



Consortium des Equipements
de Calcul Intensif
en Fédération Wallonie-Bruxelles

Introduction to Scientific Data Management

damien.francois@uclouvain.be
November 2020

<http://www.cec-hpc.be/training.html>



Goal of this session:



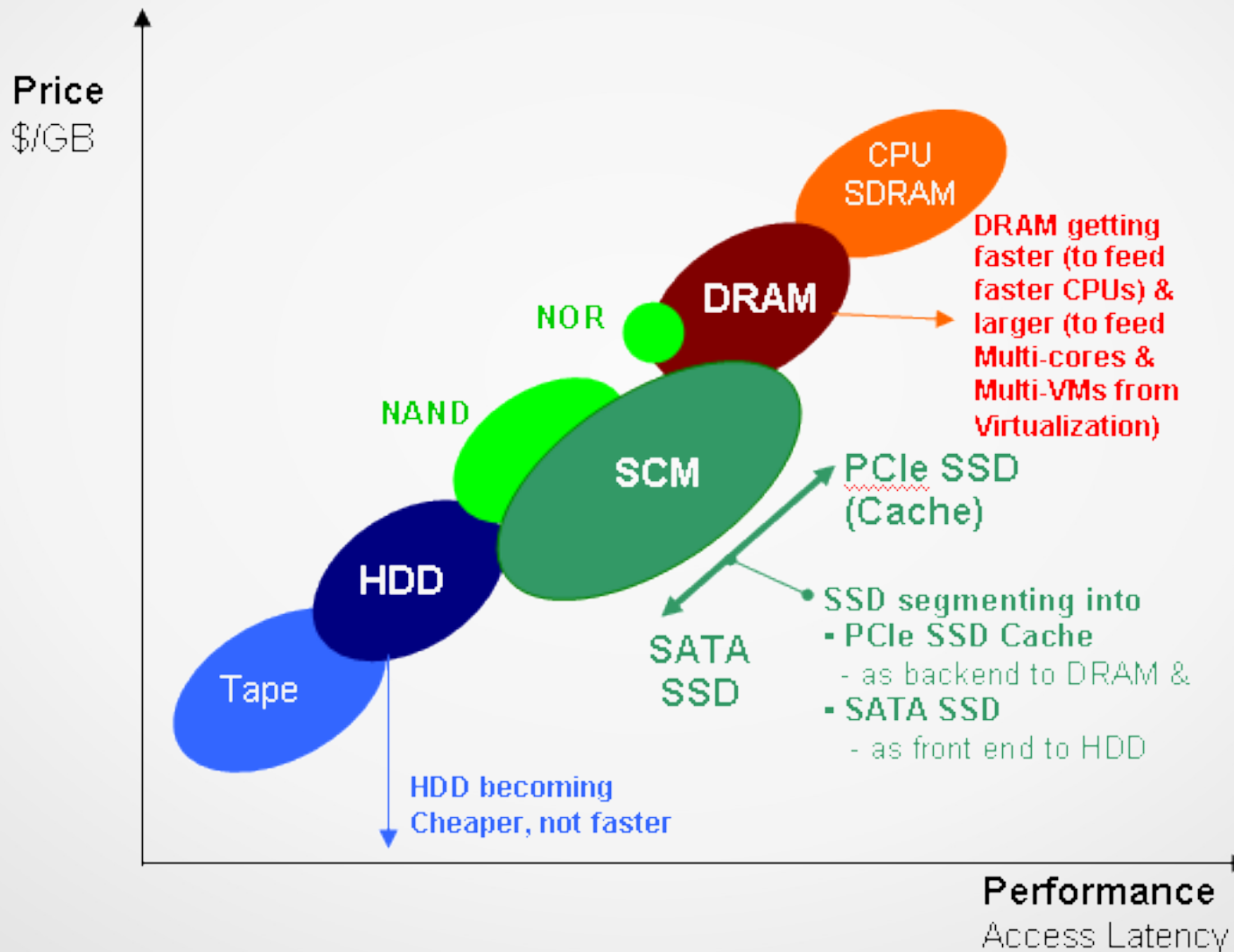
“Share tools, tips and tricks related to the **storage, transfer, and sharing** of scientific data”

1. Data storage



Filesystems – Object stores – Databases

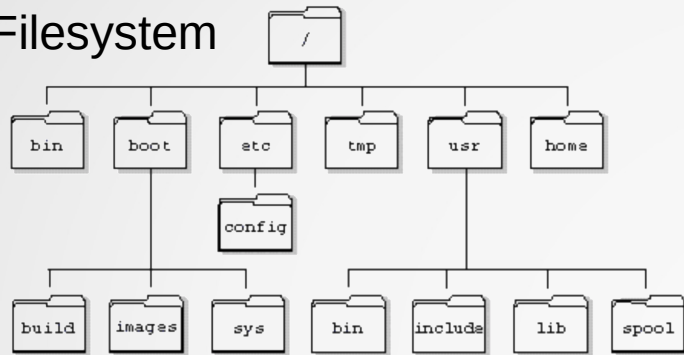
Storage Technologies



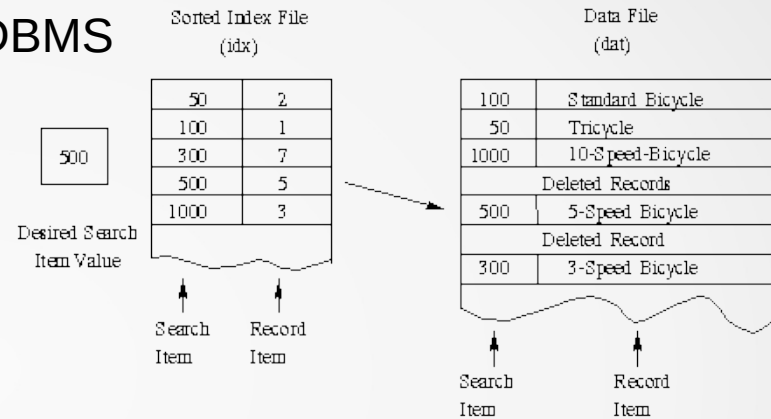
Storage paradigms



Filesystem



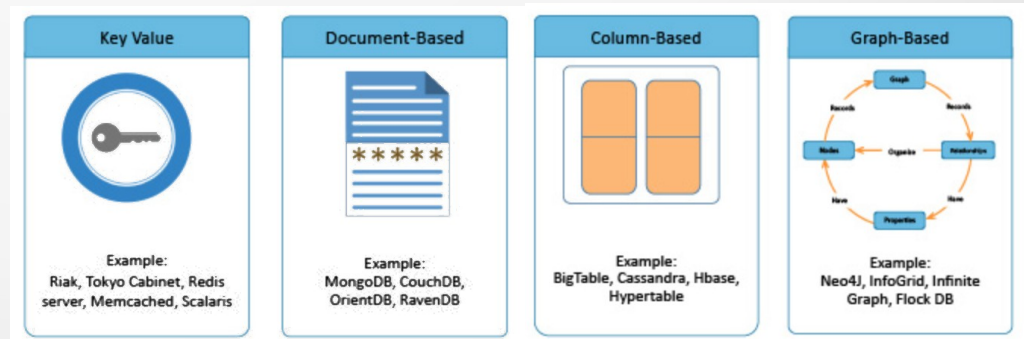
RDBMS



Objects store



NoSQL



1.1 Filesystems

(local) Filesystems



Generation 0: No system at all. There was just an arbitrary stream of data. Think punchcards, data on audiocassette, Atari 2600 ROM carts.

Generation 1: Early random access. Here, there are multiple named files on one device with no folders or other metadata. Think Apple][DOS (but not ProDOS!) as one example.

Generation 2: Early organization (aka folders). When devices became capable of holding hundreds of files, better organization became necessary. We're referring to TRS-DOS, Apple //c ProDOS, MS-DOS FAT/FAT32, etc.

Generation 3: Metadata—ownership, permissions, etc. As the user count on machines grew higher, the ability to restrict and control access became necessary. This includes AT&T UNIX, Netware, early NTFS, etc.

Generation 4: Journaling! This is the killer feature defining all current, modern filesystems—ext4, modern NTFS, UFS2, XFS, you name it. Journaling keeps the filesystem from becoming inconsistent in the event of a crash, making it much less likely that you'll lose data, or even an entire disk, when the power goes off or the kernel crashes.

Generation 5: Copy on Write snapshots, Per-block checksumming, Volume management, Far-future scalability, Asynchronous incremental replication, Online compression. Generation 5 filesystems are Btrfs and ZFS.

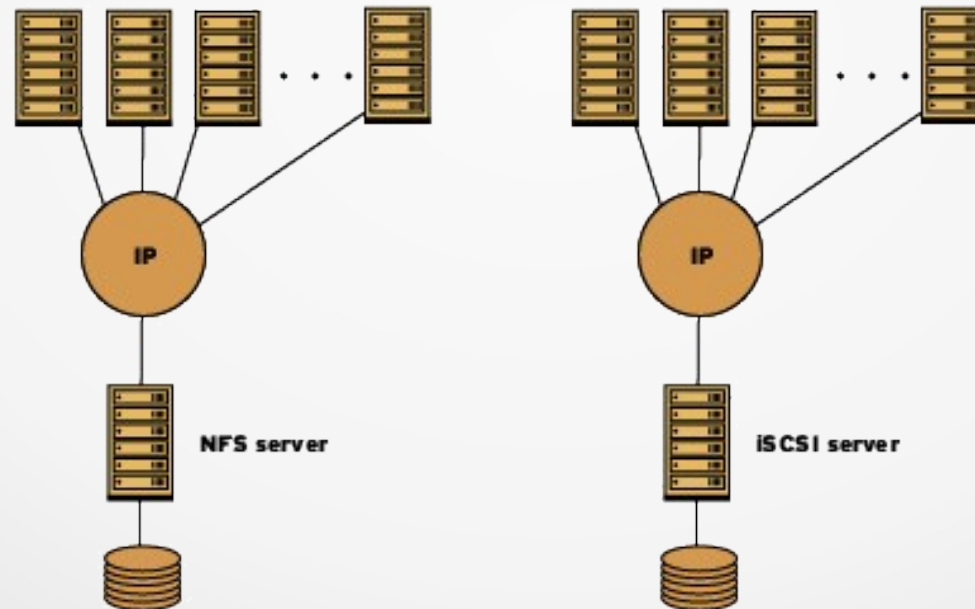


Typical usage: Operating system & Local scratch space

Network filesystem

One source many consumers

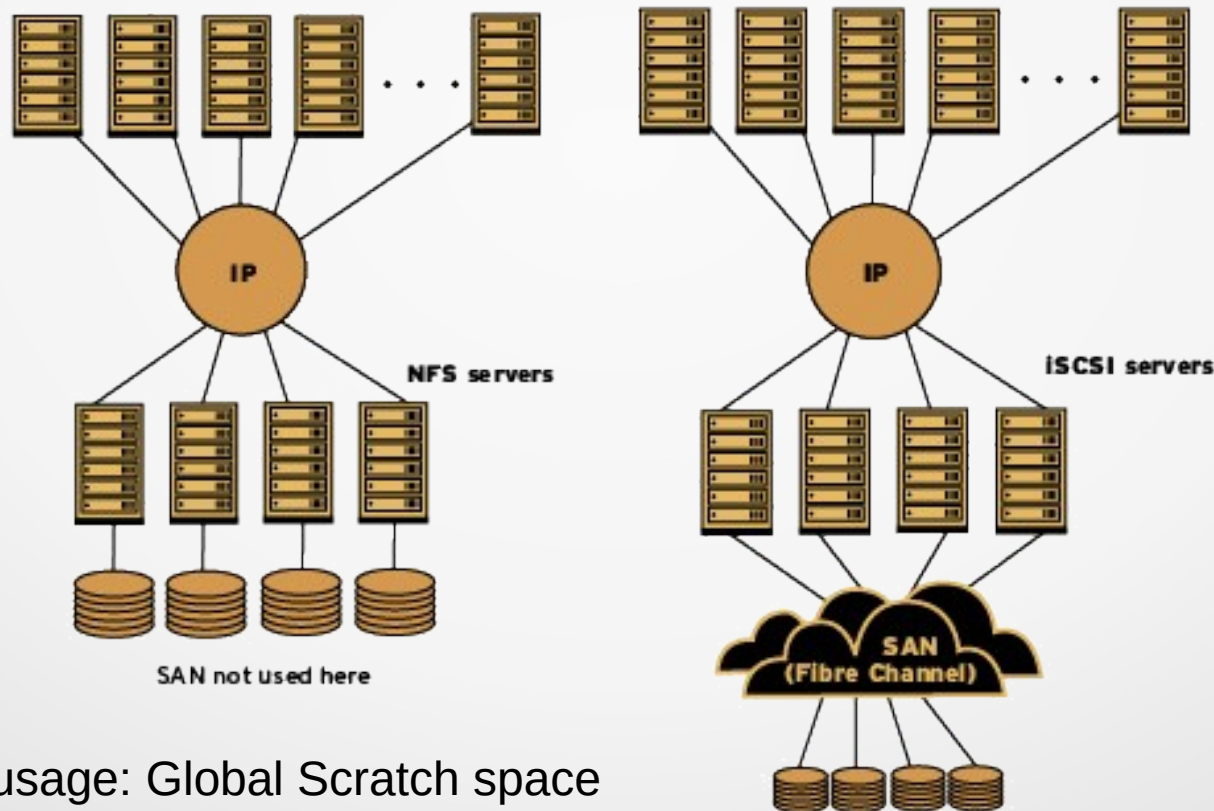
NAS: ex. NFS SAN: ex. GFS2



Typical usage: Home directories, Mass storage

Parallel / distributed filesystem

Many sources many consumers
ex: Lustre, GPFS, BeeGFS, GlusterFS



Typical usage: Global Scratch space

Special filesystems – in memory

A screenshot of a web browser window displaying the kernel.org documentation for the tmpfs filesystem. The browser's address bar shows the URL 'https://www.kernel.org/doc/Documentation/filesystems/tmpfs.txt'. The page content is rendered in a monospaced font and describes the characteristics and usage of tmpfs.

```
https://www.kernel.org/doc/Documentation/filesystems/tmpfs.txt
https://www.kernel.org/doc/Documentat
man tmpfs
CÉCI Login CÉCI Redmine LimeSurvey LimeSurvey2000 wiki-users
https://www.kernel.org/doc/Doc...
+

Tmpfs is a file system which keeps all files in virtual memory.

Everything in tmpfs is temporary in the sense that no files will be
created on your hard drive. If you unmount a tmpfs instance,
everything stored therein is lost.

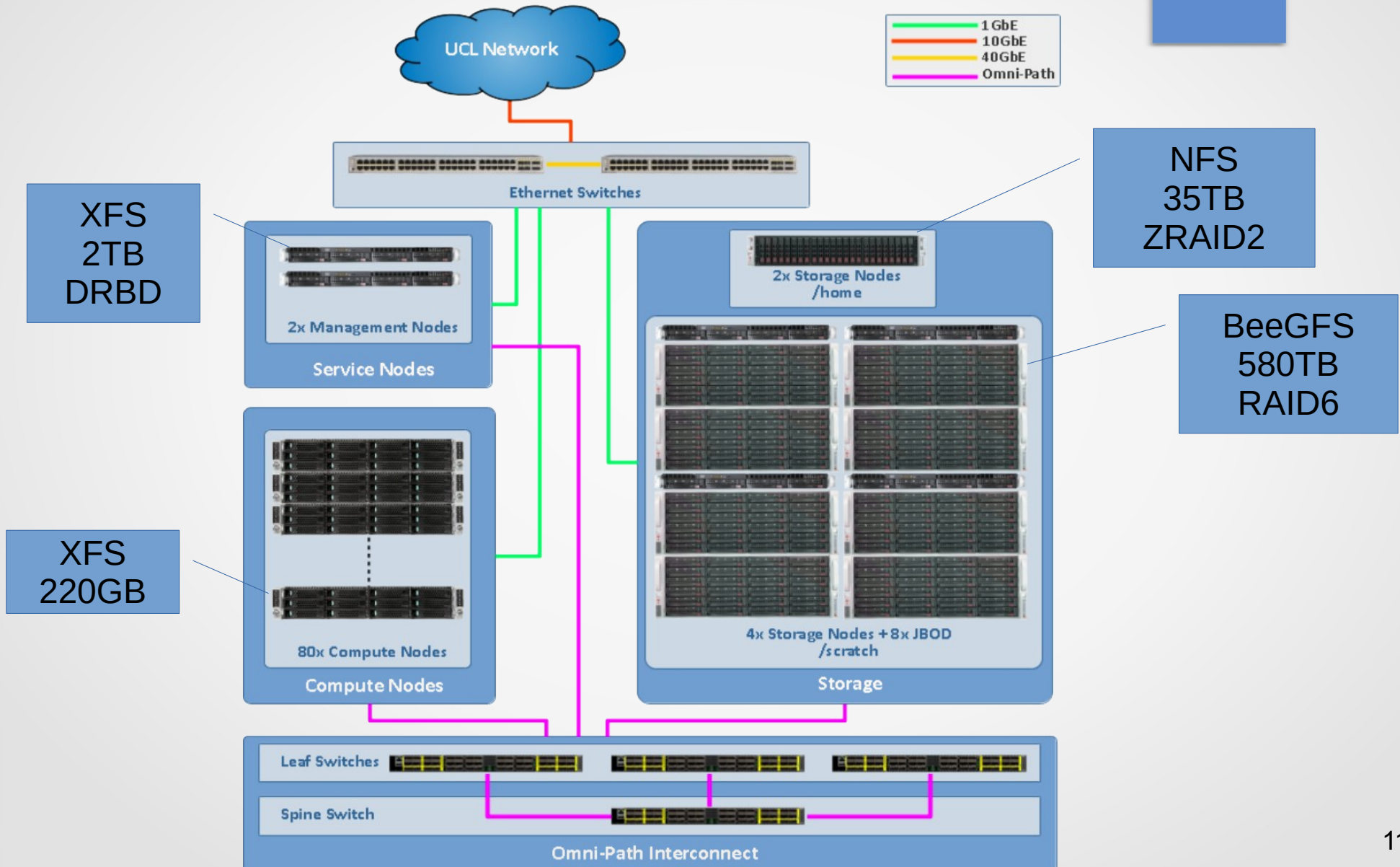
tmpfs puts everything into the kernel internal caches and grows and
shrinks to accommodate the files it contains and is able to swap
unneeded pages out to swap space. It has maximum size limits which can
be adjusted on the fly via 'mount -o remount ...'

If you compare it to ramfs (which was the template to create tmpfs)
you gain swapping and limit checking. Another similar thing is the RAM
disk (/dev/ram*), which simulates a fixed size hard disk in physical
RAM, where you have to create an ordinary filesystem on top. Ramdisks
cannot swap and you do not have the possibility to resize them.

Since tmpfs lives completely in the page cache and on swap, all tmpfs
pages currently in memory will show up as cached. It will not show up
as shared or something like that. Further on you can check the actual
RAM+swap use of a tmpfs instance with df(1) and du(1).
```

Typical usage: Temporary filesystems

Filesystems on Lemaitre3



Filesystems on Lemaitre3



Volume

```
[dfr@lemaitre3 ~]$ df -khT -x tmpfs
Filesystem                                Type      Size  Used Avail Use% Mounted on
/dev/mapper/centos_controller-root        ext4      493G  372G   96G   80% /
devtmpfs                                  devtmpfs  23G    0    23G    0% /dev
/dev/sda2                                  ext4      976M  117M  793M   13% /boot
/dev/sda1                                  vfat      50M   9.9M   41M   20% /boot/efi
vsalt:/srv/salt                           nfs4      9.8G  532M   8.7G    6% /srv/salt
vsalt:/srv/pillar                         nfs4      9.8G  532M   8.7G    6% /srv/pillar
beegfs_nodev                              beegfs    583T  203T  380T   35% /scratch
/dev/drbd1                                 xfs       3.2T  162G   3.0T    6% /trinity
gw-ucl:/CECI/gateway/home                 nfs       40T    2.5T   38T    7% /CECI/home
lm3-n:/storage                            nfs4      35T    9.1T   26T   27% /home
gw-ucl:/CECI/gateway/soft/cecisw/RedHat-7_6-85-4_Omnipath nfs      1.0T  498G  527G   49% /opt/cecisw/arch
gw-ucl:/CECI/gateway/soft                 nfs      1.0T  498G  527G   49% /CECI/soft
```

Source:

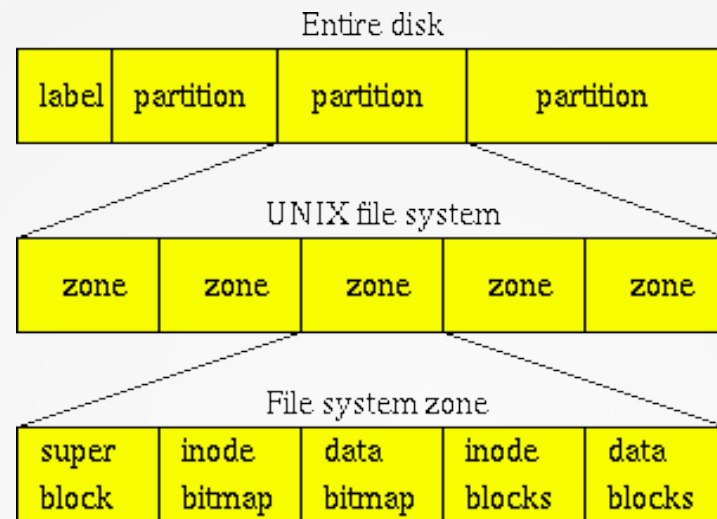
- /dev/sd... → local disk
- /dev/mapper... → LVM
- <machine>:<path> → NFS
- other (e.g. beegfs_nodev) → specific filesystem

Local mount point

Inodes

```
[dfr@lemaitre3 ~]$ df -ikhT -x tmpfs
Filesystem                                Type      Inodes  IUsed  IFree  IUse% Mounted on
/dev/mapper/centos_controller-root        ext4       32M   816K   31M    3% /
devtmpfs                                  devtmpfs   5.8M   506    5.8M    1% /dev
/dev/sda2                                  ext4       64K    32    64K    1% /boot
/dev/sda1                                  vfat        0        0        0     - /boot/efi
vsalt:/srv/salt                           nfs4      640K   13K   628K    3% /srv/salt
vsalt:/srv/pillar                         nfs4      640K   13K   628K    3% /srv/pillar
beegfs_nodev                              beegfs        0        0        0     - /scratch
/dev/drbd1                                 xfs       320M  681K  319M    1% /trinity
gw-ucl:/CECI/gateway/home                 nfs       98M    14M   85M   14% /CECI/home
lm3-n:/storage                            nfs4      51G   39M   51G    1% /home
gw-ucl:/CECI/gateway/soft/cecisw/RedHat-7_6-85-4_Omnipath nfs       98M    14M   85M   14% /opt/cecisw/arch
gw-ucl:/CECI/gateway/soft/cecisw/noarch   nfs       98M    14M   85M   14% /opt/cecisw/noarc
gw-ucl:/CECI/gateway/soft                 nfs       98M    14M   85M   14% /CECI/soft
```

A word about inodes (simplified)



directory /home/you

foo	123
bar	456
and so on...	

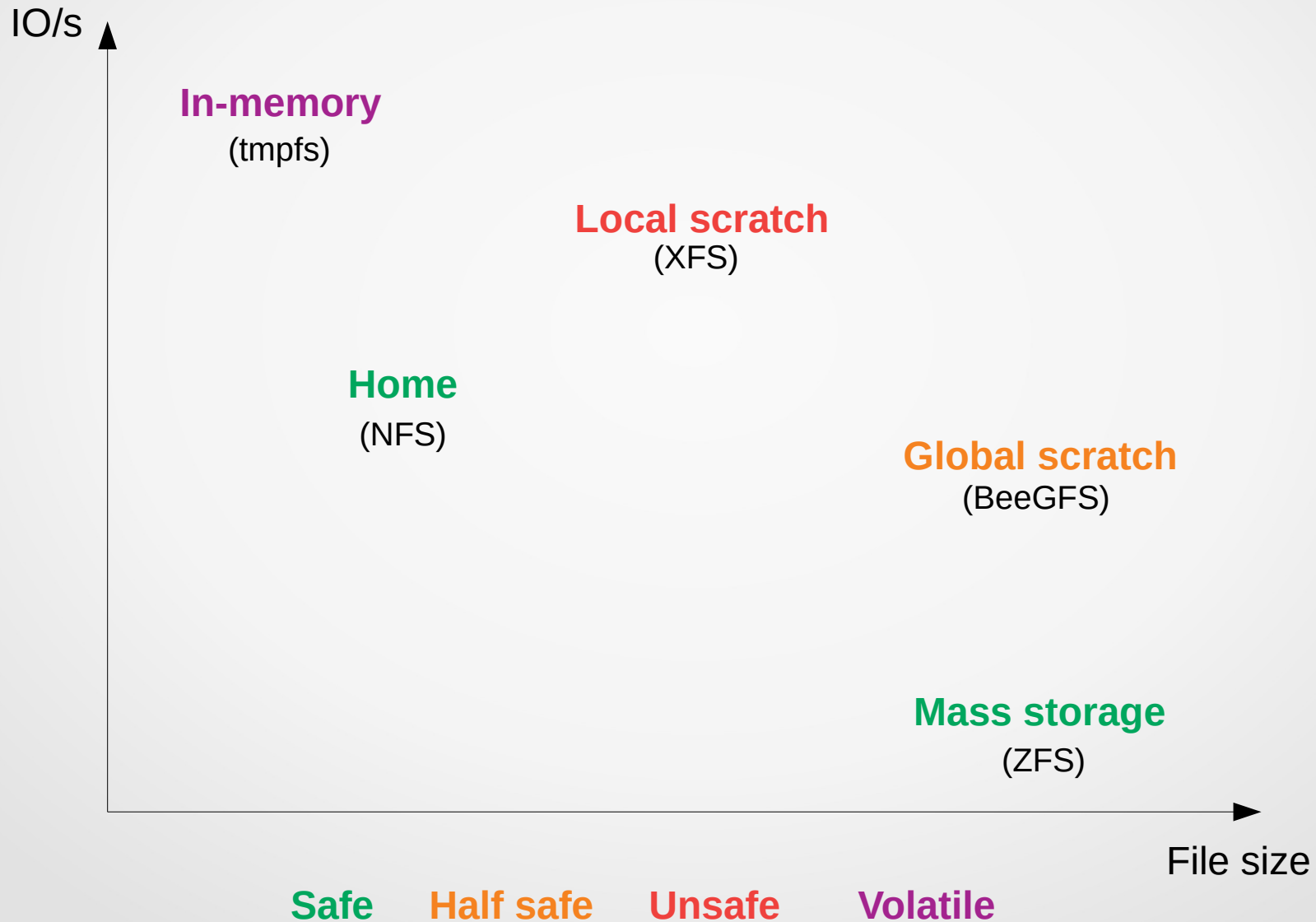
inode 123

owner/group ID
permissions
file/directory/etc.
data block #s
and so on...

blocks...

data data
data data
data data

What filesystem for what usage



File formats

Text File Formats – CSV

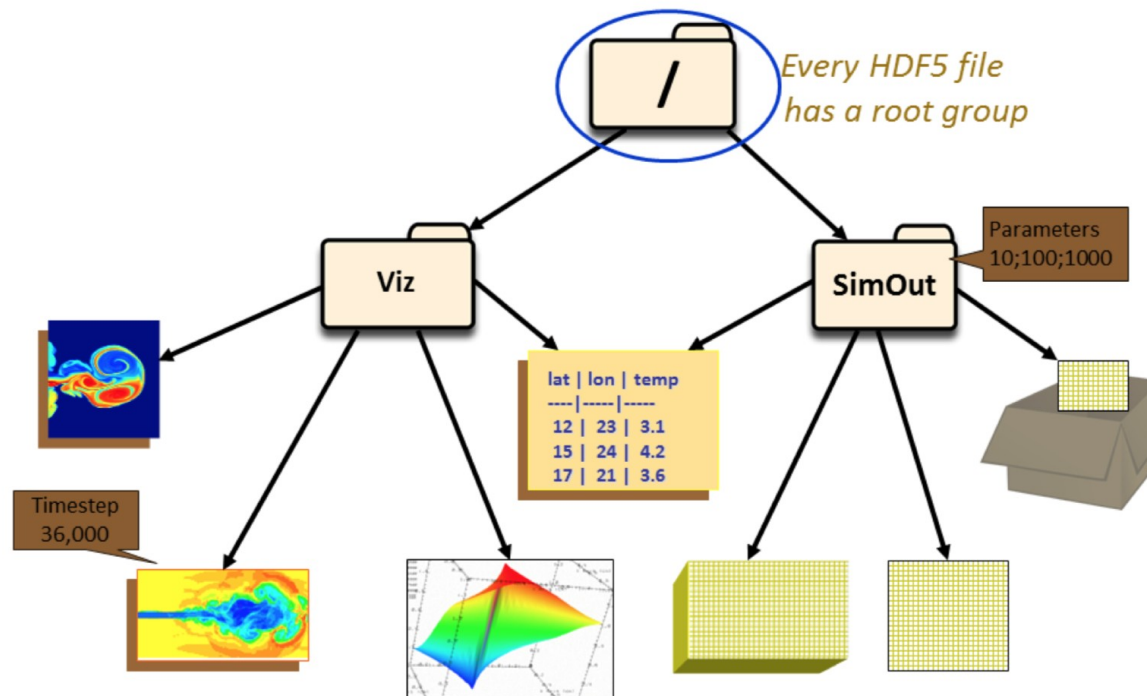
```
processmi.txt - Notepad
File Edit Format View Help
__NounName Name Handles VM WS PM NPM Path Company CPU Fileversion ProductVer
Process AcDeskBandHlpr 127 85278720 10047488 4362240 14224 C:\Program Files (x86)\Le
Process AcPrfMgrSvc 167 72732672 11051008 4956160 17128 C:\Program Files (x86)\Le
Process AcSvc 421 127533056 17551360 10117120 26960 C:\Program Files (x86)\Le
Process armsvc 84 45129728 4509696 1601536 8528 C:\Program Files (x86)\Le
Process btwdins 127 60628992 7720960 3743744 10096 C:\Program Files (x86)\Le
Process CamMute 111 40837120 5312512 1789952 9440 C:\Program Files (x86)\Le
Process CcmExec 1433 329342976 89608192 58417152 58432 C:\Program Files (x86)\Le
Process conhost 35 25186304 4136960 2048000 5168 C:\Program Files (x86)\Le
Process conhost 45 52461568 6311936 3215360 6360 C:\windows\system32\conhost.exe M
Process csrss 1145 55078912 5623808 2977792 17768 C:\Program Files (x86)\Le
Process csrss 628 96055296 17383424 20799488 28704 C:\Program Files (x86)\Le
Process CxAudMsg64 98 57937920 6680576 7385088 8832 C:\Program Files (x86)\Le
Process daemonu 413 73506816 8957952 5935104 20132 C:\Program Files (x86)\Le
Process DcaSvc 642 64344064 13996032 10178560 27904 C:\Program Files (x86)\Le
Process DcaTray 644 220614656 35368960 33832960 34628 C:\Program Files (x86)\DirectA
Process dwm 108 84504576 10313728 5709824 10024 C:\windows\system32\Dwm.exe Micros
Process EvtEng 284 105017344 20992000 12906496 22248 C:\Program Files (x86)\Le
Process EXCEL 384 316968960 42233856 27746304 48584 C:\Program Files (x86)\Microsof
Process explorer 1153 352845824 75698176 52506624 94040 C:\windows\Explorer.EXE Mi
Process fmapp 30 53719040 6168576 3743744 6960 C:\Program Files\CONEXANT\ForteConfig
Process FwcAgent 555 59432960 12111872 9289728 22248 C:\Program Files (x86)\Le
Process hkcmd 84 71577600 8040448 3964928 8048 C:\windows\system32\hkcmd.exe Intel
Process ibmpmsvc 61 44564480 4259840 2301952 6568 C:\Program Files (x86)\Le
Process Idle 0 0 24576 0 0 0 C:\windows\system32\Idle.exe
Process iexplore 774 304881664 79278080 77529088 82152 C:\Program Files (x86)\Intern
Process iexplore 639 157421568 22876160 12095488 38648 C:\Program Files (x86)\Intern
Process igfxpers 203 83292160 10944512 5844992 10808 C:\windows\System32\igfxpers.ex
Process LMS 112 41308160 5484544 2117632 9800 C:\Program Files (x86)\Le
Process lsass 1545 75075584 27377664 18477056 47272 C:\Program Files (x86)\Le
Process lsm 213 19779584 5664768 3756032 7600 C:\Program Files (x86)\Le
Process micmute 110 48328704 5672960 6172672 11120 C:\Program Files (x86)\Le
Process mscorsvw 105 73936896 14417920 8912896 11184 C:\Program Files (x86)\Le
Process MsitBLSHA 118 45617152 8912896 4481024 10328 C:\Program Files (x86)\Le
Process msitcertsvc 712 565702656 33861632 40435712 41284 C:\Program Files (x86)\Le
Process MsitTpmSvc 84 43319296 10129408 6975488 7568 C:\Program Files (x86)\Le
Process MsMpEng 503 238940160 87953408 117288960 47112 C:\Program Files (x86)\Le
Process MSOIDSVC 612 106561536 21938176 14983168 27648 C:\Program Files (x86)\Le
Process MSOIDSVC 72 36134912 4841472 2621440 6080 C:\Program Files (x86)\Le
Process msseces 341 156336128 23216128 10706944 25232 C:\Program Files\Microsoft Sec
Process NisSrv 261 80003072 4771840 9138176 18000 C:\Program Files (x86)\Le
Process nus3mon 89 75804672 5914624 2224128 10328 C:\Program Files (x86)\Renesas El
```

Binary File Formats – CDF, HDF



September 23, 2016

Introduction to HDF5



There are two groups in the HDF5 file depicted above: Viz and SimOut. Under the Viz group are a variety of images and a table that is shared with the SimOut group. The SimOut group contains a 3-dimensional array, a 2-dimensional array and a link to a 2-dimensional array in another HDF5 file.

Binary File Formats – CDF, HDF



```
/* *****  
 * Copyright by The HDF Group. *  
 * Copyright by the Board of Trustees of the University of Illinois. *  
 * All rights reserved. *  
 * *  
 * This file is part of HDF5. The full HDF5 copyright notice, including *  
 * terms governing use, modification, and redistribution, is contained in *  
 * the COPYING file, which can be found at the root of the source code *  
 * distribution tree, or in https://support.hdfgroup.org/ftp/HDF5/releases. *  
 * If you do not have access to either file, you may request a copy from *  
 * help@hdfgroup.org. *  
 * ***** */  
  
/*  
 * This example illustrates how to write and read data in an existing  
 * dataset. It is used in the HDF5 Tutorial.  
 */  
  
#include "hdf5.h"  
#define FILE "dset.h5"  
  
int main() {  
  
    hid_t    file_id, dataset_id; /* identifiers */  
    herr_t   status;  
    int      i, j, dset_data[4][6];  
  
    /* Initialize the dataset. */  
    for (i = 0; i < 4; i++)  
        for (j = 0; j < 6; j++)  
            dset_data[i][j] = i * 6 + j + 1;  
  
    /* Open an existing file. */  
    file_id = H5Fopen(FILE, H5F_ACC_RDWR, H5P_DEFAULT);  
  
    /* Open an existing dataset. */  
    dataset_id = H5Dopen2(file_id, "/dset", H5P_DEFAULT);  
  
    /* Write the dataset. */  
    status = H5Dwrite(dataset_id, H5T_NATIVE_INT, H5S_ALL, H5S_ALL, H5P_DEFAULT,  
                     dset_data);  
  
    status = H5Dread(dataset_id, H5T_NATIVE_INT, H5S_ALL, H5S_ALL, H5P_DEFAULT,  
                    dset_data);  
  
    /* Close the dataset. */  
    status = H5Dclose(dataset_id);  
  
    /* Close the file. */  
    status = H5Fclose(file_id);  
}
```

Binary File Formats – CDF, HDF



```
dfr@hmem00:~/hdf5 $ cat res.txt
1014
7795
19769
16872
11252
22757
1773
28983
22600
27925
dfr@hmem00:~/hdf5 $ cat res.h5conf
PATH res
INPUT-CLASS TEXTFP
RANK 1
DIMENSION-SIZES 10
OUTPUT-CLASS FP
OUTPUT-SIZE 64
OUTPUT-ARCHITECTURE IEEE
OUTPUT-BYTE-ORDER LE
dfr@hmem00:~/hdf5 $ h5import res.txt -c res.h5conf -o res.hf5
dfr@hmem00:~/hdf5 $ h5dump res.hf5
HDF5 "res.hf5" {
  GROUP "/" {
    DATASET "res" {
      DATATYPE  H5T_IEEE_F64LE
      DATASPACE  SIMPLE { ( 10 ) / ( 10 ) }
      DATA {
        (0): 1014, 7795, 19769, 16872, 11252, 22757, 1773, 28983, 22600, 27925
      }
    }
  }
}
```

What file format for what usage



- Meta data
 - Configuration file: INI, YAML
 - Result with context information: JSON
- Data
 - Small data (kB): CSV
 - Medium data (MB): compressed CSV
 - Large data (GB): netCDF, HDF5, DXMF
 - Huge data (TB): Database, Object store

("loss of innocence")

Use dedicated libraries to write and read them

1.2 Object storage

Object storage

- Object: data (e.g. file) + **custom** meta data



- Often built on erasure coding (“software RAID”)
- Scale out easily
- Useful for web applications but coming to scientific world
- Access with REST API (through HTTP)

S3 Python example



```
import boto.s3.connection

access_key = [REDACTED]
secret_key = [REDACTED]
conn = boto.connect_s3(
    aws_access_key_id=access_key,
    aws_secret_access_key=secret_key,
    host='192.168.64.51', port=7480,
    is_secure=False,
    calling_format=boto.s3.connection.OrdinaryCallingFormat(),
)

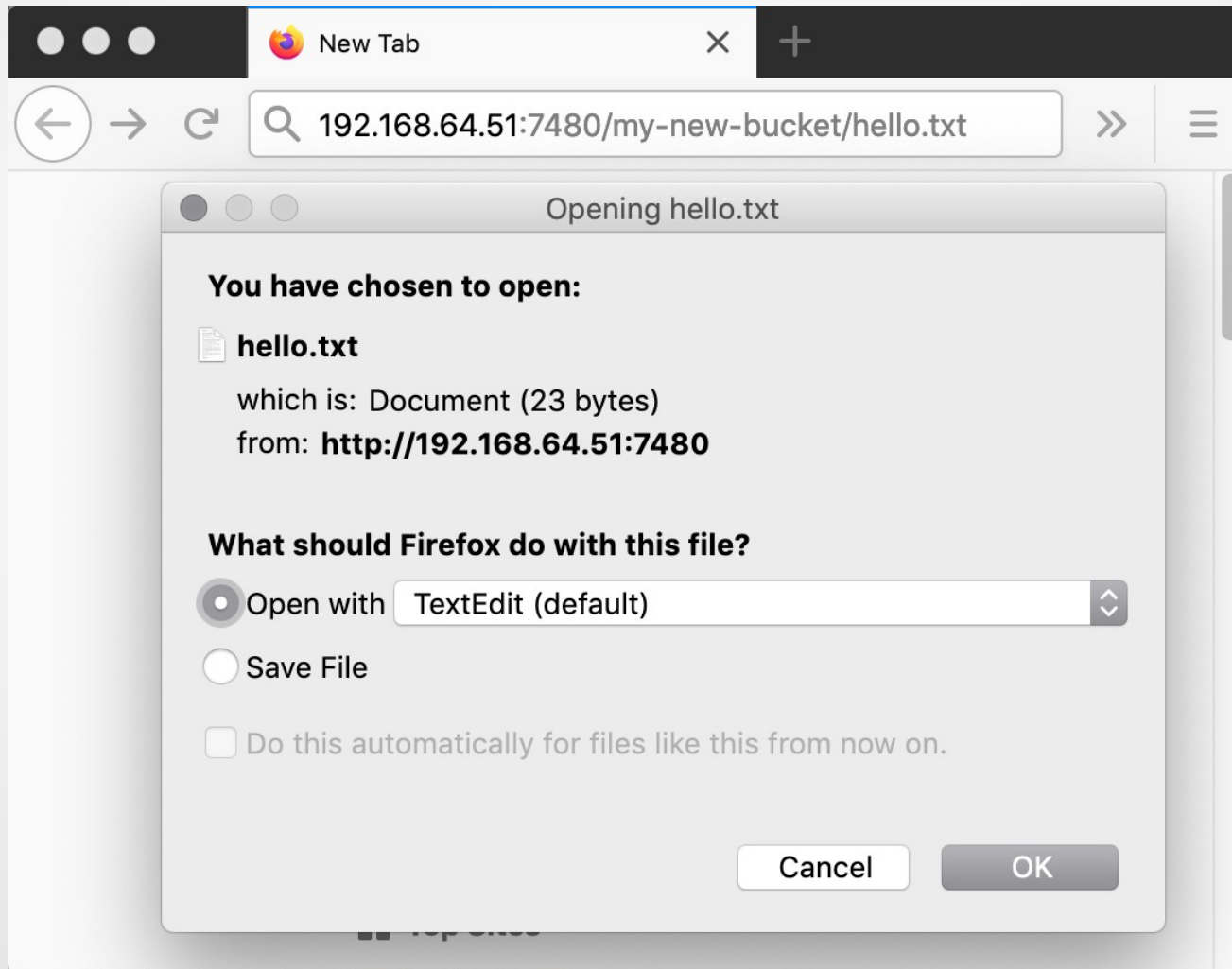
bucket = conn.create_bucket('my-new-bucket')
```

```
## Create an object
key = bucket.new_key('hello.txt')
key.set_contents_from_string('Hello file in CEPH S3!')
key.set_metadata('Owner', 'dfr')

key = bucket.new_key('smstoolsS3-2.2.20.tar.gz')
key.set_contents_from_filename('smstools-2.2.20.tar.gz')

hello_key = bucket.get_key('hello.txt')
hello_url = hello_key.generate_url(3600, query_auth=False, force_http=False)
```


S3 Python example



S3 Object tagging (meta data)



Object Tagging

Use object tagging to categorize storage. Each tag is a key-value pair. Consider the following tagging examples:

- Suppose an object contains protected health information (PHI) data. You might tag the object using the following key-value pair, as shown following:

```
PHI=True
```



or

```
Classification=PHI
```



- Suppose you store project files in your S3 bucket. You might tag these objects with a key called `Project` and a value, as shown following:

```
Project=Blue
```



- You can add multiple tags to an object, as shown following:

```
Project=x  
Classification=confidential
```



1.3 Databases

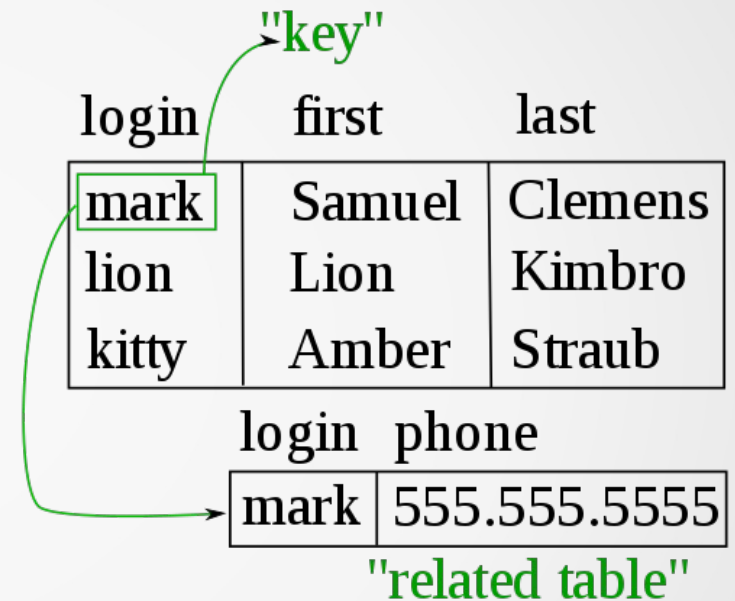
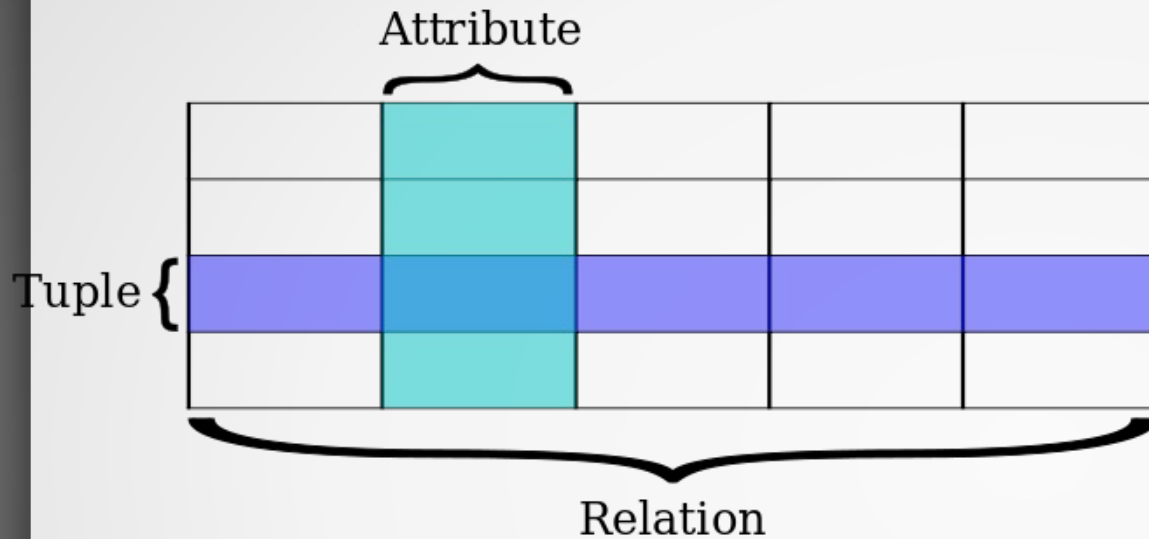
RDBMS



- Mostly needed for categorical data and alphanumeric data (not suited for matrices, but good for end-results)
- Indexes make finding a data element is very fast (and computing sums, maxima, etc.)
- Encodes relations between data (constraints, etc)
- Atomicity, Consistency, Isolation, and Durability



Tables and SQL



```
create table Users (login varchar(255), first varchar(255), last varchar(255));
```

```
insert into Users values ("mark", "Samuel", "Clemens");
```

```
select first,last from Users where login='lion';
```

```
select login, phone from Users join PhoneNb on Users.login=PhoneNb.login;
```

File-based RDBMS



- The features of a relational database without the need for a complete setup
- Simply based on files
- Command line interface + API (Python, etc.)
- Still much more efficient than writing a million small files

NoSQL



key-value

Amazon
DynamoDB (Beta)

ORACLE
BERKELEY DB **11g**



redis

graph



column



document



- Mostly needed for unstructured, semi-structured, and polymorphic data
- Scaling out very easy
- Basic Availability, Soft-state, Eventual consistency

TinyDB

- The features of a relational database without the need for a complete setup
- Simply based on files
- Command line interface + API (Python, etc.)
- Still much more efficient than writing a million small files

When to use a database?



- when you have a large number of small files
- when you perform a lot of direct writes in a large file
- when you want to keep structure/relations between data
- when software crashes have a non-negligible probability
- when files are updated by several processes

Example: Danger of NFS

A screenshot of a web browser window displaying the 'Linux NFS faq' page from nfs.sourceforge.net. The browser's address bar shows the URL 'nfs.sourceforge.net'. The page content includes a reference to Callaghan's 'NFS Illustrated', followed by question A9: 'Why does opening files with O_APPEND on multiple clients cause the files to become corrupted?'. The answer explains that NFS does not support atomic append writes and discusses cache consistency models. It also includes a link to read more about the NFS cache consistency model and an alternative NFS protocol feature. Question A10 is partially visible at the bottom: 'What does it mean when my application fails because of an ESTALE error?'.

see Callaghan's "[NFS Illustrated](#)."

A9. Why does opening files with `O_APPEND` on multiple clients cause the files to become corrupted?

A. The NFS protocol does not support atomic append writes, so append writes are *never* atomic on NFS for any platform.

Most NFS clients, including the Linux NFS client in kernels newer than 2.4.20, support "close to open" cache consistency, which provides good performance and meets the sharing needs of most applications. This style of cache consistency does not provide strict coherence of the file size attribute among multiple clients, which would be necessary to ensure that append writes are always placed at the end of a file.

Read all about the NFS cache consistency model [here](#).

Alternately, the NFS protocol could include a specific atomic append write operation, but today's versions of the protocol do not. The designers of the NFS protocol felt that atomic append writes would be rarely used, so they never added the feature. Even with such a feature, keeping the file size attribute up to date would be challenging.

A10. What does it mean when my application fails because of an ESTALE error?

A. The NFS protocol does not refer to files and directories by name or by path; it

Example: Danger of NFS



```
dfr@hmem00
x dfr@hmem00 %1
dfr@hmem00:~ $ salloc --nodes 4
salloc: Granted job allocation 1126834
dfr@hmem00:~ $ srun hostname
hmem17.cism.ucl.ac.be
hmem10.cism.ucl.ac.be
hmem11.cism.ucl.ac.be
hmem13.cism.ucl.ac.be
dfr@hmem00:~ $ srun bash -c 'for ((i=0 ; i<99999 ; i++)) ; do printf "$SLURM_PROCID %06d\n" $i >> testfile ; done'
dfr@hmem00:~ $ head -3 testfile
0 000000
0 000001
0 000002
dfr@hmem00:~ $ egrep -v '^[0-9 ]{8}$' testfile
0 0051 005959
2 0030720 012016
4701
1 00 026581
0 028921 2 021730
1 033380 031932
0 03382 01 03540 02 02721 00 033897
1 0361610 034566
32 0305901 038633
3 0051 039651
3 0 038342 1 040169
0 2 032693 007499
0 0432 038297
1 050 0533 033796
2 00 06613 056720
2 06401 03 058000
1 07178030 0673361 071781
3 01 07963 00 074553
2 1 080223 072629
```

When NOT to use a database?

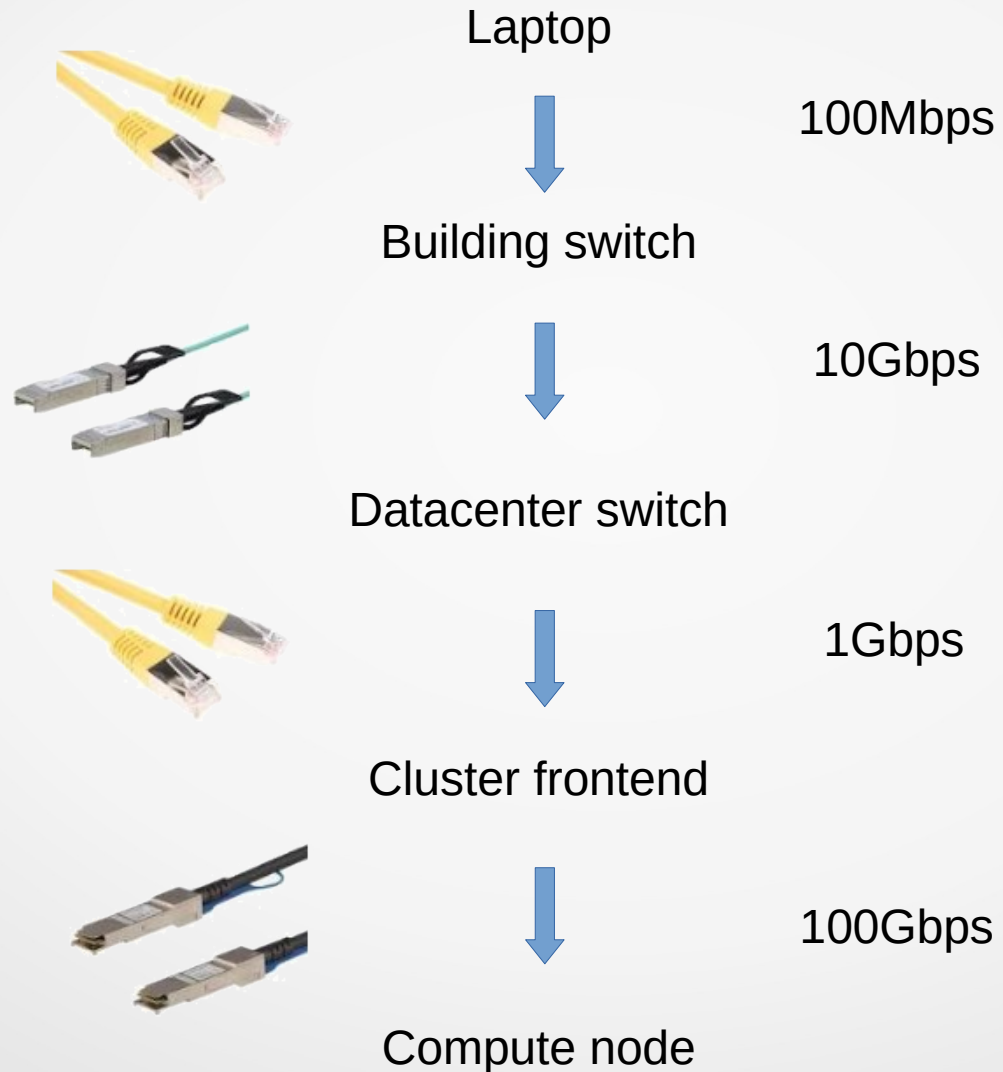


- only sequential access
- simple matrices/vectors, etc.
- direct access on fixed-size records and no structure

2.

Data transfer

Network technologies



SCP



```
scp local_file.txt remote_host:/home
```

- Most direct way to copy data from/to UNIX/Linux machines
- Inefficient (sequential and synchronous)
- “outdated, inflexible and not readily fixed” (OpenSSH 8.0 release notes)

SFTP



```
dfr@ncois:~ $ sftp lm3
Connected to lm3.
sftp> ?
Available commands:
bye                               Quit sftp
cd path                           Change remote directory to 'path'
chgrp grp path                    Change group of file 'path' to 'grp'
chmod mode path                  Change permissions of file 'path' to 'mode'
chown own path                   Change owner of file 'path' to 'own'
df [-hi] [path]                  Display statistics for current directory or
                                filesystem containing 'path'
exit                              Quit sftp
get [-afPpRr] remote [local]     Download file
reget [-fPpRr] remote [local]    Resume download file
reput [-fPpRr] [local] remote    Resume upload file
help                             Display this help text
lcd path                          Change local directory to 'path'
lls [ls-options] [path]         Display local directory listing
lmkdir path                       Create local directory
ln [-s] oldpath newpath         Link remote file (-s for symlink)
lpwd                              Print local working directory
ls [-lafhlNrSt] [path]          Display remote directory listing
lumask umask                     Set local umask to 'umask'
mkdir path                        Create remote directory
progress                          Toggle display of progress meter
put [-afPpRr] local [remote]    Upload file
pwd                               Display remote working directory
quit                              Quit sftp
rename oldpath newpath          Rename remote file
rm path                          Delete remote file
rmdir path                       Remove remote directory
symlink oldpath newpath         Symlink remote file
version                          Show SFTP version
!command                        Execute 'command' in local shell
!                                Escape to local shell
?                                Synonym for help
sftp> █
```


RSYNC



```
rsync(1) rsync(1)
NAME
    rsync - faster, flexible replacement for rcp
SYNOPSIS
    rsync [OPTION]... SRC [SRC]... DEST
    rsync [OPTION]... SRC [SRC]... [USER@]HOST:DEST
    rsync [OPTION]... SRC [SRC]... [USER@]HOST::DEST
    rsync [OPTION]... SRC [SRC]... rsync://[USER@]HOST[:PORT]/DEST
    rsync [OPTION]... SRC
    rsync [OPTION]... [USER@]HOST:SRC [DEST]
    rsync [OPTION]... [USER@]HOST::SRC [DEST]
    rsync [OPTION]... rsync://[USER@]HOST[:PORT]/SRC [DEST]
DESCRIPTION
    rsync is a program that behaves in much the same way that rcp does, but has many more
    options and uses the rsync remote-update protocol to greatly speed up file transfers
    when the destination file is being updated.

    The rsync remote-update protocol allows rsync to transfer just the differences
    between two sets of files across the network connection, using an efficient checksum-
    search algorithm described in the technical report that accompanies this package.

    Some of the additional features of rsync are:
    o support for copying links, devices, owners, groups, and permissions
    o exclude and exclude-from options similar to GNU tar
    o a CVS exclude mode for ignoring the same files that CVS would ignore
    o can use any transparent remote shell, including ssh or rsh
    o does not require super-user privileges
    o pipelining of file transfers to minimize latency costs
    o support for anonymous or authenticated rsync daemons (ideal for mirroring)
GENERAL
    Rsync copies files either to or from a remote host, or locally on the current host
    (it does not support copying files between two remote hosts).

    There are two different ways for rsync to contact a remote system: using a remote-
    shell program as the transport (such as ssh or rsh) or contacting an rsync daemon
    directly via TCP. The remote-shell transport is used whenever the source or destina-
    tion path contains a single colon (:) separator after a host specification. Contact-
```

Update only what changed: rsync



```
rsync [OPTIONS]... SRC [SRC]... [USER@]HOST:DEST
```

- Always use: -az
- Other interesting arguments:
 - -v and --progress
 - --include or --exclude
 - --delete and/or --remove-source-file
 - --dry-run
 - --size-only or --checksum
- Works well with GNU parallel

Resuming transfers



- When nothing changed but the transfer was interrupted
 - append: do not re-check partially transmitted files and resume the transfer where it was abandoned assuming first transfer attempt was with scp or with rsync --inplace

```
dfr@manneback — bash
dfr@manneback:~$ scp Accelerators.tgz lm9:/tmp
Accelerators.tgz          58% 323MB 66.7MB/s  00:03 ETA^
CKilled by signal 2.
dfr@manneback:~$ time rsync Accelerators.tgz lm9:/tmp

real    0m8.846s
user    0m3.633s
sys     0m0.150s
dfr@manneback:~$ scp Accelerators.tgz lm9:/tmp
Accelerators.tgz          54% 299MB 80.5MB/s  00:03 ETA^
CKilled by signal 2.
dfr@manneback:~$ time rsync --append Accelerators.tgz lm9:/tmp

real    0m1.074s
user    0m0.123s
sys     0m0.020s
dfr@manneback:~$
```

- From local to remote

```
$ tar zvf - /path/to/data | ssh server "cat > /srv/data_server1.tar.gz"
```

- From remote to local

```
$ ssh server tar czf - /path/to/data/ > ./data_server.tar.gz
```

- Avoid a lot of communication overhead linked to inodes
- Use the pv command to get a progress bar

Parallel data transfer: bbcp



- Better use of the bandwidth than SCP
- Needs to be installed on both sides (easy to install)
- Needs friendly firewalls

```
dfr@manneback:~$ time scp Accelerators.tgz lm9:/dev/null
Accelerators.tgz
100% 551MB 50.1MB/s 00:11
real    0m13.849s
user    0m5.043s
sys     0m2.330s
dfr@manneback:~$ time bbcp -P 2 Accelerators.tgz lm9:/dev/null/t
bbcp: Creating /dev/null/t
bbcp: 151113 22:06:00 39% done; 112.3 MB/s
bbcp: 151113 22:06:02 79% done; 112.0 MB/s
real    0m9.000s
user    0m0.232s
sys     0m0.926s
```

Parallel rsync



- Parsyncfp

```
[root@compute-3-11 calof]# parsyncfp --NP=8 --chunk=10G -i ib0 --startdir /mnt/calof lander-calof /dfs3/staff/hmangala/
INFO: You've specified what looks like an Infiniband interface [ib0]...
INFO: .. and you have 'perfquery installed, so RDMA bytes will be reported as well.
WARN: About to remove all the old cached chunkfiles from [/root/.parsyncfp/fpcache].
Enter ^C to stop this.
If you specified '--nowait', cache will be cleared in 3s regardless.
Otherwise, hit [Enter] and I'll clear them.
Press [ENTER] to continue.

INFO: The fpart chunk files [/root/.parsyncfp/fpcache/f*] are cleared .. continuing.
INFO: Forking fpart. Check [/root/.parsyncfp/fpcache/fpart.log.20.10.31_2018-12-12] for errors if it hangs.
INFO: Starting the 1st [8] rsyncs ..
```

Time	Elapsed time(m)	1m Load	[ib0] TCP / RDMA	MB/s out	Running PIDs	Susp'd PIDs	Chunks [UpTo] of [ToDo]	[2018-12-12]
20.10.39	0.07	3.00	0.00 / 132.12		5	<> 0	[8] of [50]	
20.10.43	0.12	3.08	0.00 / 426.94		8	<> 0	[13] of [50]	
20.10.47	0.18	3.40	0.00 / 818.96		8	<> 0	[15] of [50]	
20.10.50	0.23	3.40	0.00 / 848.65		8	<> 0	[15] of [50]	
20.10.54	0.30	3.77	0.00 / 835.68		8	<> 0	[15] of [50]	
20.10.57	0.35	3.86	0.00 / 781.66		8	<> 0	[15] of [50]	
20.11.01	0.42	3.86	0.00 / 734.39		8	<> 0	[15] of [50]	
20.11.04	0.47	4.12	0.00 / 618.96		8	<> 0	[15] of [50]	
20.11.08	0.53	4.51	0.00 / 572.82		8	<> 0	[15] of [50]	
20.11.11	0.58	4.51	0.00 / 534.97		8	<> 0	[15] of [50]	
20.11.15	0.65	4.79	0.00 / 537.75		8	<> 0	[15] of [50]	
20.11.18	0.70	5.12	0.00 / 525.65		8	<> 0	[15] of [50]	

Parallel rsync



- GNU Parallel + rsync

EXAMPLE: Parallelizing rsync

rsync is a great tool, but sometimes it will not fill up the available bandwidth. Running multiple **rsync** in parallel can fix this.

```
cd src-dir
find . -type f |
  parallel -j10 -X rsync -zR -Ha ./{} fooserver:/dest-dir/
```

Adjust **-j10** until you find the optimal number.

rsync -R will create the needed subdirectories, so all files are not put into a single dir. The **/** is needed so the resulting command looks similar to:

```
rsync -zR ../sub/dir/file fooserver:/dest-dir/
```

The **/** is what **rsync -R** works on.

If you are unable to push data, but need to pull them and the files are called digits.png (e.g. 000000.png) you might be able to do:

```
seq -w 0 99 | parallel rsync -Havessh fooserver:src/*{}.png destdir/
```

Limit the depth of `find` can help speed things up

3. Data sharing



with other users (Unix permissions, Encryption)
with external users (Owncloud, Dataverse)

Data sharing with other users

Sharing with the group

```
dfr@manneback — bash
dfr@manneback:~/share $ id
uid=106(dfr) gid=205(grppan) groups=205(grppan),236(grparcv),277(grpvlor)
dfr@manneback:~/share $ mkdir testdir
dfr@manneback:~/share $ chmod 750 testdir
dfr@manneback:~/share $ ll
total 0
drwxr-x--- 2 dfr grppan 6 Nov 16 14:06 testdir
dfr@manneback:~/share $ chgrp grparcv testdir/
dfr@manneback:~/share $ ll
total 0
drwxr-x--- 2 dfr grparcv 6 Nov 16 14:06 testdir
dfr@manneback:~/share $ touch testdir/testfile1
dfr@manneback:~/share $ ll testdir/
total 0
-rw-r--r-- 1 dfr grppan 0 Nov 16 14:07 testfile1
dfr@manneback:~/share $ chmod g+s testdir/
dfr@manneback:~/share $ touch testdir/testfile2
dfr@manneback:~/share $ ll testdir/
total 0
-rw-r--r-- 1 dfr grppan 0 Nov 16 14:07 testfile1
-rw-r--r-- 1 dfr grparcv 0 Nov 16 14:07 testfile2
dfr@manneback:~/share $
```

Sharing and hiding

```
dfr@manneback — bash
dfr@manneback:~$ ls -ld
drwxr-x--x 144 dfr grppan 12288 Nov 16 14:22 .
dfr@manneback:~$ ls -l minimal.c
-rw-r--r-- 1 dfr grppan 43 May 22 10:38 minimal.c
dfr@manneback:~$ cat minimal.c
int main()
{
    int i = 1337;
    return 0;
}
dfr@manneback:~$

tuto01@manneback:~$ ls -ldfr
ls: cannot open directory /home/pan/dfr:
Permission denied
tuto01@manneback:~$ cat -dfr/minimal.c
int main()
{
    int i = 1337;
    return 0;
}
tuto01@manneback:~$
```

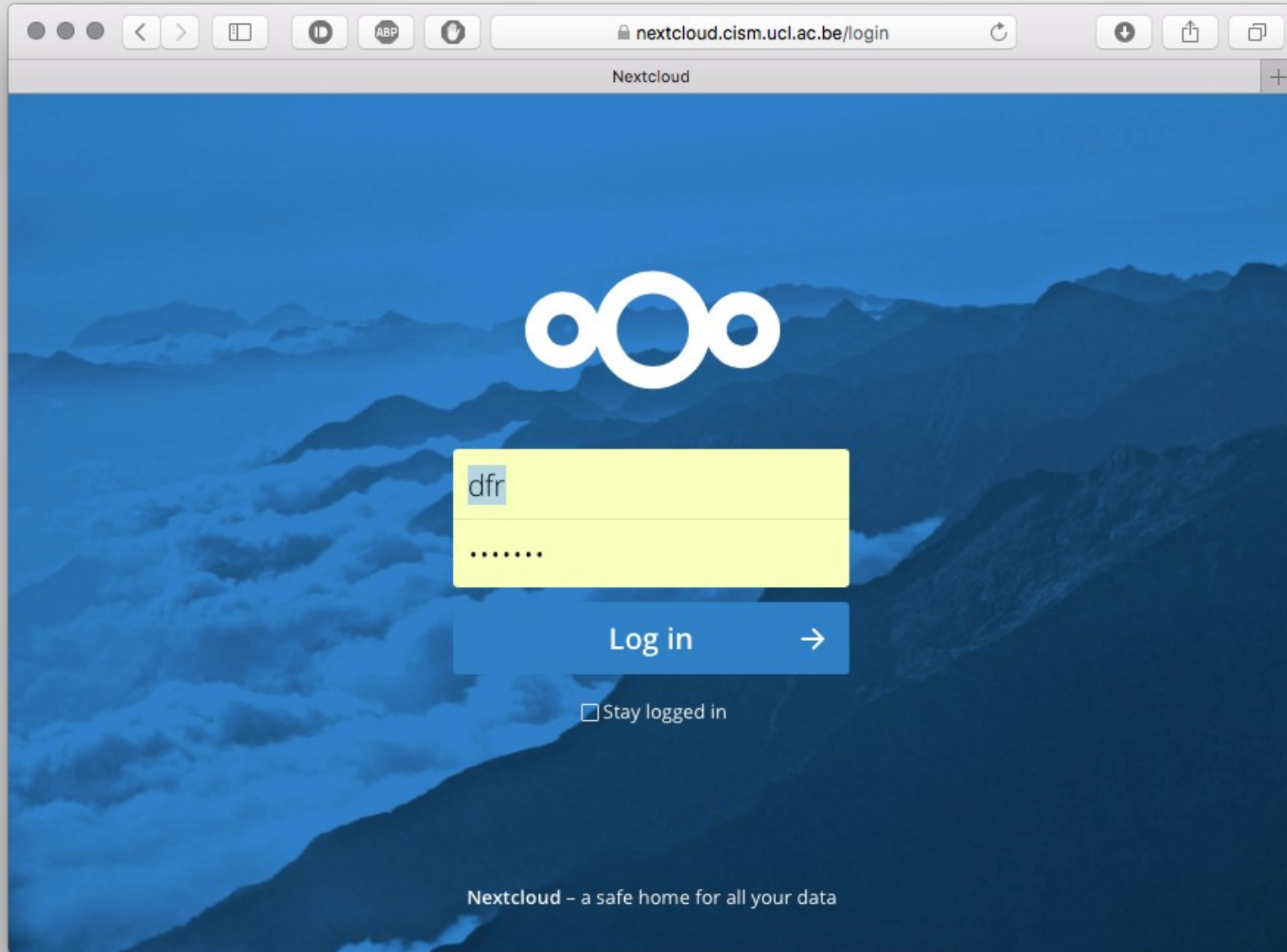
[0] 0:bash* "dfr@manneback" 14:26 16-Nov-15

Sharing and encrypting

```
dfr@manneback — bash
dfr@manneback:~$ cat minimal.c
int main()
{
    int i = 1337;
    return 0;
}
dfr@manneback:~$ openssl des3 -salt -in minimal.c -out minimal.c.des3
enter des-ede3-cbc encryption password:
Verifying - enter des-ede3-cbc encryption password:
dfr@manneback:~$ cat minimal.c.des3
Salted__0?1c_?RjD7
?      ??f?c???????X?B?6:=???????!h??.,???ni?dfr@manneback:~$
dfr@manneback:~$ openssl des3 -d -in minimal.c.des3 -out minimal.c.clear
enter des-ede3-cbc decryption password:
dfr@manneback:~$ cat minimal.c.clear
int main()
{
    int i = 1337;
    return 0;
}
dfr@manneback:~$
```

Data sharing with external users

Data sharing with external users



CISM
login

Dropbox-like



The screenshot shows the ownCloud Files interface. The browser address bar displays the URL `https://hall.cism.ucl.ac.be/owncloud/index.php/apps/files/`. The interface includes a search bar, a navigation menu on the left, and a main file list area. The file list is organized into sections: 'All files', 'Shared with you', 'Shared with others', 'Shared by link', and 'External storage'. The 'All files' section contains a table of files and folders.

Name	Size	Modified
documents	34 kB	6 months ago
green	119.3 GB	7 months ago
manneback	Pending	2 minutes ago
music	3.6 MB	9 months ago
photos	45.8 MB	6 months ago
storage	<1 kB	7 months ago
duplicity-full-signatures.20150204T143855Z.sigtar.gpg	<1 kB	7 months ago
New Document.odt	8 kB	6 months ago
ownCloudUserManual.pdf	1.7 MB	9 months ago
slides.tgz	634 kB	9 months ago

At the bottom of the interface, a status bar indicates: "Loading 'https://hall.cism.ucl.ac.be/owncloud/index.php/apps/files/', completed 16 of 17 items".

External SFTP connectors



The screenshot shows the 'ownCloud' settings page for 'External Storage'. The browser address bar indicates the URL: `https://hall.cism.ucl.ac.be/owncloud/index.php/settings/personal`. The page title is 'ownCloud' and the user is logged in as 'dfr'. The 'External Storage' section contains a table with three entries, each with a green status indicator.

Folder name	External storage	Configuration
green	SFTP	green.cism.ucl.ac.be dfr /home/pan/dfr
manneback	SFTP	manneback.cism.ucl.a dfr /home/pan/dfr
storage	SFTP	storage.cism.ucl.ac.be dfr /home/pan/dfr

Below the table, there is an input field for 'Folder name' and an 'Add storage' button. The 'SSL root certificates' section below shows a 'Choose File' button with 'no file selected' and an 'Import Root Certificate' button.

Dropbox-like



The screenshot shows the ownCloud Files interface in a browser window. The address bar displays the URL `https://hall.cism.ucl.ac.be/owncloud/index.php/apps/files/`. The interface includes a search bar, a navigation sidebar on the left, and a main file list area. The file list is organized into sections: 'All files', 'Shared with you', 'Shared with others', 'Shared by link', and 'External storage'. The file list contains the following items:

Name	Size	Modified
documents	34 kB	6 months ago
green	119.3 GB	7 months ago
manneback	Pending	2 minutes ago
music	3.6 MB	9 months ago
photos	45.8 MB	6 months ago
storage	<1 kB	7 months ago
duplicity-full-signatures.20150204T143855Z.sigtar.gpg	<1 kB	7 months ago
New Document.odt	8 kB	6 months ago
ownCloudUserManual.pdf	1.7 MB	9 months ago
slides.tgz	634 kB	9 months ago

At the bottom of the browser window, a status bar indicates: "Loading 'https://hall.cism.ucl.ac.be/owncloud/index.php/apps/files/', completed 16 of 17 items".

My home on Manneback



A screenshot of a web browser displaying the ownCloud file manager interface. The browser address bar shows the URL: https://hall.cism.ucl.ac.be/owncloud/index.php/apps/files/?dir=%2Fmanneback. The page title is 'manneback - Files - ownCloud'. The interface shows a sidebar on the left with navigation options: 'All files', 'Shared with you', 'Shared with others', 'Shared by link', 'External storage', and 'Deleted files'. The main content area shows a directory listing for the 'manneback' folder, owned by 'fclapuyt'. The listing includes folders like 'fortran', 'gpu', 'HardwareAccelerators', 'helloworld', 'hmem', 'hybrid', 'intel', 'IntelSample', 'jhead-2.97', and 'LoopTools-2.12'. A terminal window is overlaid on the bottom right of the screenshot, showing the command 'ls -ld [fgh]*/' and its output.

```
dfr@manneback:~$ ls -ld [fgh]*/
drwxr-xr-x 9 dfr grppan 4096 Jan 28 2015 fffclapuyt/
drwxr-xr-x 2 dfr grppan  49 Feb  1 2013 fortran/
drwxr-xr-x 3 dfr grppan  51 Aug 28 2013 gpu/
drwxrwxr-x 2 dfr grppan 4096 Dec 12 2013 helloworld/
drwxr-xr-x 2 dfr grppan   6 Mar 16 2015 hmem/
drwxr-xr-x 2 dfr grppan  43 Nov  8 2013 hybrid/
```

Can create a share URL



The screenshot shows the ownCloud web interface. The browser address bar displays the URL: <https://hall.cism.ucl.ac.be/owncloud/index.php/apps/files/?dir=%2Fmanneback>. The page title is "manneback - Files - ownCloud". The interface shows a file named "Acc... .tgz" (551.3 MB) with a share menu open. The share menu options are:

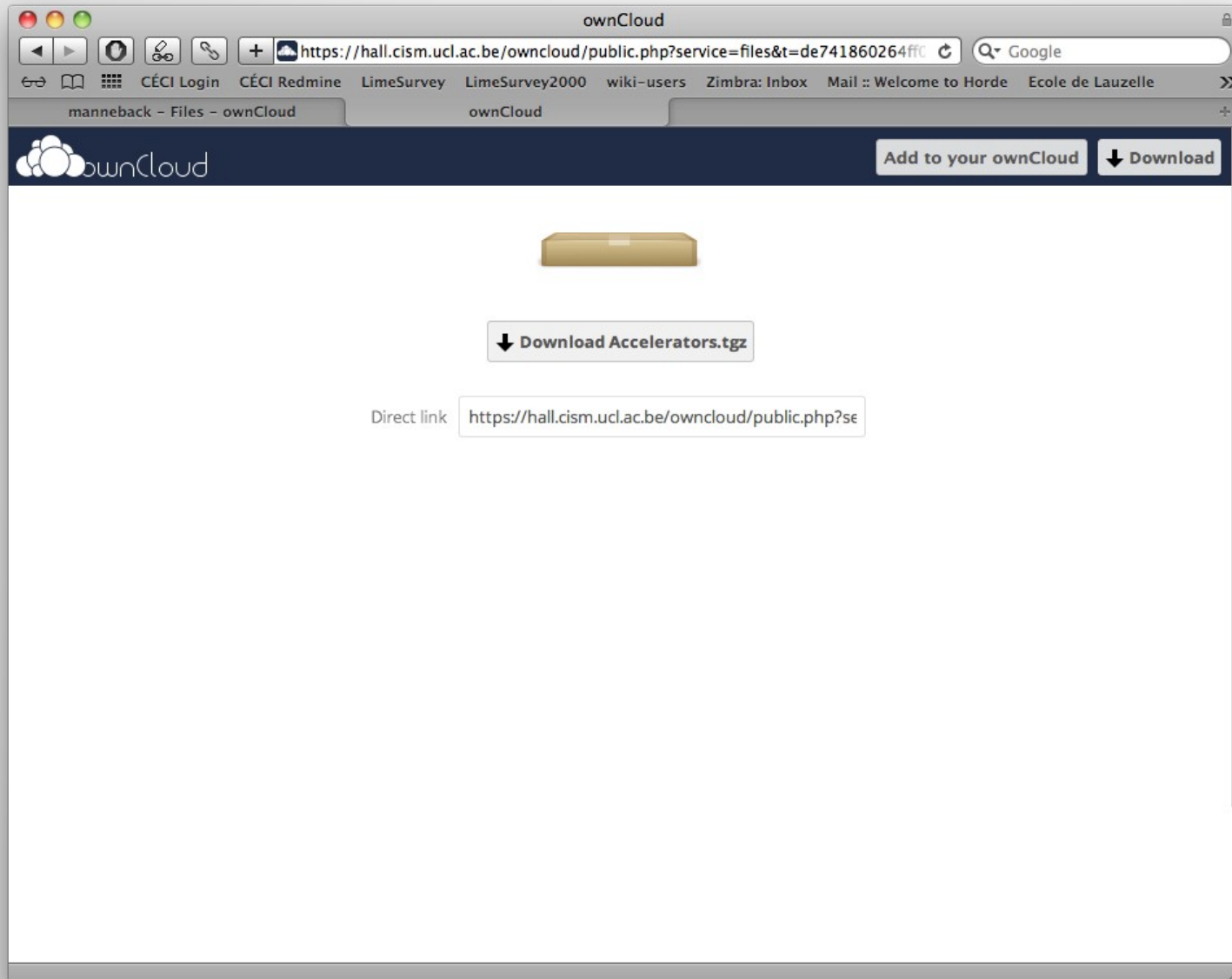
- Share with user or group ...
- Share link
- Password protect
- Set expiration date

The share link is: <https://hall.cism.ucl.ac.be/owncloud/public.php?service=files&t=>

Other files listed in the directory:

File Name	Size	Last Modified
a.sh	< 1 kB	last year
a.txt	< 1 kB	last year
Acc... .tgz	551.3 MB	years ago
arrays.c	< 1 kB	6 months ago
b.sh	< 1 kB	last year
b.txt	< 1 kB	last year

And distribute it



Open data – FAIR data



The screenshot shows the Dataverse UCLouvain website interface. At the top, there is a navigation bar with the Dataverse logo, search options, and user links. The main header features the UCLouvain logo and the text 'Open Data @ UCLouvain (uclouvain) Université catholique de Louvain'. Below this, a metrics bar shows '6 Downloads' and options for 'Contact' and 'Share'. A carousel of logos for various dataverses is displayed, including CISM, ELI, ICTEAM, and IMCN. A search bar is present with a 'Find' button and a link to 'Advanced Search'. The search results section shows '1 to 10 of 12 Results' and a 'Sort' dropdown. Two results are visible: 'Dataverse of SCEB (uclouvain)' dated Oct 7, 2019, and 'Dataverse of CP3 (IRMP)' dated Sep 5, 2019, with a link to 'Dataverse of IRMP'. The CP3 result includes a description: 'The UCL Centre for Cosmology, Particle Physics and Phenomenology (CP3) hosts research on high energy particle physics, cosmology, phenomenology and theory of the fundamental interactions. It is strong on both the experimental and theoretical fronts. The aim of the Centre is to br...'. On the left side, there are filters for 'Dataverses (9)', 'Datasets (3)', and 'Files (3)', along with sections for 'Dataverse Category', 'Publication Year', and 'Author Name'. A 'Display a menu' button is located at the bottom left of the search results area.

Open data – FAIR data



A screenshot of a web browser displaying the Dataverse website. The browser's address bar shows the URL 'dataverse.uclouvain.be/dataverse/cism'. The website header includes the 'Dataverse' logo and navigation links for 'Search', 'User Guide', 'Support', 'Sign Up', and 'Log In'. The main content area features the 'Dataverse of CISM' logo and the text 'Center for High Performance Computing and Mass Storage'. Below this, there is a search bar with the placeholder text 'Search this dataverse...', a 'Find' button, and a link to 'Advanced Search'. The search results section shows '1 to 1 of 1 Result' and a single entry titled 'Rapport d'activité 2018' by Keutgen, Thomas, dated Jun 26, 2019. The entry includes a DOI link: 'https://doi.org/10.5072/FK2/S31WN9'. On the left side of the search results, there are filters for 'Dataverses (0)', 'Datasets (1)', and 'Files (1)', as well as filters for 'Publication Year' (2019 (1)), 'Author Name' (Keutgen, Thomas (1)), 'Subject' (Other (1)), and 'Deposit Date' (2019 (1)).

Summary:



Storage: choose the right filesystem
and the right file format

Transfer: use the parallel tools when possible

Sharing: use all the potential of the UNIX
permissions and try Nextcloud and Dataverse