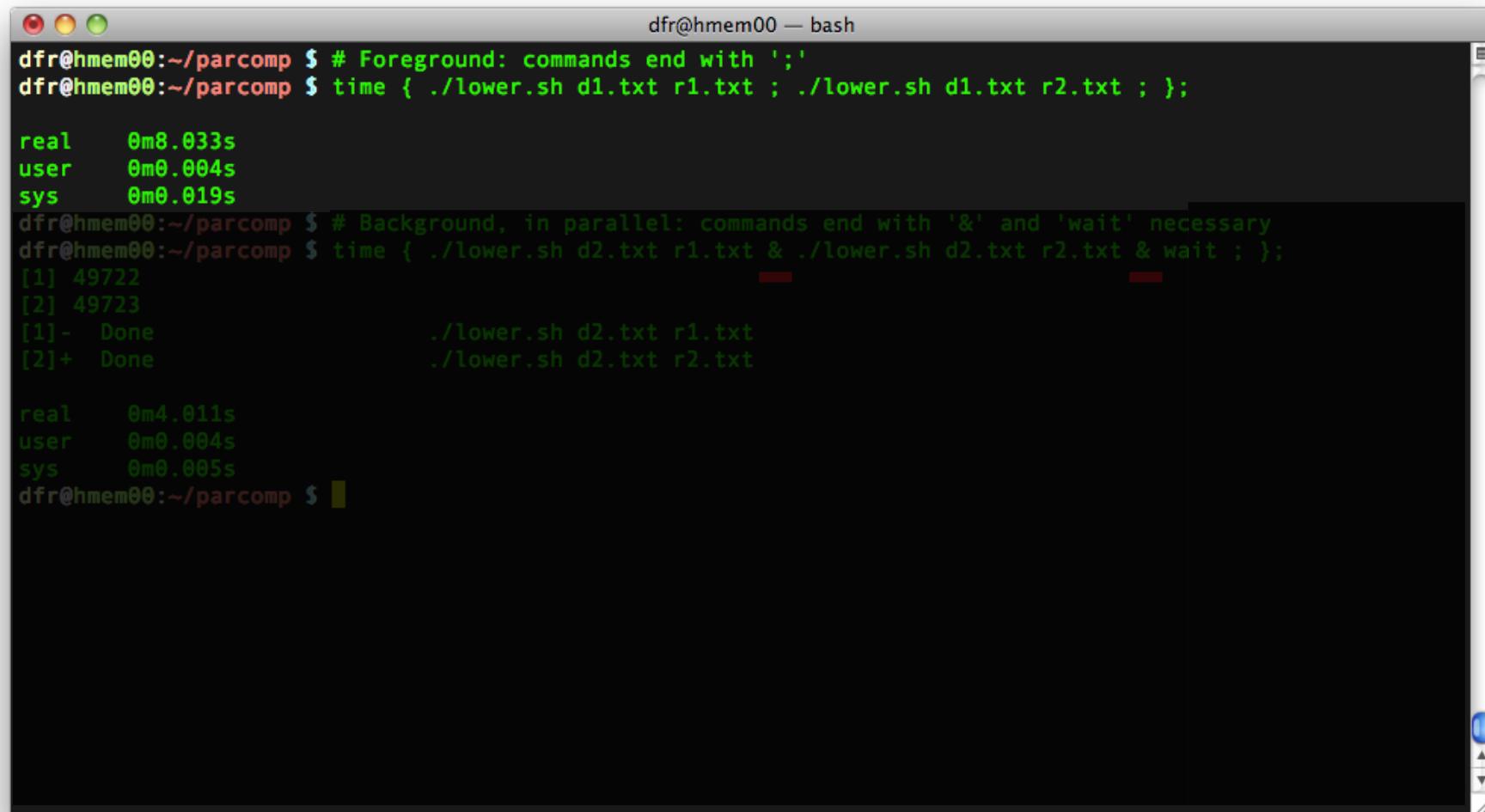
```
dfr@hmem00:~/parcomp $ # Foreground: commands end with ';'
dfr@hmem00:~/parcomp $ time { ./lower.sh d1.txt r1.txt ; ./lower.sh d1.txt r2.txt ; }

real    0m8.033s
user    0m0.004s
sys     0m0.019s
dfr@hmem00:~/parcomp $ # Background, in parallel: commands end with '&' and 'wait' necessary
dfr@hmem00:~/parcomp $ time { ./lower.sh d2.txt r1.txt & ./lower.sh d2.txt r2.txt & wait ; ;
[1] 49722
[2] 49723
[1]-  Done                  ./lower.sh d2.txt r1.txt
[2]+  Done                  ./lower.sh d2.txt r2.txt

real    0m4.011s
user    0m0.004s
sys     0m0.005s
dfr@hmem00:~/parcomp $
```

4.2. Job control & Parallel processes in Bash

Run the program twice and measure the time it takes



The screenshot shows a terminal window with a black background and white text. The window title is "dfr@hmem00 — bash". The terminal displays the following command-line session:

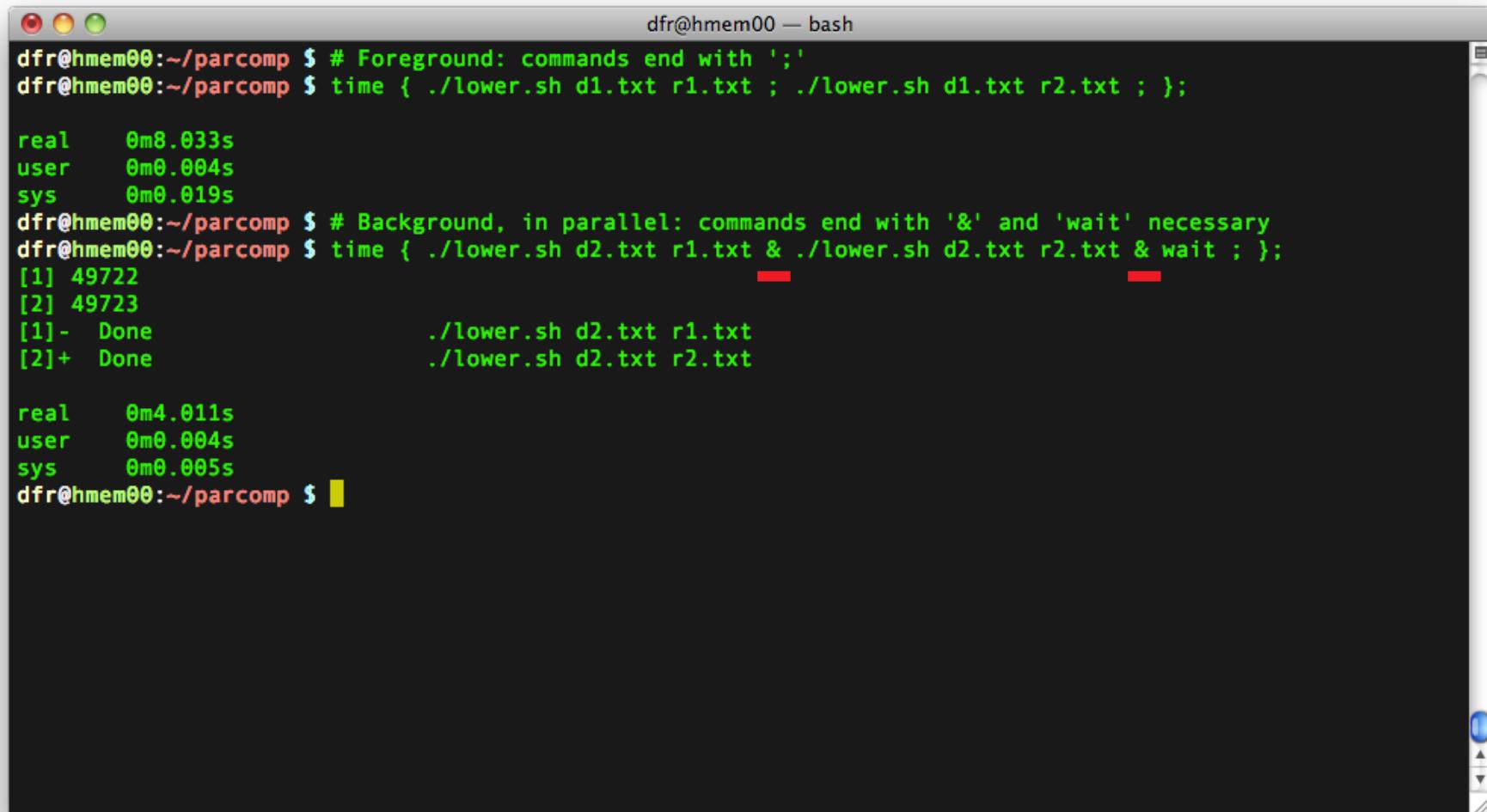
```
dfr@hmem00:~/parcomp $ # Foreground: commands end with ';'
dfr@hmem00:~/parcomp $ time { ./lower.sh d1.txt r1.txt ; ./lower.sh d1.txt r2.txt ; };

real    0m8.033s
user    0m0.004s
sys     0m0.019s
dfr@hmem00:~/parcomp $ # Background, in parallel: commands end with '&' and 'wait' necessary
dfr@hmem00:~/parcomp $ time { ./lower.sh d2.txt r1.txt & ./lower.sh d2.txt r2.txt & wait ; };
[1] 49722
[2] 49723
[1]-  Done                  ./lower.sh d2.txt r1.txt
[2]+  Done                  ./lower.sh d2.txt r2.txt

real    0m4.011s
user    0m0.004s
sys     0m0.005s
dfr@hmem00:~/parcomp $
```

4.2. Job control & Parallel processes in Bash

Run the program twice “in the background” and measure the time



The screenshot shows a terminal window with a black background and white text. The window title is "dfr@hmem00 — bash". The terminal displays the following command-line session:

```
dfr@hmem00:~/parcomp $ # Foreground: commands end with ';'
dfr@hmem00:~/parcomp $ time { ./lower.sh d1.txt r1.txt ; ./lower.sh d1.txt r2.txt ; };

real    0m8.033s
user    0m0.004s
sys     0m0.019s
dfr@hmem00:~/parcomp $ # Background, in parallel: commands end with '&' and 'wait' necessary
dfr@hmem00:~/parcomp $ time { ./lower.sh d2.txt r1.txt & ./lower.sh d2.txt r2.txt & wait ; };
[1] 49722
[2] 49723
[1]-  Done                  ./lower.sh d2.txt r1.txt
[2]+  Done                  ./lower.sh d2.txt r2.txt

real    0m4.011s
user    0m0.004s
sys     0m0.005s
dfr@hmem00:~/parcomp $
```

4.2. Job control & Parallel processes in Bash

Parallel for loop in Bash:

```
for i in {1..10}; do
  command1
  command2
done
```

```
for i in {1..10}; do
(
  command1
  command2
) &
done; wait
```

(...) & : creates a sub-shell with all commands in the bloc and start it in the background
wait : barrier to synchronize all sub-shells

4.3.1. One program and many files

The xargs command distributes data from stdin to program

```
dfr@hmem00:~/parcomp $ ls d?.txt
d1.txt  d2.txt  d3.txt  d4.txt
dfr@hmem00:~/parcomp $ ls d?.txt | xargs -n 1 echo "File: "
File:  d1.txt
File:  d2.txt
File:  d3.txt
File:  d4.txt
dfr@hmem00:~/parcomp $ time { ls d?.txt | xargs -n 1 ./lower.sh > /dev/null ; }

real    0m16.041s
user    0m0.010s
sys     0m0.006s
dfr@hmem00:~/parcomp $ time { ls d?.txt | xargs -n 1 -P 4 ./lower.sh > /dev/null ; }

real    0m4.014s
user    0m0.008s
sys     0m0.016s
dfr@hmem00:~/parcomp $
```

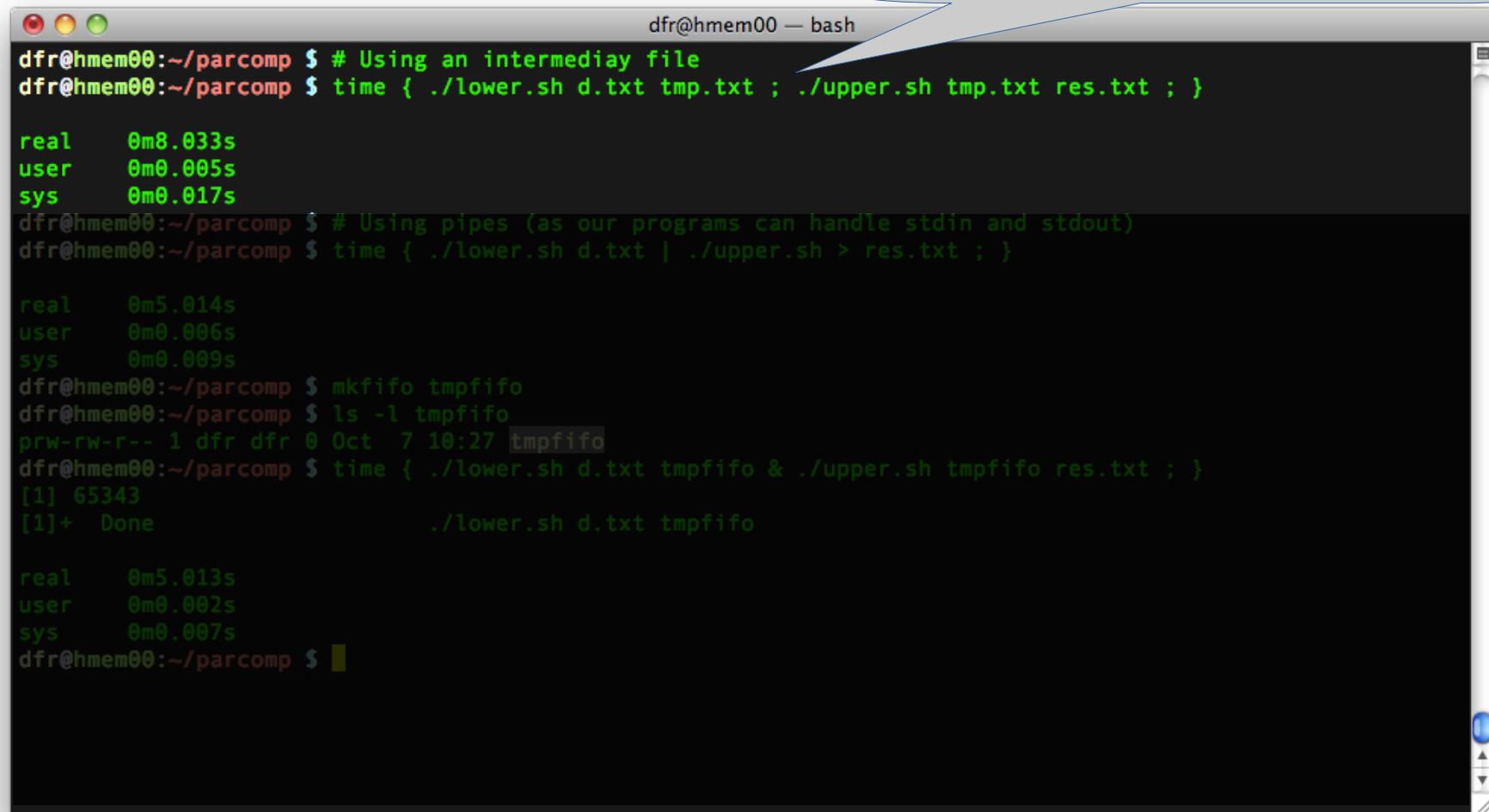
Equivalent to
.lower.sh d1.txt ;
.lower.sh d2.txt ;
.lower.sh d3.txt ;
.lower.sh d3.txt ;

Equivalent to
.lower.sh d1.txt &
.lower.sh d2.txt &
.lower.sh d3.txt &
.lower.sh d3.txt &
wait

4.3.2. Several programs and one file

Using UNIX pipes for pipelining operations

./upper.sh waits for ./lower.sh to finish
Note the intermediate file



```
dfr@hmem00:~/parcomp $ # Using an intermediary file
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt tmp.txt ; ./upper.sh tmp.txt res.txt ; }

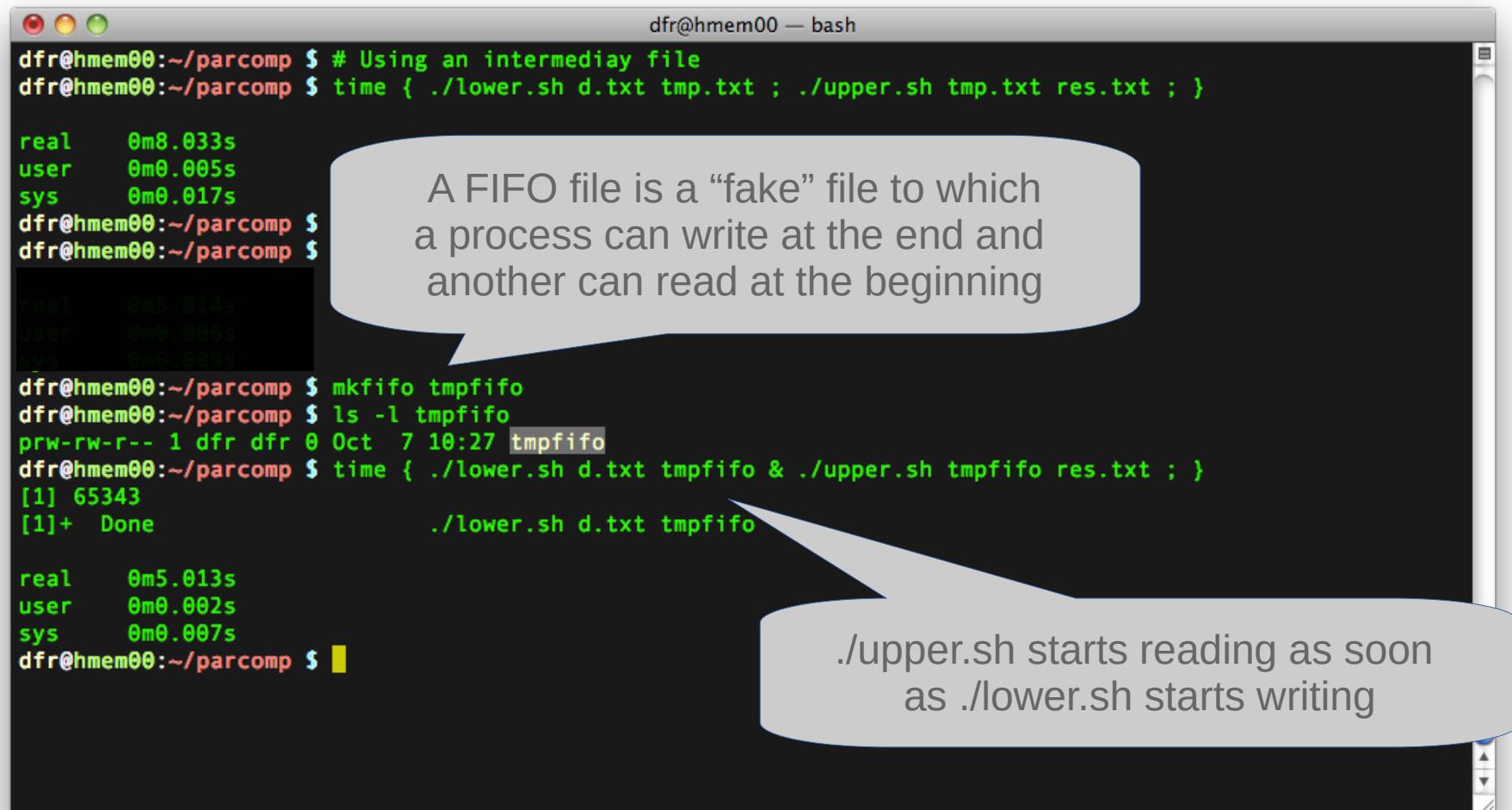
real    0m8.033s
user    0m0.005s
sys     0m0.017s
dfr@hmem00:~/parcomp $ # Using pipes (as our programs can handle stdin and stdout)
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt | ./upper.sh > res.txt ; }

real    0m5.014s
user    0m0.006s
sys     0m0.009s
dfr@hmem00:~/parcomp $ mkfifo tmpfifo
dfr@hmem00:~/parcomp $ ls -l tmpfifo
prw-rw-r-- 1 dfr dfr 0 Oct  7 10:27 tmpfifo
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt tmpfifo & ./upper.sh tmpfifo res.txt ; }
[1] 65343
[1]+  Done                      ./lower.sh d.txt tmpfifo

real    0m5.013s
user    0m0.002s
sys     0m0.007s
dfr@hmem00:~/parcomp $
```

4.3.2. Several programs and one file

Using UNIX fifos for pipelining operations



```
dfr@hmem00:~/parcomp $ # Using an intermediary file
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt tmp.txt ; ./upper.sh tmp.txt res.txt ; }

real    0m8.033s
user    0m0.005s
sys     0m0.017s
dfr@hmem00:~/parcomp $
dfr@hmem00:~/parcomp $

real    0m5.014s
user    0m0.006s
sys     0m0.009s
dfr@hmem00:~/parcomp $ mkfifo tmpfifo
dfr@hmem00:~/parcomp $ ls -l tmpfifo
prw-rw-r-- 1 dfr dfr 0 Oct  7 10:27 tmpfifo
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt tmpfifo & ./upper.sh tmpfifo res.txt ; }
[1] 65343
[1]+  Done                      ./lower.sh d.txt tmpfifo

real    0m5.013s
user    0m0.002s
sys     0m0.007s
dfr@hmem00:~/parcomp $
```

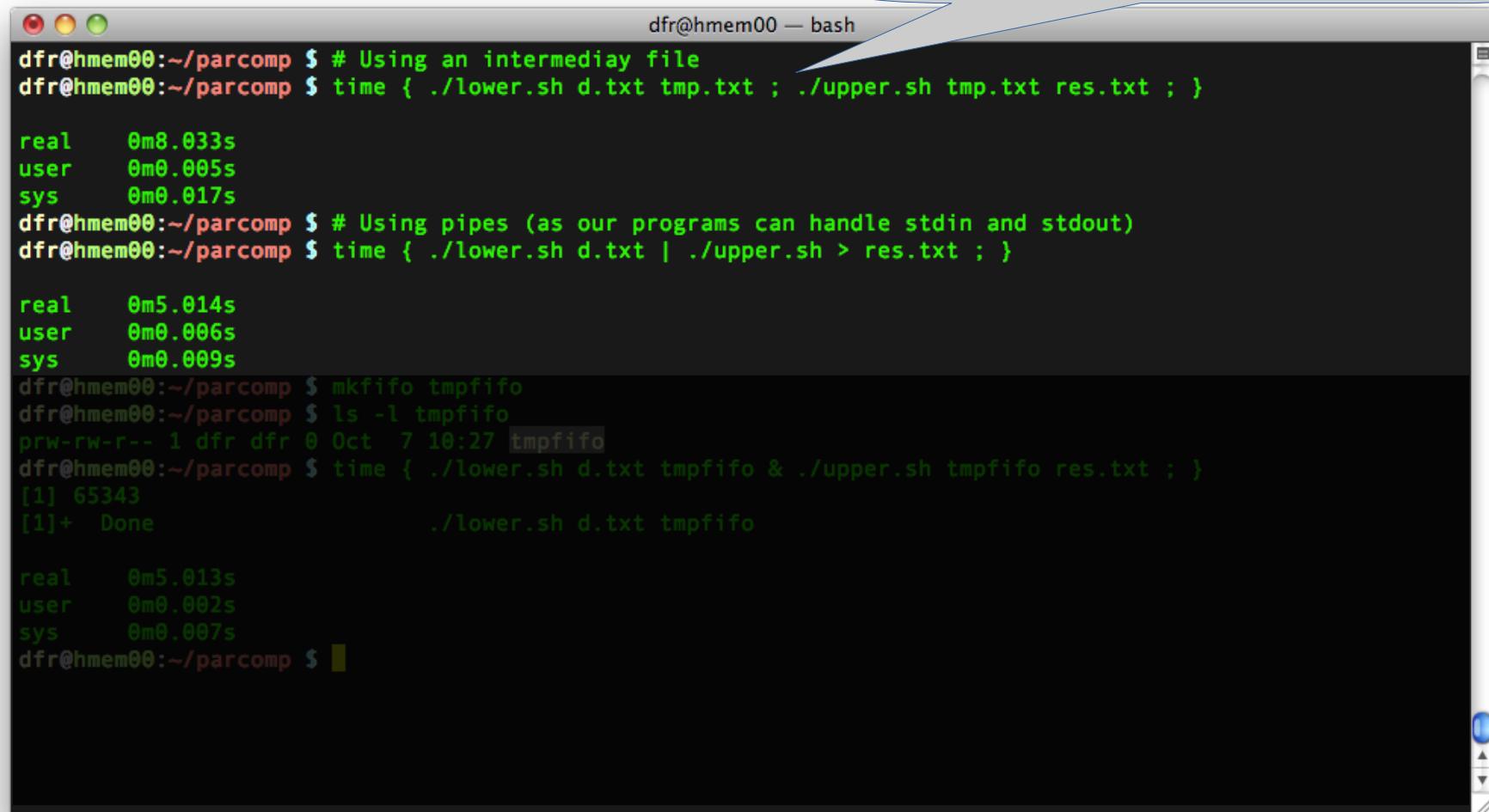
A FIFO file is a “fake” file to which a process can write at the end and another can read at the beginning

./upper.sh starts reading as soon as ./lower.sh starts writing

4.3.2. Several programs and one file

Using UNIX pipes for pipelining operations

./upper.sh waits for ./lower.sh to finish
Note the intermediate file



```
dfr@hmem00:~/parcomp $ # Using an intermediary file
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt tmp.txt ; ./upper.sh tmp.txt res.txt ; }

real    0m8.033s
user    0m0.005s
sys     0m0.017s
dfr@hmem00:~/parcomp $ # Using pipes (as our programs can handle stdin and stdout)
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt | ./upper.sh > res.txt ; }

real    0m5.014s
user    0m0.006s
sys     0m0.009s
dfr@hmem00:~/parcomp $ mkfifo tmpfifo
dfr@hmem00:~/parcomp $ ls -l tmpfifo
prw-rw-r-- 1 dfr dfr 0 Oct  7 10:27 tmpfifo
dfr@hmem00:~/parcomp $ time { ./lower.sh d.txt tmpfifo & ./upper.sh tmpfifo res.txt ; }
[1] 65343
[1]+  Done                      ./lower.sh d.txt tmpfifo

real    0m5.013s
user    0m0.002s
sys     0m0.007s
dfr@hmem00:~/parcomp $
```

4.3.3. One program and one large file

The split command distributes data from stdin to program

```
dfr@hmem00:~/parcomp $ # One process to process the whole file
dfr@hmem00:~/parcomp $ time { cat d.txt | ./lower.sh > res.txt

real    0m4.014s
user    0m0.003s
sys     0m0.009s
dfr@hmem00:~/parcomp $ # Four processes handling one line in round robin fashion
dfr@hmem00:~/parcomp $ time { cat d.txt | split --unbuffered --number r/4 --filter ./lower.sh >res.txt ; }

real    0m1.011s
user    0m0.009s
sys     0m0.021s
dfr@hmem00:~/parcomp $ !! & top -u dfr -bn1 | grep lower
time { cat d.txt | split --unbuffered --number r/4 --filter ./lower.sh >res.txt ; } & top -u dfr -bn1 | grep lower
[1] 12817
12822 dfr      20  0 103m 1252 1052 S  0.0  0.0  0:00.00 lower.sh
12823 dfr      20  0 103m 1252 1052 S  0.0  0.0  0:00.00 lower.sh
12824 dfr      20  0 103m 1252 1052 S  0.0  0.0  0:00.00 lower.sh
12825 dfr      20  0 103m 1252 1052 S  0.0  0.0  0:00.00 lower.sh
dfr@hmem00:~/parcomp $
real    0m1.011s
user    0m0.011s
sys     0m0.019s

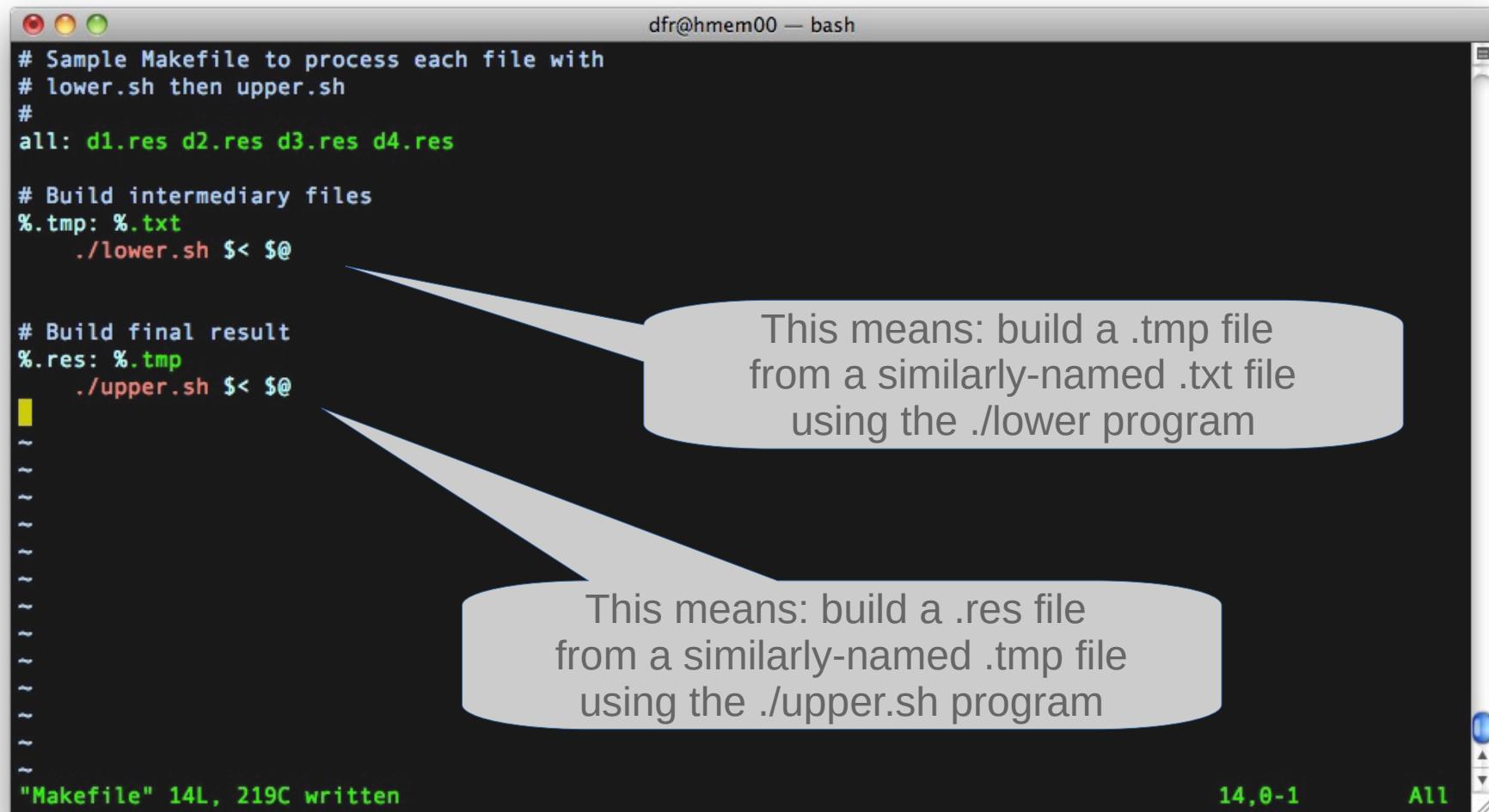
[1]+  Done                      time { cat d.txt | split --unbuffered --number r/4 --filter ./lower.sh > res
.txt; }
dfr@hmem00:~/parcomp $
```

Split the file and start 4 processes

Need recent version of Coreutils/8.22-goolf-1.4.10

4.3.4. Several programs and many files

A Makefile describes dependencies and is executed with 'make'



```
# Sample Makefile to process each file with
# lower.sh then upper.sh
#
all: d1.res d2.res d3.res d4.res

# Build intermediary files
%.tmp: %.txt
    ./lower.sh $< $@

# Build final result
%.res: %.tmp
    ./upper.sh $< $@
```

This means: build a .tmp file from a similarly-named .txt file using the ./lower program

This means: build a .res file from a similarly-named .tmp file using the ./upper.sh program

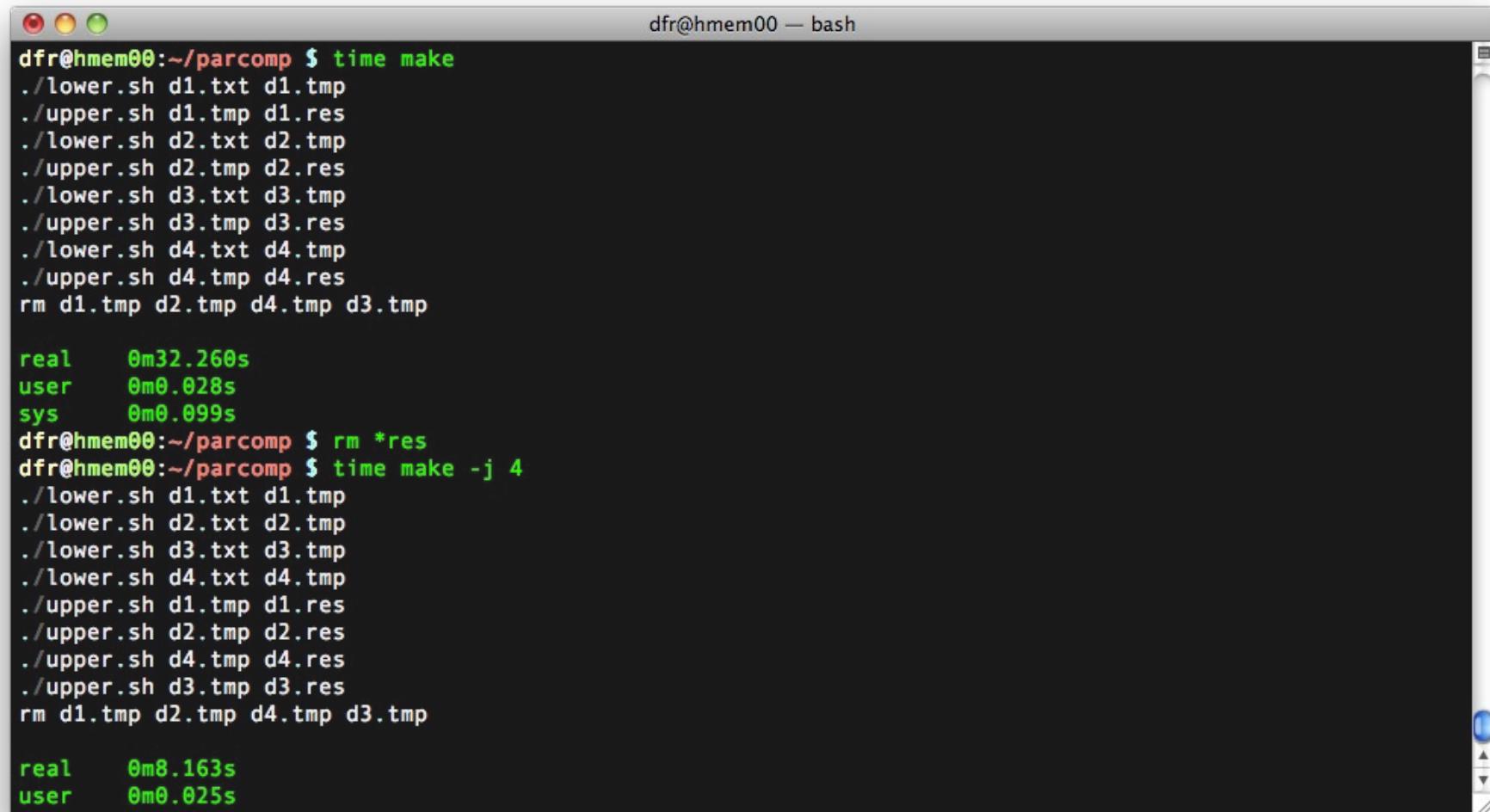
"Makefile" 14L, 219C written

14,0-1

All

4.3.4. Several programs and many files

The ‘make’ command can operate in parallel



The screenshot shows a terminal window with a dark background and light-colored text. The window title is "dfr@hmem00 — bash". The terminal displays two sessions of command execution:

```
dfr@hmem00:~/parcomp $ time make
./lower.sh d1.txt d1.tmp
./upper.sh d1.tmp d1.res
./lower.sh d2.txt d2.tmp
./upper.sh d2.tmp d2.res
./lower.sh d3.txt d3.tmp
./upper.sh d3.tmp d3.res
./lower.sh d4.txt d4.tmp
./upper.sh d4.tmp d4.res
rm d1.tmp d2.tmp d4.tmp d3.tmp

real    0m32.260s
user    0m0.028s
sys     0m0.099s

dfr@hmem00:~/parcomp $ rm *res
dfr@hmem00:~/parcomp $ time make -j 4
./lower.sh d1.txt d1.tmp
./lower.sh d2.txt d2.tmp
./lower.sh d3.txt d3.tmp
./lower.sh d4.txt d4.tmp
./upper.sh d1.tmp d1.res
./upper.sh d2.tmp d2.res
./upper.sh d4.tmp d4.res
./upper.sh d3.tmp d3.res
rm d1.tmp d2.tmp d4.tmp d3.tmp

real    0m8.163s
user    0m0.025s
```

Summary

- You have either
 - one very large file to process
 - with one program: `split`
 - with several programs: `fifo (or pipes)`
 - many files to process
 - with one program `xargs`
 - with many programs `make`

4.4. GNU Parallel

GNU Parallel

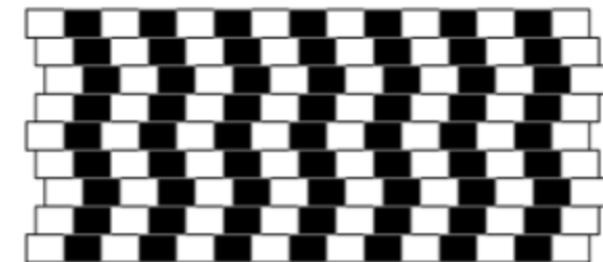
GNU **parallel** is a shell tool for executing jobs in parallel using one or more computers. A job can be a single command or a small script that has to be run for each of the lines in the input. The typical input is a list of files, a list of hosts, a list of users, a list of URLs, or a list of tables. A job can also be a command that reads from a pipe. GNU **parallel** can then split the input and pipe it into commands in parallel.

If you use **xargs** and **tee** today you will find GNU **parallel** very easy to use as GNU **parallel** is written to have the same options as **xargs**. If you write loops in shell, you will find GNU **parallel** may be able to replace most of the loops and make them run faster by running several jobs in parallel.

GNU **parallel** makes sure output from the commands is the same output as you would get had you run the commands sequentially. This makes it possible to use output from GNU **parallel** as input for other programs.

For each line of input GNU **parallel** will execute *command* with the line as arguments. If no *command* is given, the line of input is executed. Several lines will be run in parallel. GNU **parallel** can often be used as a substitute for **xargs** or **cat | bash**.

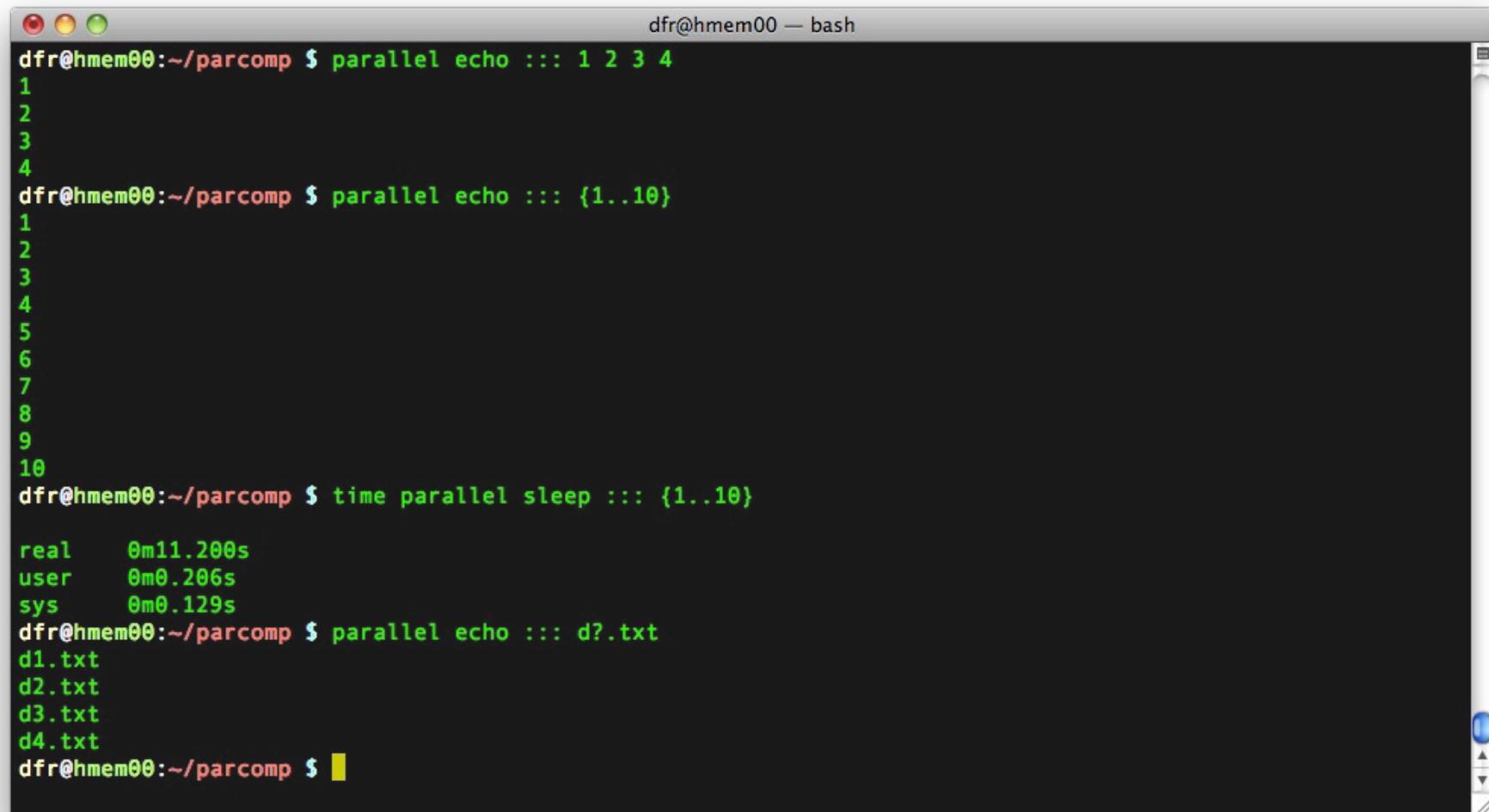
More complicated to use but very powerful
Might not be available everywhere but you can install it



For people who live life in the parallel lane.

4.4. GNU Parallel

- Syntax: `parallel command ::: argument list`



```
dfr@hmem00:~/parcomp $ parallel echo :::: 1 2 3 4
1
2
3
4
dfr@hmem00:~/parcomp $ parallel echo :::: {1..10}
1
2
3
4
5
6
7
8
9
10
dfr@hmem00:~/parcomp $ time parallel sleep :::: {1..10}

real    0m11.200s
user    0m0.206s
sys     0m0.129s
dfr@hmem00:~/parcomp $ parallel echo :::: d?.txt
d1.txt
d2.txt
d3.txt
d4.txt
dfr@hmem00:~/parcomp $
```

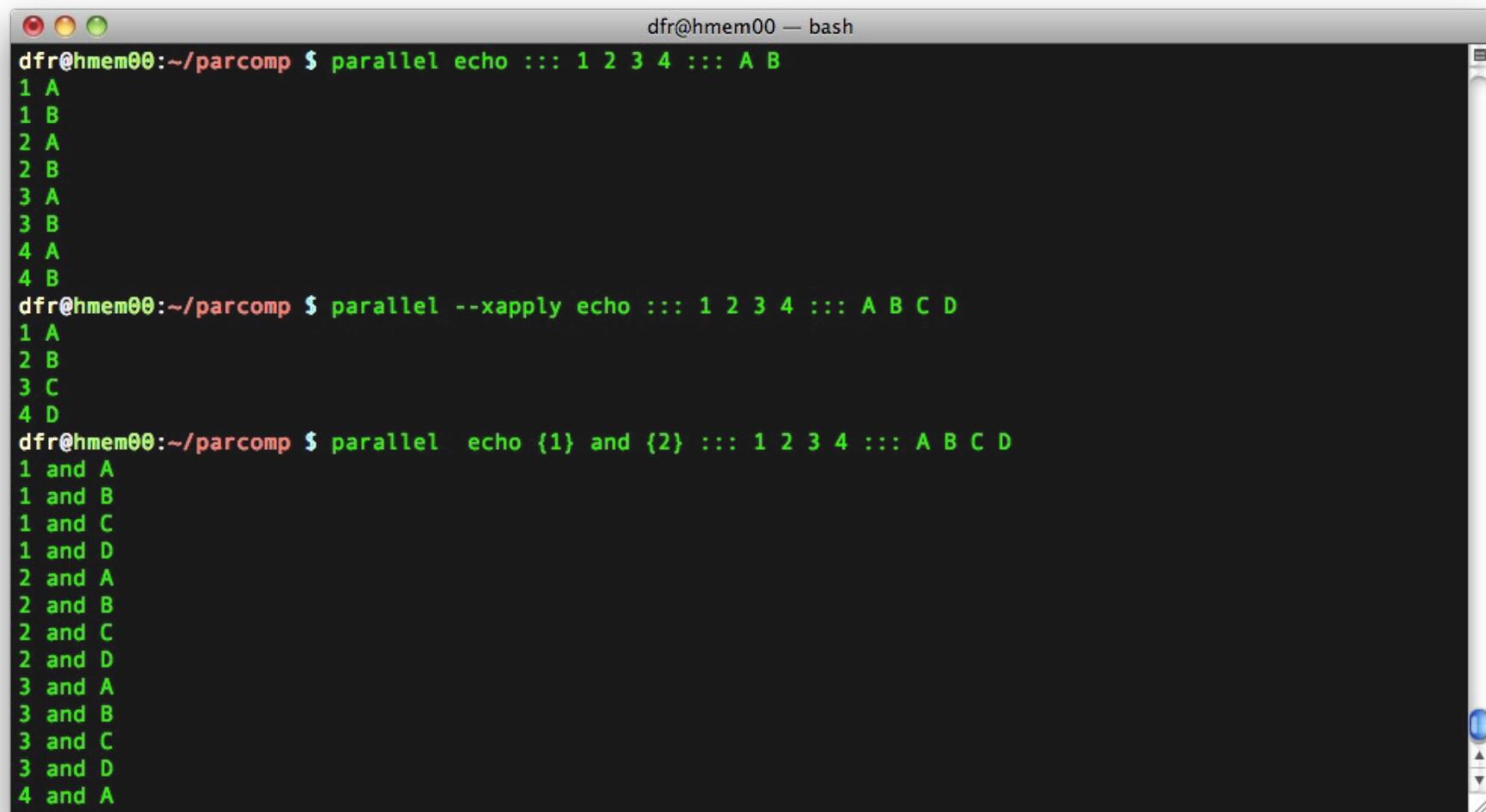
4.4. GNU Parallel

- Syntax: {} as argument placeholder.

```
d1.txt
d2.txt
d3.txt
d4.txt
dfr@hmem00:~/parcomp $ parallel echo {} {.}.res :::: d?.txt
d1.txt d1.res
d2.txt d2.res
d3.txt d3.res
d4.txt d4.res
dfr@hmem00:~/parcomp $ parallel echo {} :::: ../parcomp/d?.txt
../parcomp/d1.txt
../parcomp/d2.txt
../parcomp/d3.txt
../parcomp/d4.txt
dfr@hmem00:~/parcomp $ parallel echo {} :::: ../parcomp/d?.txt
d1.txt
d2.txt
d3.txt
d4.txt
dfr@hmem00:~/parcomp $
dfr@hmem00:~/parcomp $
dfr@hmem00:~/parcomp $
```

4.4. GNU Parallel

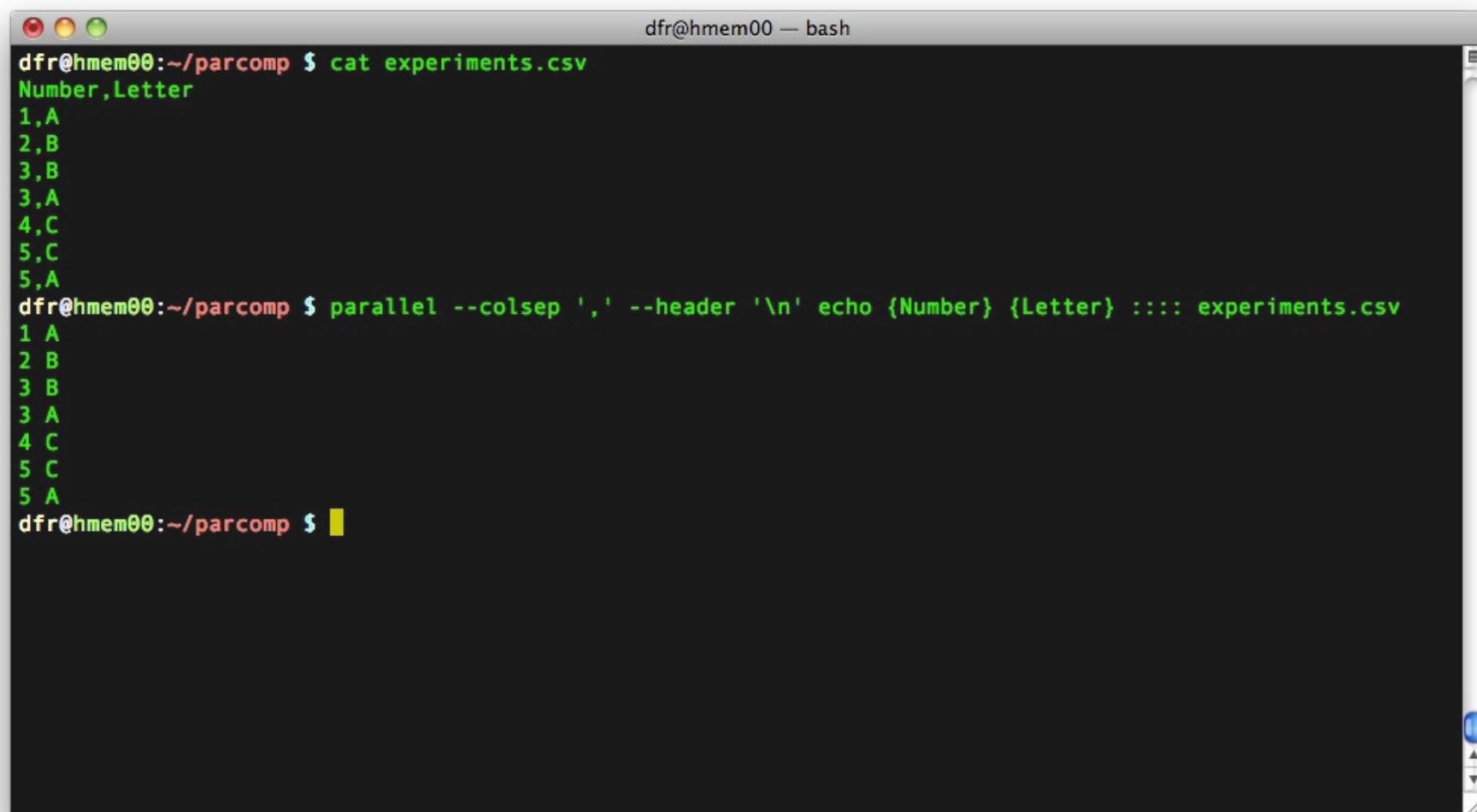
- Multiple parameters and --xapply



```
dfr@hmem00:~/parcomp $ parallel echo :::: 1 2 3 4 :::: A B
1 A
1 B
2 A
2 B
3 A
3 B
4 A
4 B
dfr@hmem00:~/parcomp $ parallel --xapply echo :::: 1 2 3 4 :::: A B C D
1 A
2 B
3 C
4 D
dfr@hmem00:~/parcomp $ parallel echo {1} and {2} :::: 1 2 3 4 :::: A B C D
1 and A
1 and B
1 and C
1 and D
2 and A
2 and B
2 and C
2 and D
3 and A
3 and B
3 and C
3 and D
4 and A
```

4.4. GNU Parallel

- When arguments are in a file : use :::: (4x ':')



The screenshot shows a terminal window with a dark background and light-colored text. The window title is "dfr@hmem00 — bash". The terminal displays the following command and its output:

```
dfr@hmem00:~/parcomp $ cat experiments.csv
Number,Letter
1,A
2,B
3,B
3,A
4,C
5,C
5,A
dfr@hmem00:~/parcomp $ parallel --colsep ',' --header '\n' echo {Number} {Letter} :::: experiments.csv
1 A
2 B
3 B
3 A
4 C
5 C
5 A
dfr@hmem00:~/parcomp $
```

4.4. GNU Parallel

Interactivity

GNU parallel can ask the user if a command should be run using **--interactive**:

```
parallel --interactive echo :::: 1 2 3
```

Output:

```
echo 1 ?...y
echo 2 ?...n
1
echo 3 ?...y
3
```

4.4. GNU Parallel

Timing

Some jobs do heavy I/O when they start. To avoid a thundering herd **GNU parallel** can delay starting new jobs. **--delay X** will make sure there is at least X seconds between each start:

```
parallel --delay 2.5 echo Starting {}\\;date :::: 1 2 3
```

Output:

```
Starting 1
Thu Aug 15 16:24:33 CEST 2013
Starting 2
Thu Aug 15 16:24:35 CEST 2013
Starting 3
Thu Aug 15 16:24:38 CEST 2013
```

4.4. GNU Parallel

Progress information

Based on the runtime of completed jobs **GNU parallel** can estimate the total runtime:

```
parallel --eta sleep ::: 1 3 2 2 1 3 3 2 1
```

Output:

```
Computers / CPU cores / Max jobs to run
1:local / 2 / 2

Computer:jobs running/jobs completed/%of started jobs/
  Average seconds to complete
ETA: 2s 0left 1.11avg  local:0/9/100%/1.1s
```

4.4. GNU Parallel

With a joblog **GNU parallel** can be stopped and later pickup where it left off. It is important that the input of the completed jobs is unchanged.

```
parallel --joblog /tmp/log exit ::: 1 2 3 0
cat /tmp/log
parallel --resume --joblog /tmp/log exit ::: 1 2 3 0 0 0
cat /tmp/log
```

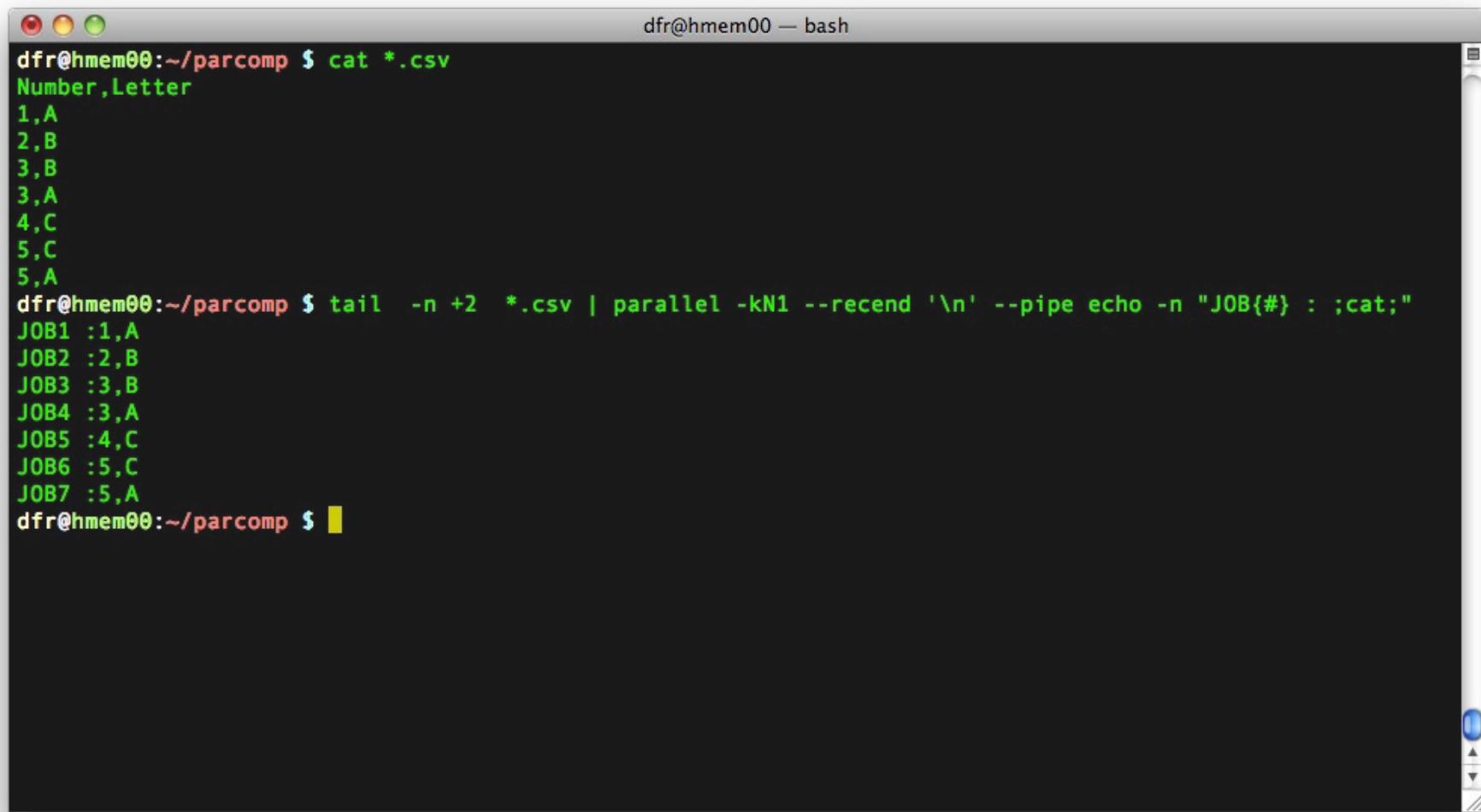
Output:

Seq	Host	Starttime	Runtime	Send	Receive	Exitval	Signal	Command
1	:	1376580069.544	0.008	0	0	1	0	exit 1
2	:	1376580069.552	0.009	0	0	2	0	exit 2
3	:	1376580069.560	0.012	0	0	3	0	exit 3
4	:	1376580069.571	0.005	0	0	0	0	exit 0

Seq	Host	Starttime	Runtime	Send	Receive	Exitval	Signal	Command
1	:	1376580069.544	0.008	0	0	1	0	exit 1
2	:	1376580069.552	0.009	0	0	2	0	exit 2
3	:	1376580069.560	0.012	0	0	3	0	exit 3
4	:	1376580069.571	0.005	0	0	0	0	exit 0
5	:	1376580070.028	0.009	0	0	0	0	exit 0
6	:	1376580070.038	0.007	0	0	0	0	exit 0

4.4. GNU Parallel

- Split a file with --pipe



The screenshot shows a terminal window with a dark background and light-colored text. The window title is "dfr@hmem00 — bash". The terminal content is as follows:

```
dfr@hmem00:~/parcomp $ cat *.csv
Number,Letter
1,A
2,B
3,B
3,A
4,C
5,C
5,A
dfr@hmem00:~/parcomp $ tail -n +2 *.csv | parallel -kN1 --recend '\n' --pipe echo -n "JOB{#} : ;cat;"
JOB1 :1,A
JOB2 :2,B
JOB3 :3,B
JOB4 :3,A
JOB5 :4,C
JOB6 :5,C
JOB7 :5,A
dfr@hmem00:~/parcomp $
```

Summary

1. General concepts, definitions, challenges
2. Hardware for parallel computing
3. Programming models
4. User tools